

ESCHER

Expressive SCHeduling with Ephemeral Resources

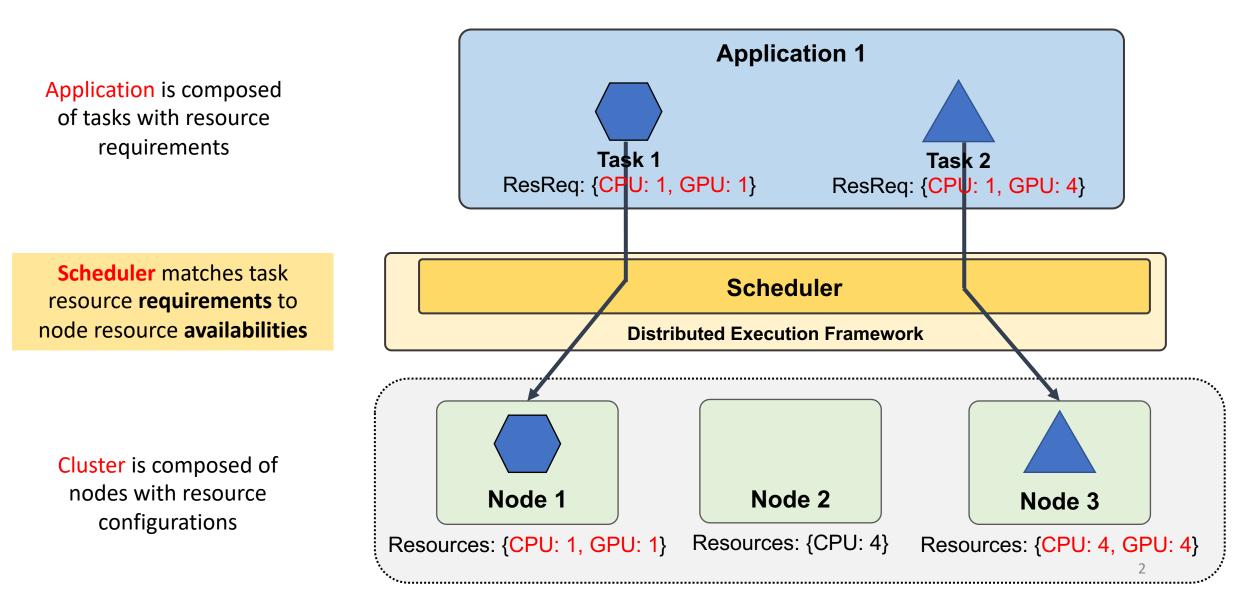
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Typical Distributed Application



Example - Distributed Training

Scheduling Requirements

1. Gang Scheduling

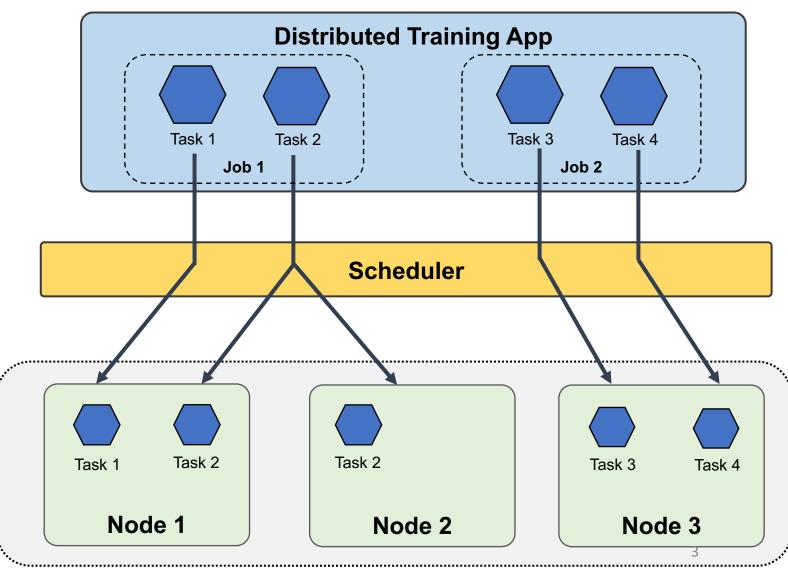
• Scheduling co-dependent tasks requires all-or-none semantics

