

# Quantum computing Set 1 Solutions

## Problem 1

①

$$E = (+1) P_0 (\text{measuring } |0\rangle) + (-1) P_1 (\text{measuring } |1\rangle)$$

$$= |\langle 0|\psi\rangle|^2 - |\langle 1|\psi\rangle|^2$$

$$= \langle \psi|0\rangle\langle 0|\psi\rangle - \langle \psi|1\rangle\langle 1|\psi\rangle$$

$$= \langle \psi| (|0\rangle\langle 0| - |1\rangle\langle 1|) |\psi\rangle$$

$$= \langle \psi| Z |\psi\rangle.$$

② Since  $X = |+\rangle\langle +| - |-\rangle\langle -|$ , by analogy, measure in  $|+\rangle, |-\rangle$  basis and associate  $|+\rangle$  with  $+1$  and  $|-\rangle$  with  $-1$ .

## Problem 2

In the first interpretation of CHSH, on input  $x, y \in \{0, 1\}$  our score was  $E(u_x \oplus v_y \oplus xy)$

$u_x, v_y \in \{0, 1\}$ . To switch interpretation, we can scale our score to  $E(2(u_x \oplus v_y + xy) - 1)$

changing the game. Notice  $2(a) - 1 = -(-1)^a$  for  $a \in \{0, 1\}$ .

So it's equivalent to our score being

$$E\left(-(-1)^{u_x \oplus v_y \oplus xy}\right).$$

Let's redefine  $u_x, v_y \in \{\pm 1\}$  by  $u_x = (-1)^{u_x}$ ,  $v_y = (-1)^{v_y}$ .

Expanding  $E$  over  $\Pr(x, y) = \frac{1}{4}$  and scaling by 4 our score is equivalent to

$$\sum_{x, y} E(-(-1)^{xy} u_x v_y) = E(u_0 v_0) + E(u_1 v_0) + E(u_0 v_1) - E(u_1 v_1).$$

## Problem 2 (cont.)

Our transformation of the score was  $4(2(\cdot) - 1) = 8(\cdot) - 4$ .

So the quantum score on the original CHSH is now

$$\begin{aligned} 8 \cos^2 \pi/8 - 4 &= 8 \cos^2 \pi/8 - 4 \cos^2 \pi/8 - 4 \sin^2 \pi/8 \\ &= 4(\cos^2 \pi/8 - \sin^2 \pi/8) = 2\sqrt{2}. \end{aligned}$$

In any det. classical strategy, either  $v_0 + v_1 = 0$  or  $v_0 - v_1 = 0$ . Therefore

$$u_0(v_0 + v_1) + u_1(v_0 - v_1) \leq 2.$$

A randomized strategy is a linear combination of det. strategies.

### Problem 3

① Let  $a_x, b_x, c_x, a_y, b_y, c_y$  be the deterministic answers by Alice, Bob and Charlie for inputs  $x, y$ . Then in order to satisfy w.p. 1,

$$a_x b_x c_x = 1$$

$$a_y b_y c_x = -1$$

$$a_y b_x c_y = -1$$

$$a_x b_y c_y = -1$$

Multiplying all 4 equations together

$$\underbrace{a_x^2 a_y^2 b_x^2 b_y^2 c_x^2 c_y^2}_{\geq 0} = -1 \quad \leftarrow \text{impossible.}$$

For  $3/4$  strategy, everyone always answers  $-1$ .

### Problem 3 (cont.)

② wlog assume  $a_x, b_x, c_x, a_y, b_y, c_y$  are now fns  $\Omega \rightarrow \{\pm 1\}$  for some set  $\Omega$  and assume they share a random sample  $\sigma \in \Omega$ . To win with pr 1,  $\forall \sigma \in \Omega$  with  $p_\sigma(\sigma) > 0$ ,  $a_x(\sigma), b_x(\sigma), \dots, c_y(\sigma)$  must satisfy the equations on the previous page which is impossible. Therefore, randomness does not yield pr 1 success.

③ Since the eigenvalues of a tensor product are the products of the eigenvalues of the terms,

$\langle X \otimes X \otimes X \rangle_\gamma = \Pr(\text{win}) - \Pr(\text{lose}) = 2\Pr(\text{win}) - 1$  since we win if the product of the answers is 1. Suffices then to verify  $\langle X \otimes X \otimes X \rangle_\gamma = 1$ .

④ Analogously, check that  $\langle X \otimes Y \otimes Y \rangle_\gamma = -1$ .

⑤ Let  $\pi_\sigma : (a_1, a_2, a_3) \mapsto (a_{\sigma(1)}, a_{\sigma(2)}, a_{\sigma(3)})$ . Easy to check  $\pi_\sigma(|\gamma\rangle) = |\gamma\rangle$ .

Notice that  $Y \otimes X \otimes Y = \pi_{(12)} X \otimes Y \otimes Y \pi_{(12)}$ . Therefore,

$$\langle Y \otimes X \otimes Y \rangle_\gamma = \langle \gamma | \pi_{(12)} X \otimes Y \otimes Y \pi_{(12)} | \gamma \rangle = \langle \gamma | X \otimes Y \otimes Y | \gamma \rangle = -1.$$

### Problem 3 (cont.)

Therefore, wins with pr 1. Similar argument holds for  $Y \otimes Y \otimes X$ .

