

The *Taylor Series* of an infinitely differentiable vector-valued function $y(t)$ of a scalar t is

$$y(t) = y(0) + t \cdot y'(0) + t^2 \cdot y''(0)/2 + t^3 \cdot y'''(0)/6 + t^4 \cdot y''''(0)/24 + \dots$$

These derivatives can be computed for a solution of the *Initial Value Problem*

$$“ y(0) = y_0 \text{ is given, and } y'(t) = f(y(t)) \text{ for all } t \geq 0 ”$$

from the derivatives of the given vector-valued function $f(y)$. In fact, from the *Chain Rule*,

$$y' = f, \quad y'' = f' \cdot f, \quad y''' = f'' \cdot f \cdot f + f'^2 \cdot f, \quad y'''' = f''' \cdot f \cdot f \cdot f + 3f'' \cdot f' \cdot f \cdot f + f' \cdot f'' \cdot f \cdot f + f'^3 \cdot f, \quad \dots$$

Then each derivative of $y(t)$ can be evaluated at $t = 0$ by evaluating each derivative of $f(y)$ at $y = y(0) = y_0$. Note that the higher derivatives of f are *symmetric multilinear operators*; for

instance, $f''(y)$ is a *symmetric bilinear* operator: $f'' \cdot u \cdot v = f'' \cdot v \cdot u$ is a vector-valued linear function of each vector u and v separately. Because linear operators do not necessarily

commute, $f'' \cdot f' \cdot f \cdot f \neq f' \cdot f'' \cdot f \cdot f$ in general, though they are equal if y 's vector space is one-dimensional. If y 's vector space is N -dimensional, then y and f can be represented by

column vectors each with N components; f' by a matrix with N^2 components; f'' by an array with N^3 components of which at most $(N+1)N^2/2$ can be distinct; f''' by an array of N^4 components ... Higher derivatives' arrays become huge when N is large.

Normally the Taylor series would be used to obtain $y(t+h)$ from $y(t)$ for any sufficiently small stepsize h :

$$y(t+h) = y(t) + h \cdot y'(t) + h^2 \cdot y''(t)/2 + h^3 \cdot y'''(t)/6 + h^4 \cdot y''''(t)/24 + \dots,$$

in which the derivatives $dy(t+h)/dh$ etc. are computed at $h = 0$ from the same formulas

$$y' = f, \quad y'' = f' \cdot f, \quad y''' = f'' \cdot f \cdot f + f'^2 \cdot f, \quad y'''' = f''' \cdot f \cdot f \cdot f + 3f'' \cdot f' \cdot f \cdot f + f' \cdot f'' \cdot f \cdot f + f'^3 \cdot f, \quad \dots$$

as before except that now $f(y)$ and its derivatives are computed at $y = y(t)$.

A similar process generates a formal series for any one-step numerical method's formula that advances an approximate solution $y = y(t)$ through one step h to $Y = Y(t+h)$, but now we differentiate with respect to h instead of t to get

$$Y(t+h) = y + h \cdot Y' + h^2 \cdot Y''/2 + h^3 \cdot Y'''/6 + h^4 \cdot Y''''/24 + \dots$$

in which $y = y(t)$ and the derivatives Y' etc. are derivatives of $Y(t+h)$ with respect to h evaluated at $h = 0$. These derivatives depend upon the numerical method's formula. For

example, take the (implicit) *Trapezoidal Rule* $Y = y + h \cdot (f(y) + f(Y))/2$. Now, at $t+h$,

$$Y' = (f(y) + f(Y))/2 + h \cdot f'(Y) \cdot Y'/2,$$

$$Y'' = f'(Y) \cdot Y' + h \cdot (f''(Y) \cdot Y' \cdot Y' + f'(Y) \cdot Y'')/2,$$

$$Y''' = 3f''(Y) \cdot Y' \cdot Y'/2 + f'(Y) \cdot Y''/2 + h \cdot (f'''(Y) \cdot Y' \cdot Y' \cdot Y' + \dots)/2, \quad \text{etc.}$$