2. Co-location

 Tasks of a job share parameter updates and must be placed on the same node for performance

3. Anti-affinity

 Avoid interference and resource contention by spreading jobs evenly spread across nodes



Example - Distributed Training

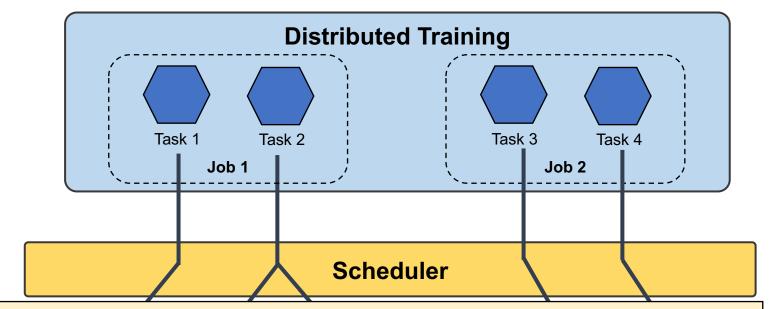
Scheduling Requirements

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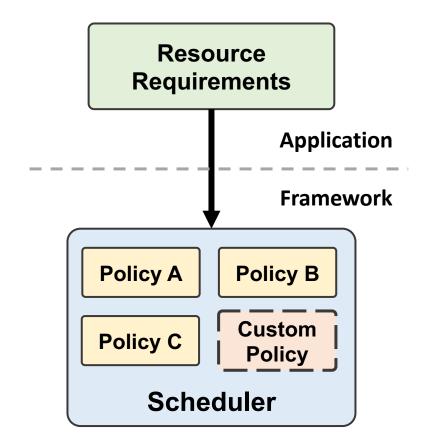
Supporting custom scheduling constraints requires *evolvable* schedulers

eveniy spreau across noues	TASK I TASK Z	IdSK Z	Task 3 Task 4
	Node 1	Node 2	Node 3

Evolvability in Monolithic Schedulers

Kubernetes, YARN

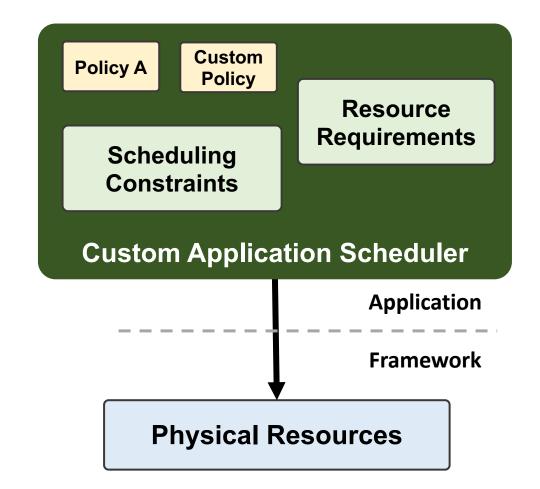
- Applications state resource requirements
- Scheduler provides a fixed set of supported policies
 - E.g., Affinity, anti-affinity
- Challenging to evolve
 - Implementing custom policies requires modifying the core scheduler
 - Can take months to add support
 - Difficult to maintain must commit to maintaining branch



Evolvability in Two-level Schedulers

Mesos, Omega

- Physical resources are exposed to applications
- Applications implement end-to-end scheduling
- Highly flexible, but application must implement a scheduler:
 - Resource state tracking
 - Task queueing
 - Fault tolerance



