Graphical models, message-passing algorithms, and convex optimization

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Tutorial slides based on joint paper with Michael Jordan
Paper at: www.eecs.berkeley.edu/~wainwrig/WaiJorVariational03.ps
Introduction

• graphical models are used and studied in various applied statistical and computational fields:
  – machine learning and artificial intelligence
  – computational biology
  – statistical signal/image processing
  – communication and information theory
  – statistical physics
  – ..... 

• based on correspondences between graph theory and probability theory

• important but difficult problems:
  – computing likelihoods, marginal distributions, modes
  – estimating model parameters and structure from (noisy) data
Outline

1. Introduction and motivation
   (a) Background on graphical models
   (b) Some applications and challenging problems
   (c) Illustrations of some message-passing algorithms

2. Exponential families and variational methods
   (a) What is a variational method (and why should I care)?
   (b) Graphical models as exponential families
   (c) Variational representations from conjugate duality

3. Exact techniques as variational methods
   (a) Gaussian inference on arbitrary graphs
   (b) Belief-propagation/sum-product on trees (e.g., Kalman filter; $\alpha$-$\beta$ alg.)
   (c) Max-product on trees (e.g., Viterbi)

4. Approximate techniques as variational methods
   (a) Mean field and variants
   (b) Belief propagation and extensions on graphs with cycles
   (c) Semidefinite constraints and convex relaxations
Undirected graphical models

Based on correspondences between graphs and random variables.

- given an undirected graph $G = (V, E)$, each node $s$ has an associated random variable $X_s$
- for each subset $A \subseteq V$, define $X_A := \{X_s, s \in A\}$.

Maximal cliques (123), (345), (456), (47)  

- a clique $C \subseteq V$ is a subset of vertices all joined by edges  
- a vertex cutset is a subset $S \subset V$ whose removal breaks the graph into two or more pieces
Factorization and Markov properties

The graph $G$ can be used to impose constraints on the random vector $X = X_V$ (or on the distribution $p$) in different ways.

**Markov property:** $X$ is Markov w.r.t $G$ if $X_A$ and $X_B$ are conditionally indpt. given $X_S$ whenever $S$ separates $A$ and $B$.

**Factorization:** The distribution $p$ factorizes according to $G$ if it can be expressed as a product over cliques:

$$p(x) = \frac{1}{Z} \prod_{C \in C} \exp \{ \theta_C(x_C) \}$$

compatibility function on clique $C$

**Theorem:** (Hammersley-Clifford) For strictly positive $p(\cdot)$, the Markov property and the Factorization property are equivalent.
Example: Hidden Markov models

(a) Hidden Markov model

(b) Coupled HMM

- HMMs are widely used in various applications
  - discrete $X_t$: computational biology, speech processing, etc.
  - Gaussian $X_t$: control theory, signal processing, etc.

- frequently wish to solve *smoothing* problem of computing
  $p(x_t \mid y_1, \ldots, y_T)$

- exact computation in HMMs is tractable, but coupled HMMs require
  algorithms for approximate computation (e.g., structured mean field)
Example: Statistical signal and image processing

(a) Natural image  
(b) Lattice  
(c) Multiscale quadtree

- frequently wish to compute log likelihoods (e.g., for classification), or marginalss/modes (e.g., for denoising, deblurring, de-convolution, coding)

- exact algorithms available for tree-structured models; approximate techniques (e.g., belief propagation and variants) required for more complex models
Example: Graphical codes for communication

**Goal:** Achieve reliable communication over a noisy channel.

- wide variety of applications: satellite communication, sensor networks, computer memory, neural communication
- error-control codes based on careful addition of redundancy, with their fundamental limits determined by Shannon theory
- key implementational issues: efficient construction, encoding and decoding
- very active area of current research: graphical codes (e.g., turbo codes, low-density parity check codes) and iterative message-passing algorithms (belief propagation; max-product)
Graphical codes and decoding (continued)

Parity check matrix

\[ H = \begin{bmatrix}
1 & 0 & 1 & 0 & 1 & 0 & 1 \\
0 & 1 & 1 & 0 & 0 & 1 & 1 \\
0 & 0 & 0 & 1 & 1 & 1 & 1 \\
\end{bmatrix} \]

Codeword: \[ [0101010] \]
Non-codeword: \[ [0000011] \]

- **Decoding**: requires finding maximum likelihood codeword:

\[ \hat{x}_{ML} = \arg \max_x p(y | x) \quad \text{s.t.} \quad Hx = 0 \pmod{2}. \]

- use of belief propagation as an approximate decoder has revolutionized the field of error-control coding
Example: Computer vision

- disparity for stereo vision: estimate depth in scenes based on two (or more) images taken from different positions
- *global approaches*: disparity map based on optimization in an MRF

\[
\begin{align*}
\theta_t(d_t) & \quad \theta_{st}(d_s, d_t) \\
& \quad \theta_s(d_s)
\end{align*}
\]

- grid-structured graph \( G = (V, E) \)
- \( d_s \equiv \) disparity at grid position \( s \)
- \( \theta_s(d_s) \equiv \) image data fidelity term
- \( \theta_{st}(d_s, d_t) \equiv \) disparity coupling

- optimal disparity map \( \tilde{d} \) found by solving MAP estimation problem for this Markov random field

- computationally intractable (NP-hard) in general, but iterative message-passing algorithms (e.g., belief propagation) solve many practical instances
Challenging computational problems

Frequently, it is of interest to compute various quantities associated with an undirected graphical model:

(a) the log normalization constant $\log Z$

(b) local marginal distributions or other local statistics

(c) modes or most probable configurations

Relevant dimensions often grow rapidly in graph size $\rightarrow$ major computational challenges.

**Example:** Consider a naive approach to computing the normalization constant for binary random variables:

$$Z = \sum_{\mathbf{x} \in \{0,1\}^n} \prod_{C \in \mathcal{C}} \exp \{ \theta_C(x_C) \}$$

Complexity scales exponentially as $2^n$. 
Gibbs sampling in the Ising model

- binary variables on a graph $G = (V, E)$ with pairwise interactions:

$$p(x; \theta) \propto \exp \left\{ \sum_{s \in V} \theta_s x_s + \sum_{(s,t) \in E} \theta_{st} x_s x_t \right\}$$

- Update $x^{(m+1)}_s$ stochastically based on values $x^{(m)}_{\mathcal{N}(s)}$ at neighbors:

1. Choose $s \in V$ at random.
2. Sample $u \sim U(0,1)$ and update

$$x^{(m+1)}_s = \begin{cases} 1 & \text{if } u \leq \left\{ 1 + \exp[-(\theta_s + \sum_{t \in \mathcal{N}(s)} \theta_{st} x^{(m)}_t)] \right\}^{-1} \\ 0 & \text{otherwise} \end{cases}$$

- sequence $\{x^{(m)}\}$ converges (in a stochastic sense) to a sample from $p(x; \theta)$
Mean field updates in the Ising model

- binary variables on a graph $G = (V, E)$ with pairwise interactions:

\[ p(x; \theta) \propto \exp \left\{ \sum_{s \in V} \theta_s x_s + \sum_{(s,t) \in E} \theta_{st} x_s x_t \right\} \]

- simple (deterministic) message-passing algorithm involving variational parameters $\nu_s \in (0, 1)$ at each node

```
1. Choose $s \in V$ at random.
2. Update $\nu_s$ based on neighbors $\{\nu_t, t \in \mathcal{N}(s)\}$:

\[ \nu_s \leftarrow \left\{ 1 + \exp\left[ - \left( \theta_s + \sum_{t \in \mathcal{N}(s)} \theta_{st} \nu_t \right) \right] \right\}^{-1} \]
```