Here every instance of Y or its derivatives is evaluated at $t+h$. For the Taylor series we set $h = 0$ in the foregoing formulas and substitute for derivatives of Y in succession to get

$$Y' = f, \quad Y'' = f' \cdot f, \quad Y''' = (3f'' \cdot f + f'^2) \cdot f/2, \quad \dots$$

in which now the derivatives of Y are evaluated at t , and f and its derivatives are evaluated at $y(t)$. Hence, the computed approximation

$$Y(t+h) = y(t) + h \cdot f + h^2 \cdot f' \cdot f/2 + h^3 \cdot (3f'' \cdot f + f'^2) \cdot f/12 + \dots$$

can be compared with the local solution

$$y(t+h) = y(t) + h \cdot f + h^2 \cdot f' \cdot f/2 + h^3 \cdot (f'' \cdot f + f'^2) \cdot f/6 + \dots$$

to reveal the 2nd-order Trapezoidal Rule's local truncation error

$$y(t+h) - Y(t+h) = h^3 \cdot (f'^2 - f'' \cdot f) \cdot f/12 + \dots$$

Another example is the (implicit) *Midpoint Rule* $Y = y + h \cdot f((y+Y)/2)$. At $t+h$,
 $Y' = f((y+Y)/2) + h \cdot f'((y+Y)/2) \cdot Y'/2$,
 $Y'' = f'((y+Y)/2) \cdot Y' + h \cdot (f''((y+Y)/2) \cdot Y' \cdot Y'/4 + f'((y+Y)/2) \cdot Y''/2)$,
 $Y''' = 3f''((y+Y)/2) \cdot Y' \cdot Y'/4 + 3f'((y+Y)/2) \cdot Y''/2 + h \cdot (f''' \dots)$, etc.

Setting $h = 0$ and substituting for derivatives of Y in succession yields

$$Y' = f, \quad Y'' = f' \cdot f, \quad Y''' = (3f'' \cdot f + 6f'^2) \cdot f/4, \quad \dots$$

in which now the derivatives of Y are evaluated at t , and f and its derivatives are evaluated at $y(t)$. Hence, by comparison with the Taylor series for $y(t+h)$ we find the 2nd-order Midpoint Rule's local truncation error to be

$$y(t+h) - Y(t+h) = h^3 \cdot (f'' \cdot f - 2f'^2) \cdot f/24 + \dots$$

The foregoing manipulations are tedious enough that only a computerized algebra system should perform them. However, programming *Maple* or *Mathematica* or *Derive* to perform them has been more difficult than it should be. At least some of the difficulty arises, I think, because these languages disallow declaration of a variable's linguistic *Type*. Besides the derivatives' non-commutative partially associative multiplication, their multi-linear symmetry has to be taken into account in order to achieve correct simplifications of expressions like the fifth derivative

$$y^{(5)} = f'''' \cdot f \cdot f \cdot f \cdot f + 6f'''' \cdot f' \cdot f \cdot f \cdot f + 4f'' \cdot f'' \cdot f \cdot f \cdot f + 4f'' \cdot f'^2 \cdot f \cdot f + \\ + 3f'' \cdot f' \cdot f \cdot f \cdot f + f' \cdot f'''' \cdot f \cdot f \cdot f + 3f' \cdot f'' \cdot f' \cdot f \cdot f + f'^2 \cdot f'' \cdot f \cdot f + f'^4 \cdot f.$$

(I hope I've gotten it right.)

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The foregoing *formal* (because their convergence is undetermined) Taylor series expansions do not reveal an important property possessed by the computed solution Y of the Initial Value Problem when it is obtained from a *Reflexive* formula, which is a formula in which $Y(t+h)$ and $y(t)$ are merely swapped when the sign of h is reversed. The Midpoint and Trapezoidal Rules' formulas are reflexive. The composition of T/h steps of a reflexive formula to approximate the true solution $y(T)$ at a fixed T , but using any sufficiently small stepsize h so long as T/h is an integer, can be proved to produce a computed approximation $Y(T)$ that depends upon h and differs from $y(T)$ by an error

$$y(T) - Y(T) = c_2 h^2 + c_4 h^4 + c_6 h^6 + \dots$$

whose formal expansion in powers of h contains only even powers. The expansion need not converge for any $h > 0$; instead it is an *Asymptotic* expansion whose behavior is conveyed by its first few terms ever more accurately as $h \rightarrow 0$. Recomputations of $Y(T)$ with diminishing stepsizes $h, h/2, h/4, h/8, \dots$ provides a sequence to which *Richardson's Extrapolation* can be applied, as in *Romberg* integration, to achieve what amounts to higher-order convergence.