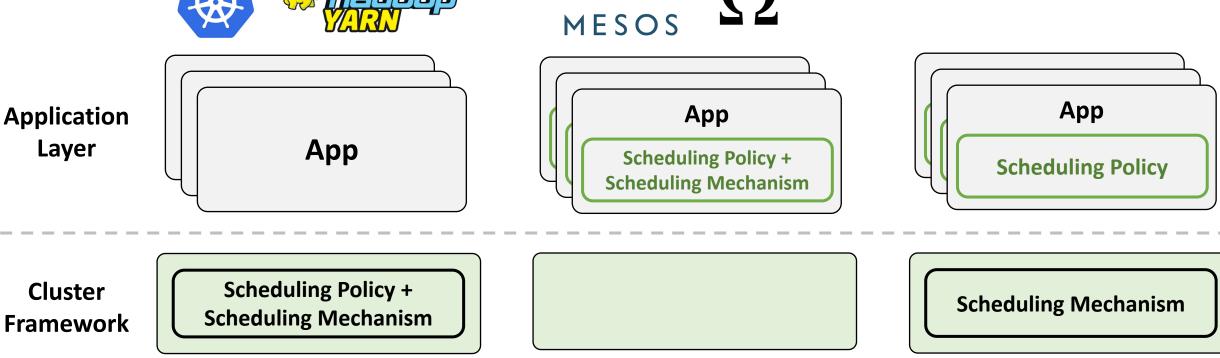
Summary of solutions today

Monolithic Schedulers

Simple, but hard to evolve



Application Layer



Two-level Schedulers

Highly evolvable, but complex



ESCHER

Simple and evolvable

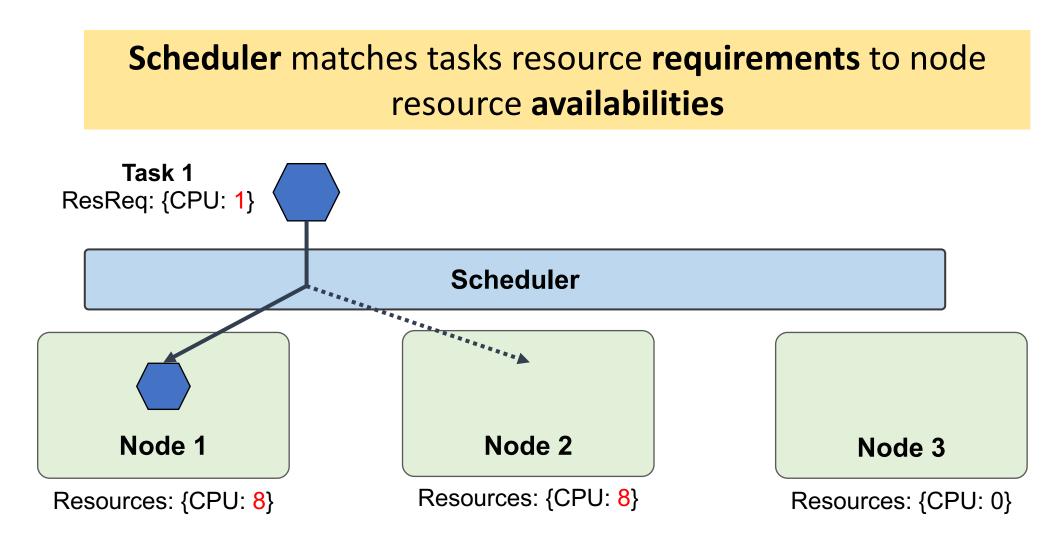
ESCHER Insights

With the following two scheduling abstractions, frameworks can allow applications to express a wide range of scheduling policies:

1. A simple resource matching scheduler

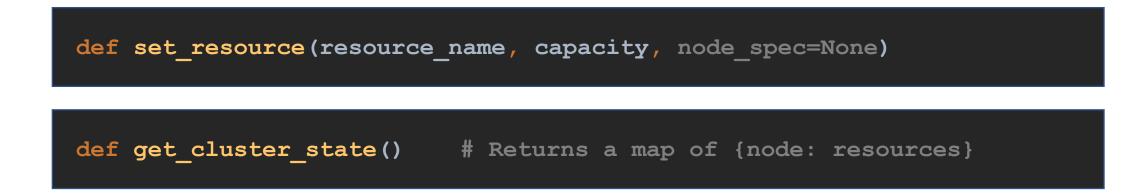
2. An API for applications to create resources at runtime