Questions:

- principled derivation?
- convergence and accuracy?
Sum and max-product algorithms: On trees

Exact for trees, but approximate for graphs with cycles.

\[ M_{ts} \equiv \text{message from node } t \text{ to } s \]
\[ \Gamma(t) \equiv \text{neighbors of node } t \]

**Sum-product**: for marginals

(generalizes \( \alpha - \beta \) algorithm; Kalman filter)

**Max-product**: for MAP configurations

(generalizes Viterbi algorithm)

\[
\text{Update: } M_{ts}(x_s) \leftarrow \sum_{x'_t \in \mathcal{X}_t} \left\{ \exp \left[ \theta_{st}(x_s, x'_t) + \theta_t(x'_t) \right] \prod_{v \in \Gamma(t) \setminus s} M_{vt}(x_t) \right\}
\]

\[
\text{Marginals: } p(x_s; \theta) \propto \exp\{\theta_t(x_t)\} \prod_{t \in \Gamma(s)} M_{ts}(x_s).
\]
Sum and max-product: On graphs with cycles

- what about applying same updates on graph with cycles?
- updates need not converge (effect of cycles)
- seems naive, but remarkably successful in many applications

Questions:
- meaning of these updates for graphs with cycles?
- convergence? accuracy of resulting “marginals”?

$M_{ts} \equiv$ message from node $t$ to $s$
$\Gamma(t) \equiv$ neighbors of node $t$

<table>
<thead>
<tr>
<th>Sum-product:</th>
<th>Max-product:</th>
</tr>
</thead>
<tbody>
<tr>
<td>for marginals</td>
<td>for modes</td>
</tr>
</tbody>
</table>

\[ T_w \]
\[ T_u \]
\[ T_v \]
\[ M_{wt} \]
\[ M_{ut} \]
\[ M_{vt} \]
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Variational methods

- "variational": umbrella term for optimization-based formulation of problems, and methods for their solution
- historical roots in the calculus of variations
- modern variational methods encompass a wider class of methods (e.g., dynamic programming; finite-element methods)

Variational principle: Representation of a quantity of interest \( \hat{u} \) as the solution of an optimization problem.

1. allows the quantity \( \hat{u} \) to be studied through the lens of the optimization problem
2. approximations to \( \hat{u} \) can be obtained by approximating or relaxing the variational principle
Illustration: A simple variational principle

Goal: Given a vector $\mathbf{y} \in \mathbb{R}^n$ and a symmetric matrix $Q > 0$, solve the linear system $Q \mathbf{u} = \mathbf{y}$.

Unique solution $\hat{\mathbf{u}}(\mathbf{y}) = Q^{-1} \mathbf{y}$ can be obtained by matrix inversion.

Variational formulation: Consider the function $J_\mathbf{y} : \mathbb{R}^n \to \mathbb{R}$ defined by

$$J_\mathbf{y}(\mathbf{u}) := \frac{1}{2} \mathbf{u}^T Q \mathbf{u} - \mathbf{y}^T \mathbf{u}.$$ 

It is strictly convex, and the minimum is uniquely attained:

$$\hat{\mathbf{u}}(\mathbf{y}) = \arg \min_{\mathbf{u} \in \mathbb{R}^n} J_\mathbf{y}(\mathbf{u}) = Q^{-1} \mathbf{y}.$$ 

Various methods for solving linear systems (e.g., conjugate gradient) exploit this variational representation.
Useful variational principles for graphical models?

Consider an undirected graphical model:

\[
p(x) = \frac{1}{Z} \prod_{C \in \mathcal{C}} \exp \{ \theta_C(x_C) \}.
\]

Core problems that arise in many applications:

(a) computing the log normalization constant \( \log Z \)

(b) computing local marginal distributions (e.g., \( p(x_s) = \sum_{x_t, t \neq s} p(x) \))

(c) computing modes or most likely configurations \( \hat{x} \in \arg \max_x p(x) \)

Approach: Develop variational representations of all of these problems by exploiting results from:

(a) exponential families \( \text{ (e.g., Brown, 1986) } \)

(b) convex duality \( \text{ (e.g., Rockafellar, 1973) } \)
Maximum entropy formulation of graphical models

- suppose that we have measurements $\hat{\mu}$ of the average values of some (local) functions $\phi_\alpha : \mathcal{X}^n \to \mathbb{R}$
- in general, will be many distributions $p$ that satisfy the measurement constraints $\mathbb{E}_p[\phi_\alpha(x)] = \hat{\mu}$
- will consider finding the $p$ with maximum “uncertainty” subject to the observations, with uncertainty measured by entropy

$$H(p) = - \sum_x p(x) \log p(x).$$

**Constrained maximum entropy problem:** Find $\hat{p}$ to solve

$$\max_{p \in \mathcal{P}} H(p) \quad \text{such that} \quad \mathbb{E}_p[\phi_\alpha(x)] = \hat{\mu}$$

- elementary argument with Lagrange multipliers shows that solution takes the exponential form

$$\hat{p}(x; \theta) \propto \exp \left\{ \sum_{\alpha \in \mathcal{I}} \theta_\alpha \phi_\alpha(x) \right\}.$$
Exponential families

\[ \phi_\alpha : \mathcal{X}^n \to \mathbb{R} \quad \equiv \quad \text{sufficient statistic} \]

\[ \phi = \{ \phi_\alpha, \alpha \in \mathcal{I} \} \quad \equiv \quad \text{vector of sufficient statistics} \]

\[ \theta = \{ \theta_\alpha, \alpha \in \mathcal{I} \} \quad \equiv \quad \text{parameter vector} \]

\[ \nu \quad \equiv \quad \text{base measure (e.g., Lebesgue, counting)} \]

- parameterized family of densities (w.r.t. \( \nu \)):

\[
p(x; \theta) = \exp \left\{ \sum_\alpha \theta_\alpha \phi_\alpha(x) - A(\theta) \right\}
\]

- cumulant generating function (log normalization constant):

\[
A(\theta) = \log \left( \int \exp \{ \langle \theta, \phi(x) \rangle \} \nu(dx) \right)
\]

- set of valid parameters \( \Theta := \{ \theta \in \mathbb{R}^d | A(\theta) < +\infty \} \).

- will focus on regular families for which \( \Theta \) is open.
Examples: Scalar exponential families

<table>
<thead>
<tr>
<th>Family</th>
<th>$\mathcal{X}$</th>
<th>$\nu$</th>
<th>$\log p(x; \theta)$</th>
<th>$A(\theta)$</th>
</tr>
</thead>
<tbody>
<tr>
<td>Bernoulli</td>
<td>${0, 1}$</td>
<td>Counting</td>
<td>$\theta x - A(\theta)$</td>
<td>$\log[1 + \exp(\theta)]$</td>
</tr>
<tr>
<td>Gaussian</td>
<td>$\mathbb{R}$</td>
<td>Lebesgue</td>
<td>$\theta_1 x + \theta_2 x^2 - A(\theta)$</td>
<td>$\frac{1}{2} [\theta_1 + \log \frac{2\pi e}{-\theta_2}]$</td>
</tr>
<tr>
<td>Exponential</td>
<td>$(0, +\infty)$</td>
<td>Lebesgue</td>
<td>$\theta (-x) - A(\theta)$</td>
<td>$-\log \theta$</td>
</tr>
<tr>
<td>Poisson</td>
<td>${0, 1, 2 \ldots}$</td>
<td>Counting (h(x) = 1/x!)</td>
<td>$\theta x - A(\theta)$</td>
<td>$\exp(\theta)$</td>
</tr>
</tbody>
</table>
Graphical models as exponential families

- choose random variables $X_s$ at each vertex $s \in V$ from an arbitrary exponential family (e.g., Bernoulli, Gaussian, Dirichlet etc.)
- exponential family can be the same at each node (e.g., multivariate Gaussian), or different (e.g., mixture models).