## Problem 4

(1) This actually holds for general  $d$ -dim maximally entangled state:  $\frac{1}{\sqrt{d}} \sum_i |i\rangle|i\rangle$ .

$$U \otimes \mathbb{I} |\Phi\rangle = \frac{1}{\sqrt{d}} \sum_i U|i\rangle \otimes |i\rangle$$

$$= \frac{1}{\sqrt{d}} \sum_{i,j} |j\rangle \underbrace{\langle j|U|i\rangle}_{\mathbb{I}} \otimes |i\rangle$$

$$= \frac{1}{\sqrt{d}} \sum_{i,j} |j\rangle \otimes |i\rangle \langle j|U|i\rangle$$

$$= \frac{1}{\sqrt{d}} \sum_{i,j} |j\rangle \otimes |i\rangle \underbrace{\langle i|U^\dagger|j\rangle}_{\mathbb{I}}$$

$$= \frac{1}{\sqrt{d}} \sum_j |j\rangle U^\dagger |j\rangle$$

$$= \mathbb{I} \otimes U^\dagger |\Phi\rangle.$$

↙ b.c.  $(\langle j|U|i\rangle)^\dagger = |i\rangle^\dagger U^\dagger \langle j|^\dagger$   
 $= \langle i|U^\dagger|j\rangle$

## Problem 4 (cont.)

Let  $\text{Ctrl}_{12}-U = \text{CNOT}_{12} \circ U$  and  $\text{Ctrl}_{13}-U^\top = \text{CNOT}_{13} \circ U^\top$ .  $\Rightarrow$

$$\begin{aligned}\text{Ctrl}_{12}-U |0\rangle|\Phi\rangle &= |0\rangle \mathbb{I}_2 |\Phi\rangle \\ &= |0\rangle|\Phi\rangle \\ &= \text{Ctrl}_{13}-U^\top |0\rangle|\Phi\rangle\end{aligned}$$

$$\begin{aligned}\text{Ctrl}_{12}-U |1\rangle|\Phi\rangle &= |1\rangle U_2 |\Phi\rangle \\ &= |1\rangle U_3^\top |\Phi\rangle \\ &= \text{Ctrl}_{13}-U^\top |1\rangle|\Phi\rangle\end{aligned}$$

Holds by linearity for general state on 1<sup>st</sup> qubit.



## Problem 4 (cont.)

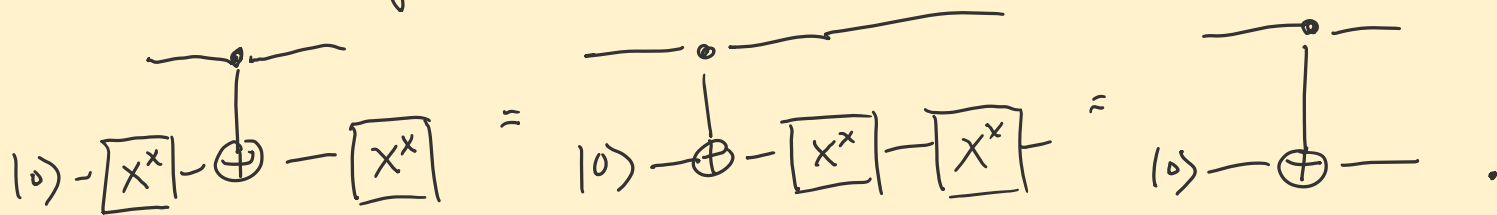
② We apply (part 1) to the EPR state to move the CNOT to the 3<sup>rd</sup> qubit.

Now we are measuring one qubit of the EPR state in the computational basis. The

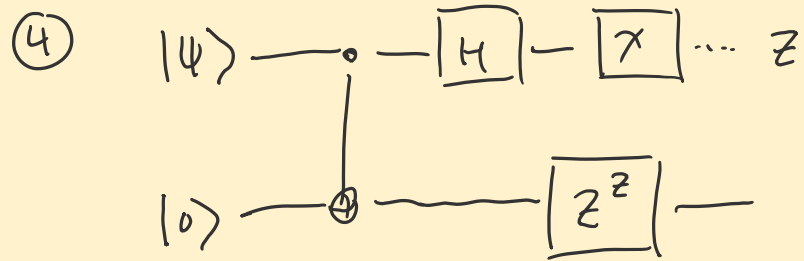
principle of deferred measurement cuts both ways so the circuit is equivalent to with

or  $\frac{1}{2}$  running the ckt in eq.(3) for each  $x \in \{0,1\}$ .

③  $|x\rangle = |0\rangle - \boxed{X^x}$ . As CNOT = Ctrl-X, its easy to see that commutation gives us



## Problem 4 (cont.)



is equivalent to  $|\psi\rangle|0\rangle =$

$$(\alpha|0\rangle + \beta|1\rangle) \otimes |0\rangle \rightarrow \alpha|00\rangle + \beta|11\rangle$$

$$\rightarrow \frac{1}{\sqrt{2}} (\alpha|00\rangle + \alpha|10\rangle + \beta|01\rangle - \beta|11\rangle)$$

$$= \frac{1}{\sqrt{2}} (|0\rangle|\psi\rangle + |1\rangle Z|\psi\rangle)$$

$$= \frac{1}{\sqrt{2}} \sum_z |z\rangle Z^z |\psi\rangle$$

meaning  $Z$  collapses the state to  $Z^z |\psi\rangle$  so applying  $Z^z$  produces  $|\psi\rangle$ .