Abstraction 1 - Resource Matching Scheduler



Abstraction 2 – Create Resources on-the-fly

Applications should be able to create resources and get cluster state at runtime through an API

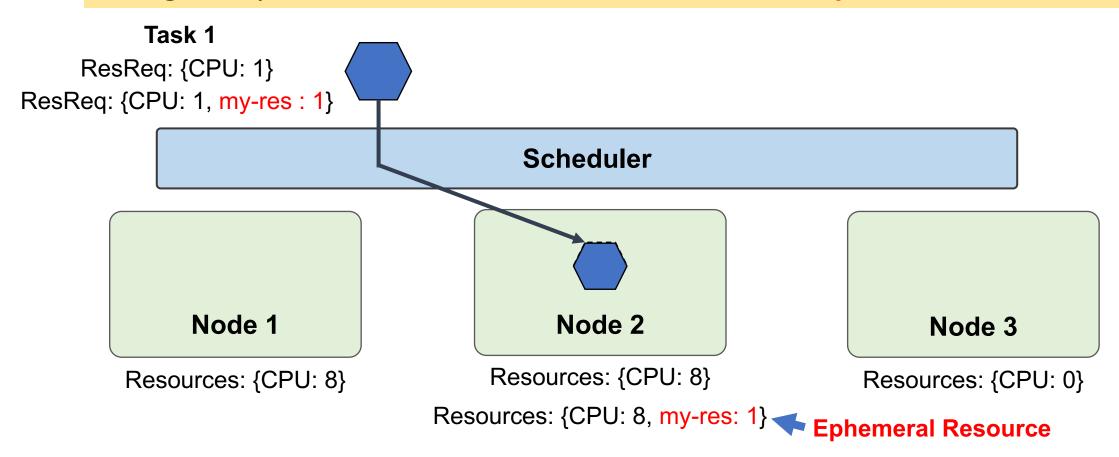


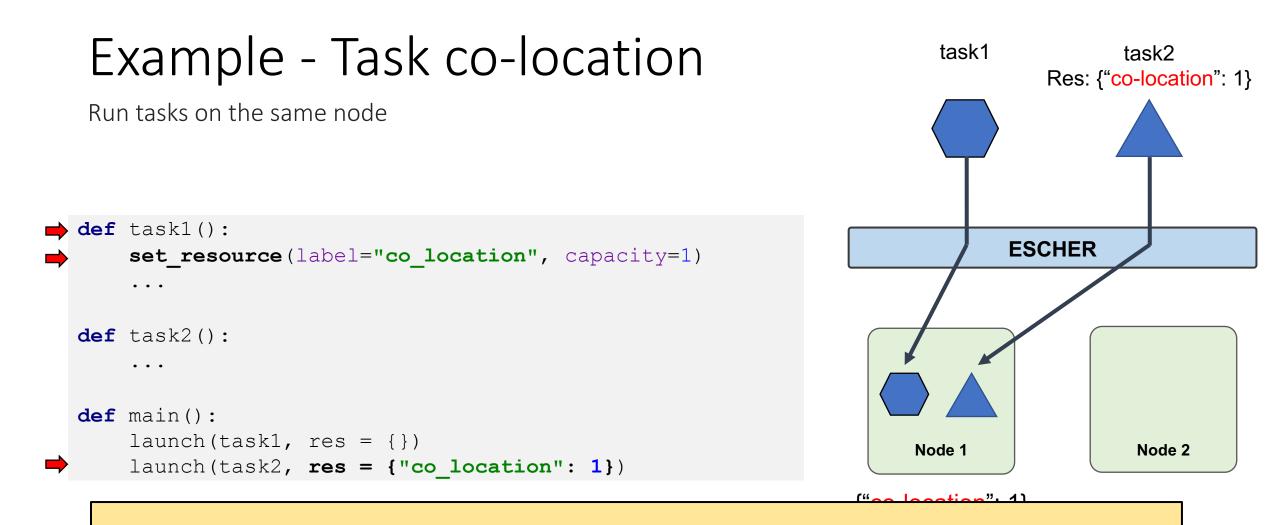
- Can specify resource availability constraints for resource creation
- If not node_spec not specified, resource created locally