Key requirement: The collection $\phi$ of sufficient statistics must respect the structure of $G$. 
Example: Discrete Markov random field

\[ \theta_t(x_t) \rightarrow \theta_{st}(x_s, x_t) \rightarrow \theta_s(x_s) \]

Indicators:
\[ \mathbb{I}_j(x_s) = \begin{cases} 1 & \text{if } x_s = j \\ 0 & \text{otherwise} \end{cases} \]

Parameters:
\[ \theta_s = \{\theta_{s;j}, j \in \mathcal{X}_s\} \]
\[ \theta_{st} = \{\theta_{st;jk}, (j, k) \in \mathcal{X}_s \times \mathcal{X}_t\} \]

Compact form:
\[ \theta_s(x_s) := \sum_j \theta_{s;j} \mathbb{I}_j(x_s) \]
\[ \theta_{st}(x_s, x_t) := \sum_{j,k} \theta_{st;jk} \mathbb{I}_j(x_s) \mathbb{I}_k(x_t) \]

Density (w.r.t. counting measure) of the form:
\[ p(x; \theta) \propto \exp \left\{ \sum_{s \in V} \theta_s(x_s) + \sum_{(s,t) \in E} \theta_{st}(x_s, x_t) \right\} \]

Cumulant generating function (log normalization constant):
\[ A(\theta) = \log \sum_{x \in \mathcal{X}^n} \exp \left\{ \sum_{s \in V} \theta_s(x_s) + \sum_{(s,t) \in E} \theta_{st}(x_s, x_t) \right\} \]
Special case: Hidden Markov model

- Markov chain \( \{X_1, X_2, \ldots\} \) evolving in time, with noisy observation \( Y_t \) at each time \( t \)

\[
\begin{align*}
\theta_{23}(x_2, x_3) &= \log p(x_3 | x_2) \\
\theta_5(x_5) &= \log p(y_5 | x_5)
\end{align*}
\]

- an HMM is a particular type of discrete MRF, representing the conditional \( p(x | y; \theta) \)

- exponential parameters have a concrete interpretation

- the cumulant generating function \( A(\theta) \) is equal to the log likelihood 

\[
\log p(y; \theta)
\]
**Example: Multivariate Gaussian**

$U(\theta)$: Matrix of natural parameters  \hspace{1cm} $\phi(\mathbf{x})$: Matrix of sufficient statistics

\[
\begin{bmatrix}
0 & \theta_1 & \theta_2 & \ldots & \theta_n \\
\theta_1 & \theta_{11} & \theta_{12} & \ldots & \theta_{1n} \\
\theta_2 & \theta_{21} & \theta_{22} & \ldots & \theta_{2n} \\
\vdots & \vdots & \vdots & \ddots & \vdots \\
\theta_n & \theta_{n1} & \theta_{n2} & \ldots & \theta_{nn}
\end{bmatrix}
\begin{bmatrix}
1 & x_1 & x_2 & \ldots & x_n \\
x_1 & (x_1)^2 & x_1x_2 & \ldots & x_1x_n \\
x_2 & x_2x_1 & (x_2)^2 & \ldots & x_2x_n \\
\vdots & \vdots & \vdots & \ddots & \vdots \\
x_n & x_nx_1 & x_nx_2 & \ldots & (x_n)^2
\end{bmatrix}
\]

Edgewise natural parameters $\theta_{st} = \theta_{ts}$ must respect graph structure:

(a) Graph structure \hspace{1cm} (b) Structure of $[Z(\theta)]_{st} = \theta_{st}$. 
Example: Mixture of Gaussians

- can form mixture models by combining different types of random variables
- let $Y_s$ be conditionally Gaussian given the discrete variable $X_s$ with parameters $\gamma_{s;j} = (\mu_{s;j}, \sigma^2_{s;j})$:
  \[ X_s \overset{\circ}{\rightharpoonup} p(x_s; \theta_s) \]
  \[ p(y_s \mid x_s; \gamma_s) \]
  \[ Y_s \overset{\bullet}{\rightharpoonup} \]
  \[ X_s \equiv \text{mixture indicator} \]
  \[ Y_s \equiv \text{mixture of Gaussian} \]
- couple the mixture indicators $X = \{X_s, s \in V\}$ using a discrete MRF
- overall model has the exponential form

$$p(y, x; \theta, \gamma) \propto \prod_{s \in V} p(y_s \mid x_s; \gamma_s) \exp \left\{ \sum_{s \in V} \theta_s(x_s) + \sum_{(s,t) \in E} \theta_{st}(x_s, x_t) \right\}.$$
Conjugate dual functions

- conjugate duality is a fertile source of variational representations
- any function $f$ can be used to define another function $f^*$ as follows:

$$f^*(v) := \sup_{u \in \mathbb{R}^n} \{ \langle v, u \rangle - f(u) \}.$$  

- easy to show that $f^*$ is always a convex function
- how about taking the “dual of the dual”? I.e., what is $(f^*)^*$?
- when $f$ is well-behaved (convex and lower semi-continuous), we have $(f^*)^* = f$, or alternatively stated:

$$f(u) = \sup_{v \in \mathbb{R}^n} \{ \langle u, v \rangle - f^*(v) \}.$$
**Geometric view: Supporting hyperplanes**

**Question:** Given all hyperplanes in $\mathbb{R}^n \times \mathbb{R}$ with normal $(v, -1)$, what is the intercept of the one that supports $\text{epi}(f)$?

**Epigraph of $f$:**

$$\text{epi}(f) := \{(u, \beta) \in \mathbb{R}^{n+1} \mid f(u) \leq \beta\}.$$

Analytically, we require the smallest $c \in \mathbb{R}$ such that:

$$\langle v, u \rangle - c \leq f(u) \quad \text{for all} \quad u \in \mathbb{R}^n$$

By re-arranging, we find that this optimal $c^*$ is the dual value:

$$c^* = \sup_{u \in \mathbb{R}^n} \{\langle v, u \rangle - f(u)\}.$$
Example: Single Bernoulli

Random variable $X \in \{0, 1\}$ yields exponential family of the form:

$$p(x; \theta) \propto \exp \{\theta x\} \quad \text{with} \quad A(\theta) = \log [1 + \exp(\theta)].$$

Let’s compute the dual $A^*(\mu) := \sup_{\theta \in \mathbb{R}} \{\mu \theta - \log[1 + \exp(\theta)]\}$.

(Possible) stationary point: $\mu = \exp(\theta)/[1 + \exp(\theta)]$.