Scheduling with Ephemeral Resources

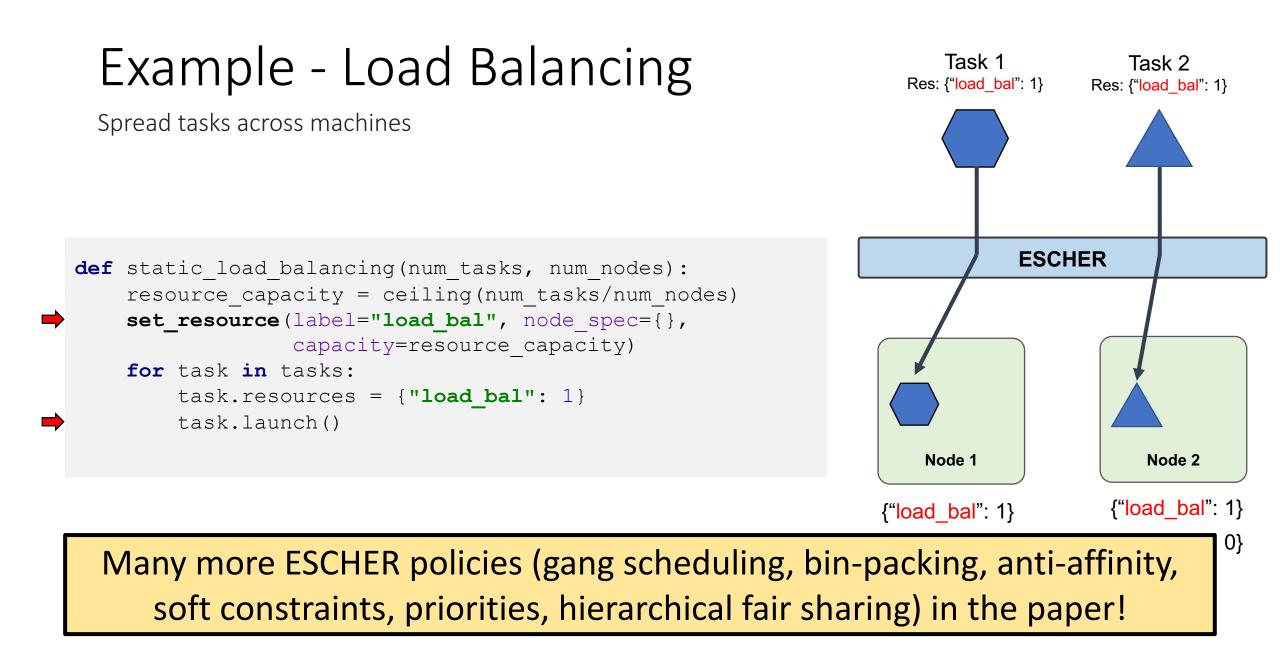
A simple resource matching scheduler can be induced to make targeted placement decisions with short-lived *ephemeral* resources