We find that:

$$A^*(\mu) = \begin{cases} 
\mu \log \mu + (1 - \mu) \log(1 - \mu) & \text{if } \mu \in [0, 1], \\
+\infty & \text{otherwise.}
\end{cases}$$

Leads to the variational representation:

$$A(\theta) = \max_{\mu \in [0,1]} \left\{ \mu \cdot \theta - A^*(\mu) \right\}.$$
More general computation of the dual $A^*$

- consider the definition of the dual function:

$$A^*(\mu) = \sup_{\theta \in \mathbb{R}^d} \{ \langle \mu, \theta \rangle - A(\theta) \}.$$ 

- taking derivatives w.r.t $\theta$ to find a stationary point yields:

$$\mu - \nabla A(\theta) = 0.$$ 

- **Useful fact:** Derivatives of $A$ yield *mean parameters*:

$$\frac{\partial A}{\partial \theta_\alpha} (\theta) = \mathbb{E}_\theta [\phi_\alpha(x)] := \int \phi_\alpha(x)p(x; \theta) \nu(x).$$

Thus, stationary points satisfy the equation:

$$\mu = \mathbb{E}_\theta [\phi(x)] \quad (1)$$
Computation of dual (continued)

• assume solution \(\theta(\mu)\) to equation (1) exists

• strict concavity of objective guarantees that \(\theta(\mu)\) attains global maximum with value

\[
A^*(\mu) = \langle \mu, \theta(\mu) \rangle - A(\theta(\mu))
= \mathbb{E}_{\theta(\mu)} \left[ \langle \theta(\mu), \phi(x) \rangle - A(\theta(\mu)) \right]
= \mathbb{E}_{\theta(\mu)}[\log p(x; \theta(\mu))]
\]

• recall the definition of entropy:

\[
H(p(x)) := -\int \left[ \log p(x) \right] p(x) \nu(dx)
\]

• thus, we recognize that \(A^*(\mu) = -H(p(x; \theta(\mu)))\) when equation (1) has a solution

**Question:** For which \(\mu \in \mathbb{R}^d\) does equation (1) have a solution \(\theta(\mu)\)?
Sets of realizable mean parameters

- for any distribution \( p(\cdot) \), define a vector \( \mu \in \mathbb{R}^d \) of mean parameters:

\[
\mu_\alpha := \int \phi_\alpha(x)p(x)\nu(dx)
\]

- now consider the set \( \mathcal{M}(G; \phi) \) of all realizable mean parameters:

\[
\mathcal{M}(G; \phi) = \{ \mu \in \mathbb{R}^d \mid \mu_\alpha = \int \phi_\alpha(x)p(x)\nu(dx) \text{ for some } p(\cdot) \}\]

- for discrete families, we refer to this set as a marginal polytope, denoted by \( \text{MARG}(G; \phi) \)
Examples of $\mathcal{M}$: Gaussian MRF

$\phi(x)$ Matrix of sufficient statistics

$$
\begin{bmatrix}
1 & x_1 & x_2 & \ldots & x_n \\
x_1 & (x_1)^2 & x_1x_2 & \ldots & x_1x_n \\
x_2 & x_2x_1 & (x_2)^2 & \ldots & x_2x_n \\
\vdots & \vdots & \vdots & \ddots & \vdots \\
x_n & x_nx_1 & x_nx_2 & \ldots & (x_n)^2
\end{bmatrix}
$$

$U(\mu)$ Matrix of mean parameters

$$
\begin{bmatrix}
1 & \mu_1 & \mu_2 & \ldots & \mu_n \\
\mu_1 & \mu_{11} & \mu_{12} & \ldots & \mu_{1n} \\
\mu_2 & \mu_{21} & \mu_{22} & \ldots & \mu_{2n} \\
\vdots & \vdots & \vdots & \ddots & \vdots \\
\mu_n & \mu_{n1} & \mu_{n2} & \ldots & \mu_{nn}
\end{bmatrix}
$$

- Gaussian mean parameters are specified by a single semidefinite constraint as $\mathcal{M}_{Gauss} = \{\mu \in \mathbb{R}^{n+\binom{n}{2}} | U(\mu) \succeq 0\}$.

Scalar case:

$$
U(\mu) = \begin{bmatrix}
1 & \mu_1 \\
\mu_1 & \mu_{11}
\end{bmatrix}
$$
Examples of $\mathcal{M}$: Discrete MRF

- **sufficient statistics:**
  \[
  \mathbb{I}_j(x_s) \quad \text{for } s = 1, \ldots, n, \quad j \in \mathcal{X}_s \\
  \mathbb{I}_{jk}(x_s, x_t) \quad \text{for } (s, t) \in E, \quad (j, k) \in \mathcal{X}_s \times \mathcal{X}_t
  \]

- **mean parameters** are simply marginal probabilities, represented as:
  \[
  \mu_s(x_s) := \sum_{j \in \mathcal{X}_s} \mu_{s;j} \mathbb{I}_j(x_s), \quad \mu_{st}(x_s, x_t) := \sum_{(j, k) \in \mathcal{X}_s \times \mathcal{X}_t} \mu_{st;jk} \mathbb{I}_{jk}(x_s, x_t)
  \]

- denote the set of realizable $\mu_s$ and $\mu_{st}$ by $\text{MARG}(G)$
- refer to it as the *marginal polytope*
- extremely difficult to characterize for general graphs
Geometry and moment mapping

For suitable classes of graphical models in exponential form, the gradient map $\nabla A$ is a bijection between $\Theta$ and the interior of $\mathcal{M}$.

(e.g., Brown, 1986; Efron, 1978)
Variational principle in terms of mean parameters

- The conjugate dual of $A$ takes the form:

$$A^*(\mu) = \begin{cases} -H(p(x; \theta(\mu))) & \text{if } \mu \in \text{int } \mathcal{M}(G; \phi) \\ +\infty & \text{if } \mu \notin \text{cl } \mathcal{M}(G; \phi). \end{cases}$$

**Interpretation:**
- $A^*(\mu)$ is finite (and equal to a certain negative entropy) for any $\mu$ that is globally realizable
- if $\mu \notin \text{cl } \mathcal{M}(G; \phi)$, then the max. entropy problem is *infeasible*

- The cumulant generating function $A$ has the representation:

$$A(\theta) = \sup_{\mu \in \mathcal{M}(G; \phi)} \{\langle \theta, \mu \rangle - A^*(\mu)\},$$

**cumulant generating func.** max. ent. problem over $\mathcal{M}$

- in contrast to the “free energy” approach, solving this problem provides both the value $A(\theta)$ and the exact mean parameters $\widehat{\mu}_\alpha = \mathbb{E}_\theta[\phi_\alpha(x)]$
Alternative view: Kullback-Leibler divergence

- Kullback-Leibler divergence defines “distance” between probability distributions:

\[
D(p \parallel q) := \int \left[ \log \frac{p(x)}{q(x)} \right] p(x) \nu(dx)
\]

- for two exponential family members \( p(x; \theta^1) \) and \( p(x; \theta^2) \), we have

\[
D(p(x; \theta^1) \parallel p(x; \theta^2)) = A(\theta^2) - A(\theta^1) - \langle \mu^1, \theta^2 - \theta^1 \rangle
\]

- substituting \( A(\theta^1) = \langle \theta^1, \mu^1 \rangle - A^*(\mu^1) \) yields a mixed form:

\[
D(p(x; \theta^1) \parallel p(x; \theta^2)) \equiv D(\mu^1 \parallel \theta^2) = A(\theta^2) + A^*(\mu^1) - \langle \mu^1, \theta^2 \rangle
\]

Hence, the following two assertions are equivalent:

\[
A(\theta^2) = \sup_{\mu^1 \in \mathcal{M}(G;\phi)} \{ \langle \theta^2, \mu^1 \rangle - A^*(\mu^1) \}
\]

\[
0 = \inf_{\mu^1 \in \mathcal{M}(G;\phi)} D(\mu^1 \parallel \theta^2)
\]
Challenges

1. In general, mean parameter spaces $\mathcal{M}$ can be very difficult to characterize (e.g., multidimensional moment problems).

2. Entropy $A^*(\mu)$ as a function of only the mean parameters $\mu$ typically lacks an explicit form.

Remarks:

1. Variational representation clarifies why certain models are tractable.

2. For intractable cases, one strategy is to solve an approximate form of the optimization problem.
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   (c) Illustrations of some message-passing algorithms

2. Exponential families and variational methods
   (a) What is a variational method (and why should I care)?
   (b) Graphical models as exponential families
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3. Exact techniques as variational methods
   (a) Gaussian inference on arbitrary graphs
   (b) Belief-propagation/sum-product on trees (e.g., Kalman filter; $\alpha$-$\beta$ alg.)
   (c) Max-product on trees (e.g., Viterbi)

4. Approximate techniques as variational methods
   (a) Mean field and variants
   (b) Belief propagation and extensions
   (c) Semidefinite constraints and convex relaxations
**A(i): Multivariate Gaussian (fixed covariance)**

Consider the set of all Gaussians with fixed *inverse* covariance $Q > 0$.

- potentials $\phi(x) = \{x_1, \ldots, x_n\}$ and natural parameter $\theta \in \Theta = \mathbb{R}^n$.
- cumulant generating function:

$$A(\theta) = \log \int_{\mathbb{R}^n} \exp \left\{ \sum_{s=1}^{n} \theta_s x_s \right\} \exp \left\{ -\frac{1}{2} x^T Q x \right\} dx$$

- completing the square yields $A(\theta) = \frac{1}{2} \theta^T Q^{-1} \theta + \text{constant}$
- straightforward computation leads to the dual

$$A^*(\mu) = \frac{1}{2} \mu^T Q \mu - \text{constant}$$

- putting the pieces back together yields the variational principle

$$A(\theta) = \sup_{\mu \in \mathbb{R}^n} \left\{ \theta^T \mu - \frac{1}{2} \mu^T Q \mu \right\} + \text{constant}$$

- optimum is uniquely obtained at the familiar Gaussian mean $\hat{\mu} = Q^{-1} \theta$. 
A(ii): Multivariate Gaussian (arbitrary covariance)

- matrices of sufficient statistics, natural parameters, and mean parameters:

\[
\begin{align*}
\phi(x) &= \begin{bmatrix} 1 \\ x \end{bmatrix} \begin{bmatrix} 1 & x \end{bmatrix}, & U(\theta) &= \begin{bmatrix} 0 & [\theta_s] \\ [\theta_s] & [\theta_{st}] \end{bmatrix} & U(\mu) &= \mathbb{E}\left\{ \begin{bmatrix} 1 \\ x \end{bmatrix} \begin{bmatrix} 1 & x \end{bmatrix} \right\}
\end{align*}
\]

- cumulant generating function:

\[
A(\theta) = \log \int \exp \left\{ \langle U(\theta), \phi(x) \rangle \right\} dx
\]

- computing the dual function:

\[
A^*(\mu) = -\frac{1}{2} \log \det U(\mu) - \frac{n}{2} \log 2\pi e,
\]

- exact variational principle is a log-determinant problem:

\[
A(\theta) = \sup_{U(\mu) \succ 0, [U(\mu)]_{11}=1} \left\{ \langle U(\theta), U(\mu) \rangle + \frac{1}{2} \log \det U(\mu) \right\} + \frac{n}{2} \log 2\pi e
\]

- solution yields the normal equations for Gaussian mean and covariance.
B: Belief propagation/sum-product on trees

- discrete variables $X_s \in \{0, 1, \ldots, m_s - 1\}$ on a tree $T = (V, E)$

- sufficient statistics: indicator functions for each node and edge
  \[
  \mathbb{I}_j(x_s) \quad \text{for} \quad s = 1, \ldots, n, \quad j \in \mathcal{X}_s
  \]
  \[
  \mathbb{I}_{jk}(x_s, x_t) \quad \text{for} \quad (s, t) \in E, \quad (j, k) \in \mathcal{X}_s \times \mathcal{X}_t.
  \]

- exponential representation of distribution:
  \[
  p(x; \theta) \propto \exp \left\{ \sum_{s \in V} \theta_s(x_s) + \sum_{(s,t) \in E} \theta_{st}(x_s, x_t) \right\}
  \]
  where $\theta_s(x_s) := \sum_{j \in \mathcal{X}_s} \theta_{s;j} \mathbb{I}_j(x_s)$ (and similarly for $\theta_{st}(x_s, x_t)$)

- mean parameters are simply marginal probabilities, represented as:
  \[
  \mu_s(x_s) := \sum_{j \in \mathcal{X}_s} \mu_{s;j} \mathbb{I}_j(x_s), \quad \mu_{st}(x_s, x_t) := \sum_{(j, k) \in \mathcal{X}_s \times \mathcal{X}_t} \mu_{st;jk} \mathbb{I}_{jk}(x_s, x_t)
  \]

- the marginals must belong to the following marginal polytope:
  \[
  \text{MARG}(T) := \{ \mu \geq 0 \mid \sum_{x_s} \mu_s(x_s) = 1, \sum_{x_t} \mu_{st}(x_s, x_t) = \mu_s(x_s) \}
  \]
Decomposition of entropy for trees

- by the junction tree theorem, any tree can be factorized in terms of its marginals $\mu \equiv \mu(\theta)$ as follows:

$$p(x; \theta) = \prod_{s \in V} \mu_s(x_s) \prod_{(s,t) \in E} \frac{\mu_{st}(x_s, x_t)}{\mu_s(x_s)\mu_t(x_t)}$$

- taking logs and expectations leads to an entropy decomposition

$$H(p(x; \theta)) = -A^*(\mu(\theta)) = \sum_{s \in V} H_s(\mu_s) - \sum_{(s,t) \in E} I_{st}(\mu_{st})$$

where

Single node entropy: $H_s(\mu_s) := -\sum_{x_s} \mu_s(x_s) \log \mu_s(x_s)$

Mutual information: $I_{st}(\mu_{st}) := \sum_{x_s, x_t} \mu_{st}(x_s, x_t) \log \frac{\mu_{st}(x_s, x_t)}{\mu_s(x_s)\mu_t(x_t)}$.

- thus, the dual function $A^*(\mu)$ has an explicit and easy form
Exact variational principle on trees

putting the pieces back together yields:

\[ A(\theta) = \max_{\mu \in \text{MARG}(T)} \{ \langle \theta, \mu \rangle + \sum_{s \in V} H_s(\mu_s) - \sum_{(s, t) \in E(T)} I_{st}(\mu_{st}) \}. \]

let’s try to solve this problem by a (partial) Lagrangian formulation

assign a Lagrange multiplier \( \lambda_{ts}(x_s) \) for each constraint

\[ C_{ts}(x_s) := \mu_s(x_s) - \sum_{x_t} \mu_{st}(x_s, x_t) = 0 \]

will enforce the normalization (\( \sum_{x_s} \mu_s(x_s) = 1 \)) and non-negativity constraints explicitly

the Lagrangian takes the form:

\[
\mathcal{L}(\mu; \lambda) = \langle \theta, \mu \rangle + \sum_{s \in V} H_s(\mu_s) - \sum_{(s, t) \in E(T)} I_{st}(\mu_{st}) \\
+ \sum_{(s, t) \in E} \left[ \sum_{x_t} \lambda_{st}(x_t) C_{st}(x_t) + \sum_{x_s} \lambda_{ts}(x_s) C_{ts}(x_s) \right]
\]
Lagrangian derivation (continued)

- taking derivatives of the Lagrangian w.r.t $\mu_s$ and $\mu_{st}$ yields

\[
\frac{\partial L}{\partial \mu_s(x_s)} = \theta_s(x_s) - \log \mu_s(x_s) + \sum_{t \in \mathcal{N}(s)} \lambda_{ts}(x_s) + C
\]

\[
\frac{\partial L}{\partial \mu_{st}(x_s, x_t)} = \theta_{st}(x_s, x_t) - \log \frac{\mu_{st}(x_s, x_t)}{\mu_s(x_s)\mu_t(x_t)} - \lambda_{ts}(x_s) - \lambda_{st}(x_t) + C'
\]

- setting these partial derivatives to zero and simplifying:

\[
\mu_s(x_s) \propto \exp \{\theta_s(x_s)\} \prod_{t \in \mathcal{N}(s)} \exp \{\lambda_{ts}(x_s)\}
\]

\[
\mu_s(x_s, x_t) \propto \exp \{\theta_s(x_s) + \theta_t(x_t) + \theta_{st}(x_s, x_t)\} \times \prod_{u \in \mathcal{N}(s) \setminus t} \exp \{\lambda_{us}(x_s)\} \prod_{v \in \mathcal{N}(t) \setminus s} \exp \{\lambda_{vt}(x_t)\}
\]

- enforcing the constraint $C_{ts}(x_s) = 0$ on these representations yields the familiar update rule for the messages $M_{ts}(x_s) = \exp(\lambda_{ts}(x_s))$:

\[
M_{ts}(x_s) \leftarrow \sum_{x_t} \exp \{\theta_t(x_t) + \theta_{st}(x_s, x_t)\} \prod_{u \in \mathcal{N}(t) \setminus s} M_{ut}(x_t)
\]
**C: Max-product algorithm on trees**

**Question:** What should be the form of a variational principle for computing modes?

**Intuition:** Consider behavior of the family \( \{ p(x; \beta \theta) \mid \beta > 0 \} \).

**Conclusion:** Problem of computing modes should be related to limiting form \( (\beta \to +\infty) \) of computing marginals.
Limiting form of the variational principle

- consider the variational principle for a discrete MRF of the form \( p(x; \beta \theta) \):
  \[
  \frac{1}{\beta} A(\beta \theta) = \frac{1}{\beta} \max_{\mu \in \text{MARG}} \{ \langle \beta \theta, \mu \rangle - A^*(\mu) \}.
  \]

- taking limits as \( \beta \to +\infty \) yields:
  \[
  \max_{x \in \mathcal{X}^N} \left\{ \sum_{s \in V} \theta_s(x_s) + \sum_{(s,t) \in E} \theta_{st}(x_s, x_t) \right\} = \max_{\mu \in \text{MARG}(G)} \{ \langle \theta, \mu \rangle \}.
  \]

  - computation of modes
  - linear program

- thus, computing the mode in a discrete MRF is equivalent to a linear program over the marginal polytope
Max-product on tree-structured MRFs

- recall the max-product (belief revision) updates:

\[ M_{ts}(x_s) \leftarrow \max_{x_t} \exp \left\{ \theta_t(x_t) + \theta_{st}(x_s, x_t) \right\} \prod_{u \in \mathcal{N}(t) \setminus s} M_{ut}(x_t) \]

- for trees, the variational principle (linear program) takes the especially simple form

\[
\max_{\mu \in \text{MARG}(T)} \left\{ \sum_{s \in V} \theta_s(x_s) \mu_s(x_s) + \sum_{(s,t) \in E} \theta_{st}(x_s, x_t) \mu_{st}(x_s, x_t) \right\}
\]

- constraint set is the marginal polytope for trees

\[
\text{MARG}(T) \ := \ \{ \mu \geq 0 \mid \sum_{x_s} \mu_s(x_s) = 1, \sum_{x_t} \mu_{st}(x_s, x_t) = \mu_s(x_s) \}
\]

- a similar Lagrangian formulation shows that max-product is an iterative method for solving this linear program (details in Wainwright & Jordan, 2003)
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**A: Mean field theory**

**Difficulty:** (typically) no explicit form for $-A^*(\mu)$ (i.e., entropy as a function of mean parameters) $\implies$ exact variational principle is intractable.

**Idea:** Restrict $\mu$ to a *subset* of distributions for which $-A^*(\mu)$ has a tractable form.

**Examples:**

(a) For product distributions $p(x) = \prod_{s \in V} \mu_s(x_s)$, entropy decomposes as $-A^*(\mu) = \sum_{s \in V} H_s(x_s)$.

(b) Similarly, for trees (more generally, decomposable graphs), the junction tree theorem yields an explicit form for $-A^*(\mu)$.

**Definition:** A subgraph $H$ of $G$ is *tractable* if the entropy has an explicit form for any distribution that respects $H$. 
Geometry of mean field

- let $H$ represent a tractable subgraph (i.e., for which $A^*$ has explicit form)

- let $\mathcal{M}_{tr}(G; H)$ represent tractable mean parameters:
  $$\mathcal{M}_{tr}(G; H) := \{\mu | \mu = \mathbb{E}_\theta[\phi(x)] \text{ s.t. } \theta \text{ respects } H\}.$$  

- under mild conditions, $\mathcal{M}_{tr}$ is a non-convex inner approximation to $\mathcal{M}$

- optimizing over $\mathcal{M}_{tr}$ (as opposed to $\mathcal{M}$) yields lower bound:
  $$A(\theta) \geq \sup_{\tilde{\mu} \in \mathcal{M}_{tr}} \{\langle \theta, \tilde{\mu} \rangle - A^*(\tilde{\mu})\}.$$
Alternative view: Minimizing KL divergence

- recall the *mixed form* of the KL divergence between $p(x; \theta)$ and $p(x; \tilde{\theta})$:

  \[
  D(\tilde{\mu} \| \theta) = A(\theta) + A^*(\tilde{\mu}) - \langle \tilde{\mu}, \theta \rangle
  \]

- try to find the “best” approximation to $p(x; \theta)$ in the sense of KL divergence

- in analytical terms, the problem of interest is

  \[
  \inf_{\tilde{\mu} \in \mathcal{M}_{tr}} D(\tilde{\mu} \| \theta) = A(\theta) + \inf_{\tilde{\mu} \in \mathcal{M}_{tr}} \left\{ A^*(\tilde{\mu}) - \langle \tilde{\mu}, \theta \rangle \right\}
  \]

- hence, finding the tightest lower bound on $A(\theta)$ is equivalent to finding the best approximation to $p(x; \theta)$ from distributions with $\tilde{\mu} \in \mathcal{M}_{tr}$
Example: Naive mean field algorithm for Ising model

- consider completely disconnected subgraph $H = (V, \emptyset)$
- permissible exponential parameters belong to subspace
  $\mathcal{E}(H) = \{ \theta \in \mathbb{R}^d \mid \theta_{st} = 0 \ \forall \ (s, t) \in E \}$
- allowed distributions take product form $p(x; \theta) = \prod_{s \in V} p(x_s; \theta_s)$, and generate
  $\mathcal{M}_{tr}(G; H) = \{ \mu \mid \mu_{st} = \mu_s \mu_t, \ \mu_s \in [0, 1] \}$.
- approximate variational principle:
  $\max_{\mu_s \in [0, 1]} \left\{ \sum_{s \in V} \theta_{s} \mu_s + \sum_{(s, t) \in E} \theta_{st} \mu_s \mu_t - \left[ \sum_{s \in V} \mu_s \log \mu_s + (1-\mu_s) \log (1-\mu_s) \right] \right\}$.

- Co-ordinate ascent: with all $\{\mu_t, t \neq s\}$ fixed, problem is strictly concave in $\mu_s$ and optimum is attained at
  $\mu_s \leftarrow \left\{ 1 + \exp\left[ -\left( \theta_s + \sum_{t \in \mathcal{N}(s)} \theta_{st} \mu_t \right) \right] \right\}^{-1}$
Example: Structured mean field for coupled HMM

- entropy of distribution that respects $H$ decouples into sum: one term for each chain.
- *structured mean field updates* are an iterative method for finding the tightest approximation (either in terms of KL or lower bound)
B: Belief propagation on arbitrary graphs

Two main ingredients:

1. Exact entropy $-A^*(\mu)$ is intractable, so let’s approximate it.

   The Bethe approximation $A_{\text{Bethe}}^*(\mu) \approx A^*(\mu)$ is based on the exact expression for trees:

   $$-A_{\text{Bethe}}^*(\mu) = \sum_{s \in V} H_s(\mu_s) - \sum_{(s, t) \in E} I_{st}(\mu_{st}).$$

2. The marginal polytope $\text{MARG}(G)$ is also difficult to characterize, so let’s use the following (tree-based) outer bound:

   $$\text{LOCAL}(G) := \{ \tau \geq 0 \mid \sum_{x_s} \tau_s(x_s) = 1, \sum_{x_t} \tau_{st}(x_s, x_t) = \tau_s(x_s) \},$$

**Note:** Use $\tau$ to distinguish these locally consistent *pseudomarginals* from globally consistent marginals.
Geometry of belief propagation

• combining these ingredients leads to the *Bethe variational principle:*  
  \[
  \max_{\tau \in \text{LOCAL}(G)} \left\{ \langle \theta, \tau \rangle + \sum_{s \in V} H_s(\mu_s) - \sum_{(s,t) \in E} I_{st}(\tau_{st}) \right\}
  \]

• belief propagation can be derived as an iterative method for solving a Lagrangian formulation of the BVP (Yedidia et al., 2002)

• belief propagation uses a *polyhedral outer approximation* to \( \mathcal{M} \)

• for any graph, \( \text{LOCAL}(G) \supseteq \text{MARG}(G) \).

• equality holds \( \iff G \) is a tree.
Illustration: Globally inconsistent BP fixed points

Consider the following assignment of pseudomarginals $\tau_s, \tau_{st}$:

![Graph with nodes and matrices]

- can verify that $\tau \in \text{LOCAL}(G)$, and that $\tau$ is a fixed point of belief propagation (with all constant messages)
- however, $\tau$ is globally inconsistent

Note: More generally: for any $\tau$ in the interior of $\text{LOCAL}(G)$, can construct a distribution with $\tau$ as a BP fixed point.
High-level perspective

- message-passing algorithms (e.g., mean field, belief propagation) are solving approximate versions of exact variational principle in exponential families

- there are two distinct components to approximations:
  (a) can use either inner or outer bounds to $\mathcal{M}$
  (b) various approximations to entropy function $-A^* (\mu)$

- mean field: non-convex inner bound and exact form of entropy

- BP: polyhedral outer bound and non-convex Bethe approximation

- Kikuchi and variants: tighter polyhedral outer bounds and better entropy approximations (Yedidia et al., 2001; Minka, 2001; Pakzad & Anantharam, 2002; Yildirim & McEliece, 2002)
Generalized belief propagation on hypergraphs

- a hypergraph is a natural generalization of a graph
- it consists of a set of vertices $V$ and a set $E$ of hyperedges, where each hyperedge is a subset of $V$
- convenient graphical representation in terms of poset diagrams

(a) Ordinary graph  (b) Hypertree (width 2)  (c) Hypergraph

- descendants and ancestors of a hyperedge $h$:

\[
\mathcal{D}^+(h) := \{ g \in E \mid g \subseteq h \}, \quad \mathcal{A}^+(h) := \{ g \in E \mid g \supseteq h \}.
\]
**Hypertree factorization and entropy**

- hypertrees are an alternative way to describe junction trees

- associated with any poset is a Möbius function \( \omega : E \times E \rightarrow \mathbb{Z} \)

\[
\omega(g, g) = 1, \quad \omega(g, h) = - \sum_{\{f | g \subseteq f \subset h\}} \omega(g, f)
\]

**Example:** For Boolean poset, \( \omega(g, h) = (-1)^{|h|\setminus|g|} \).

- use the Möbius function to define a correspondence between the collection of marginals \( \mu := \{\mu_h\} \) and new set of functions \( \varphi := \{\varphi_h\} \):

\[
\log \varphi_h(x_h) = \sum_{g \in \mathcal{D}^+(h)} \omega(g, h) \log \mu_g(x_g), \quad \log \mu_h(x_h) = \sum_{g \in \mathcal{D}^+(h)} \log \varphi_g(x_g).
\]

- any hypertree-structured distribution is guaranteed to factor as:

\[
p(x) = \prod_{h \in E} \varphi_h(x_h).
\]
Examples: Hypertree factorization

1. Ordinary tree:

\[ \varphi_s(x_s) = \mu_s(x_s) \quad \text{for any vertex } s \]
\[ \varphi_{st}(x_s, x_t) = \frac{\mu_{st}(x_s, x_t)}{\mu_s(x_s) \mu_t(x_t)} \quad \text{for any edge } (s, t) \]

2. Hypertree:

\[ \varphi_{1245} = \frac{\mu_{1245}}{\mu_5 \mu_5 \mu_5} \]
\[ \varphi_{45} = \frac{\mu_{45}}{\mu_5} \]
\[ \varphi_{5} = \mu_5 \]

Combining the pieces:

\[ p = \frac{\mu_{1245}}{\mu_5 \mu_5 \mu_5} \cdot \frac{\mu_{2356}}{\mu_5 \mu_5 \mu_5} \cdot \frac{\mu_{4578}}{\mu_5 \mu_5 \mu_5} \cdot \frac{\mu_{25}}{\mu_5} \cdot \frac{\mu_{45}}{\mu_5} \cdot \frac{\mu_{56}}{\mu_5} \cdot \frac{\mu_{58}}{\mu_5} \cdot \mu_5 = \frac{\mu_{1245} \mu_{2356} \mu_{4578}}{\mu_5 \mu_5} \]
Building augmented hypergraphs

Better entropy approximations via augmented hypergraphs.

(a) Original
(b) Clustering
(c) Full covering

(d) Kikuchi
(e) Fails single counting
C. Convex relaxations

Possible concerns with the Bethe/Kikuchi problems and variations?

(a) lack of convexity ⇒ multiple local optima, and substantial algorithmic complications

(b) failure to bound the log partition function

Goal: Techniques for approximate computation of marginals and parameter estimation based on:

(a) convex variational problems ⇒ unique global optimum

(b) relaxations of exact problem ⇒ upper bounds on $A(\theta)$

Usefulness of bounds:

(a) interval estimates for marginals

(b) approximate parameter estimation

(c) large deviations (prob. of rare events)
Bounds from “convexified” Bethe/Kikuchi problems

**Idea:** Upper bound $-A^*(\mu)$ by convex combination of tree-structured entropies.

$$-A^*(\mu) \leq -\rho(T^1)A^*(\mu(T^1)) - \rho(T^2)A^*(\mu(T^2)) - \rho(T^3)A^*(\mu(T^3))$$

- given any spanning tree $T$, define the moment-matched tree distribution:

$$p(x; \mu(T)) := \prod_{s \in V} \mu_s(x_s) \prod_{(s,t) \in E} \frac{\mu_{st}(x_s, x_t)}{\mu_s(x_s) \mu_t(x_t)}$$

- use $-A^*(\mu(T))$ to denote the associated tree entropy
- let $\rho = \{\rho(T)\}$ be a probability distribution over spanning trees
**Edge appearance probabilities**

**Experiment:** What is the probability \( \rho_e \) that a given edge \( e \in E \) belongs to a tree \( T \) drawn randomly under \( \rho \)?

![Diagram](image)

(a) Original  
(b) \( \rho(T^1) = \frac{1}{3} \)  
(c) \( \rho(T^2) = \frac{1}{3} \)  
(d) \( \rho(T^3) = \frac{1}{3} \)

In this example: \( \rho_b = 1; \quad \rho_e = \frac{2}{3}; \quad \rho_f = \frac{1}{3} \).

The vector \( \rho_e = \{ \rho_e \mid e \in E \} \) must belong to the *spanning tree polytope*, denoted \( \mathbb{T}(G) \).

(Edmonds, 1971)
Optimal bounds by tree-rewighted message-passing

Recall the constraint set of locally consistent marginal distributions:

\[
\text{LOCAL}(G) = \{ \tau \geq 0 | \sum_{x_s} \tau_s(x_s) = 1, \sum_{x_s} \tau_{st}(x_s, x_t) = \tau_t(x_t) \}.
\]

\[
\begin{align*}
&\text{normalization} \\
&\text{marginalization}
\end{align*}
\]

**Theorem:** (Wainwright, Jaakkola, & Willsky, 2002; To appear in IEEE-IT)

(a) For any given edge weights \( \rho_e = \{\rho_e\} \) in the spanning tree polytope, the optimal upper bound over all tree parameters is given by:

\[
A(\theta) \leq \max_{\tau \in \text{LOCAL}(G)} \{ \langle \theta, \tau \rangle + \sum_{s \in V} H_s(\tau_s) - \sum_{(s,t) \in E} \rho_{st} I_{st}(\tau_{st}) \}.
\]

(b) This optimization problem is strictly convex, and its unique optimum is specified by the fixed point of \( \rho_e \)-rewighted message passing:

\[
M^*_{ts}(x_s) = \kappa \sum_{x'_t \in x_t} \left\{ \exp \left[ \frac{\theta_{st}(x_s, x'_t)}{\rho_{st}} + \theta_t(x'_t) \right] \prod_{v \in \Gamma(t) \setminus s} \left[ M^*_{vt}(x_t) \right]^{\rho_{vt}} \frac{1}{[M^*_{st}(x_t)]^{(1-\rho_{ts})}} \right\}.
\]
Semidefinite constraints in convex relaxations

**Fact:** Belief propagation and its hypergraph-based generalizations all involve polyhedral (i.e., linear) outer bounds on the marginal polytope.

**Idea:** Use semidefinite constraints to generate more global outer bounds.

**Example:** For the Ising model, relevant mean parameters are $\mu_s = p(X_s = 1)$ and $\mu_{st} = p(X_s = 1, X_t = 1)$.

Define $y = [1 \ x]^T$, and consider the second-order moment matrix:

$$
\mathbb{E}[yy^T] = \begin{bmatrix}
1 & \mu_1 & \mu_2 & \ldots & \mu_n \\
\mu_1 & \mu_1 & \mu_{12} & \ldots & \mu_{1n} \\
\mu_2 & \mu_{12} & \mu_2 & \ldots & \mu_{2n} \\
\vdots & \vdots & \vdots & \ddots & \vdots \\
\mu_n & \mu_{n1} & \mu_{n2} & \ldots & \mu_n
\end{bmatrix}
$$

It must be positive semidefinite, which imposes (an infinite number of) linear constraints on $\mu_s, \mu_{st}$. 
Illustrative example

Locally consistent (pseudo)marginals

Second-order moment matrix

\[
\begin{bmatrix}
\mu_1 & \mu_{12} & \mu_{13} \\
\mu_{21} & \mu_2 & \mu_{23} \\
\mu_{31} & \mu_{32} & \mu_3
\end{bmatrix}
= 
\begin{bmatrix}
0.5 & 0.4 & 0.1 \\
0.4 & 0.5 & 0.4 \\
0.1 & 0.4 & 0.5
\end{bmatrix}
\]

Not positive-semidefinite!
Log-determinant relaxation

- based on optimizing over covariance matrices $M_1(\mu) \in \text{SDEF}_1(K_n)$

**Theorem:** Consider an outer bound $\text{OUT}(K_n)$ that satisfies:

$$\text{MARG}(K_n) \subseteq \text{OUT}(K_n) \subseteq \text{SDEF}_1(K_n)$$

For any such outer bound, $A(\theta)$ is upper bounded by:

$$\max_{\mu \in \text{OUT}(K_n)} \left\{ \langle \theta, \mu \rangle + \frac{1}{2} \log \det \left[ M_1(\mu) + \frac{1}{3} \text{blkdiag}[0, I_n] \right] \right\} + \frac{n}{2} \log\left( \frac{\pi e}{2} \right)$$

**Remarks:**

1. Log-det. problem can be solved efficiently by interior point methods.

2. Relevance for applications:
   (a) Upper bound on $A(\theta)$.
   (b) Method for computing approximate marginals.

(Wainwright & Jordan, 2003)
Results for approximating marginals

(a) Nearest-neighbor grid

(b) Fully connected

- average $\ell_1$ error in approximate marginals over 100 trials
- coupling types: repulsive (−), mixed (+/−), attractive (+)
Tree-based linear programming relaxations

- recall the equivalence between computing modes and solving an LP over the marginal polytope:

\[
\max_{x \in \mathcal{X}^N} \left\{ \sum_{s \in V} \theta_s(x_s) + \sum_{(s,t) \in E} \theta_{st}(x_s, x_t) \right\} = \max_{\mu \in \text{MARG}(G)} \{ \langle \theta, \mu \rangle \}.
\]

- our development suggests a very natural tree-based LP relaxation:

\[
\max_{x \in \mathcal{X}^N} \left\{ \sum_{s \in V} \theta_s(x_s) + \sum_{(s,t) \in E} \theta_{st}(x_s, x_t) \right\} \leq \max_{\mu \in \text{LOCAL}(G)} \{ \langle \theta, \mu \rangle \}.
\]

computation of modes

\text{tree-relaxed LP}

- this relaxation is always exact for trees

- its behavior for MRFs with cycles depends on the graph topology and strength of compatibility functions and observations
Geometry of LP relaxation

- two vertex types in relaxed polytope:
  
  integral: optimal configurations
  
  (e.g., \([0 \ 1 \ 0 \ 1 \ 0 \ 1]\))

  fractional: locally consistent
  
  (e.g., \([1 \ \frac{1}{2} \ 1 \ \frac{1}{2} \ 1 \ \frac{1}{2}]\))

- cost parameters \(\theta_s\) and \(\theta_{st}\) specify a particular direction in space

- if LP optimum for cost \(\theta\) attained at an integral vertex \(\implies\) LP relaxation is tight

- challenge: characterizing “good” directions in space
Link to message-passing algorithms

- tree-reweighted max-product algorithm linked to this tree-relaxed LP (analogous to Bethe and sum-product):
  - messages can be defined in terms of Lagrange multipliers
  - fixed point condition related to specification of dual optimum

- tree-reweighted form of max-product: \(M^*_{ts}(x_s) = \kappa \max_{x'_t \in \mathcal{X}_t} \left\{ \exp \left[ \frac{\theta_{st}(s, x'_t)}{\rho_{st}} + \theta_t(x'_t) \right] \frac{\prod_{v \in \Gamma(t) \setminus s} \left[ M^*_{vt}(x_t) \right]^{\rho_{vt}}}{\left[ M^*_{st}(x_t) \right]^{(1 - \rho_{ts})}} \right\} \).

- tree-relaxed LP and tree-reweighted max-product used in several applications
  - data association in sensor networks (Chen et al., 2003)
  - error-control decoding in communication (Feldman et al., 2005)
  - disparity computation in computer vision (Weiss et al., 2005; Kolmogorov, 2005)
Summary and future directions

• variational methods are based on converting statistical and computational problems to optimization:
  (a) complementary to sampling-based methods (e.g., MCMC)
  (b) a variety of new “relaxations” remain to be explored

• many open questions:
  (a) prior error bounds available only in special cases
  (b) extension to non-parametric settings?
  (c) hybrid techniques (variational and MCMC)
  (d) variational methods in parameter estimation
  (e) fast techniques for solving large-scale relaxations (e.g., SDPs, other convex programs)