11



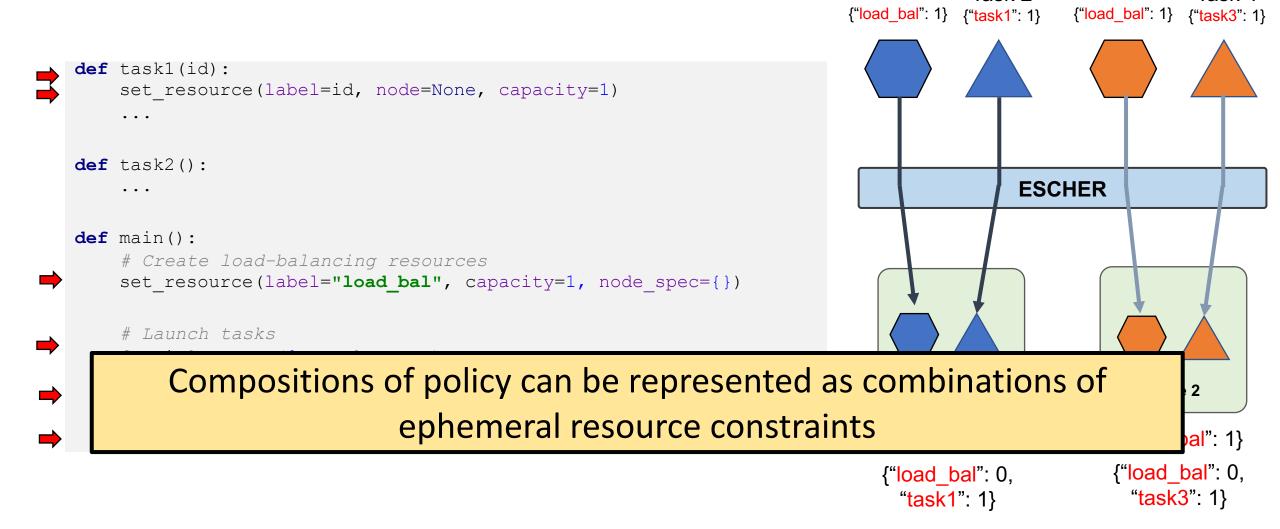


ESCHER allows declarative specification of scheduling policies by dynamically creating ephemeral resources



Policy Composition: Load Balancing & Co-location

Co-locate two tasks and spread out pairs of tasks



Task 1

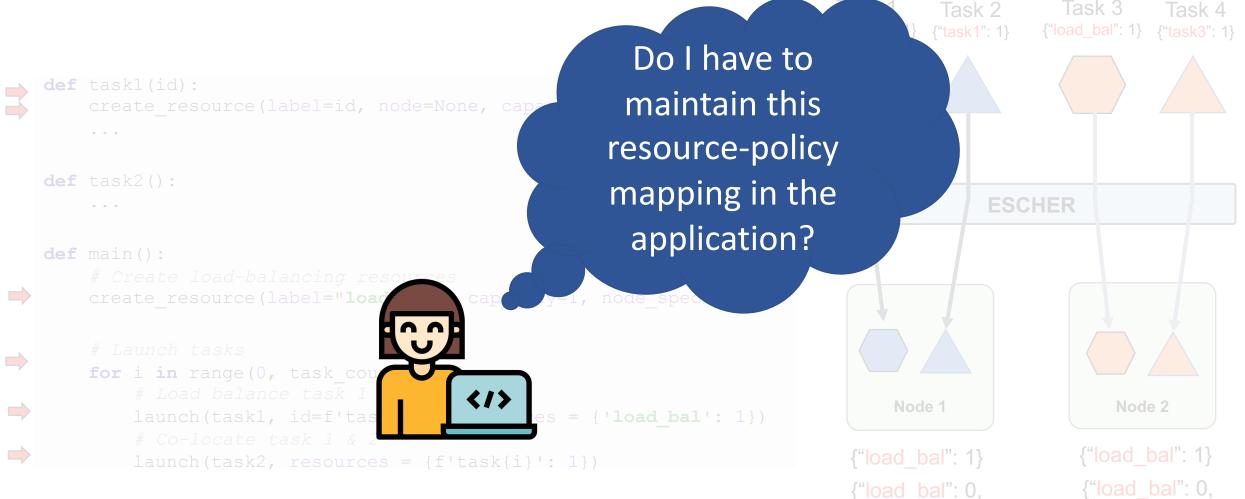
Task 2

Task 3

Task 4

Policy Composition: Load Balancing & Co-location

Co-locate two tasks and spread out pairs of tasks



"task3": 1}

"task1": 1}

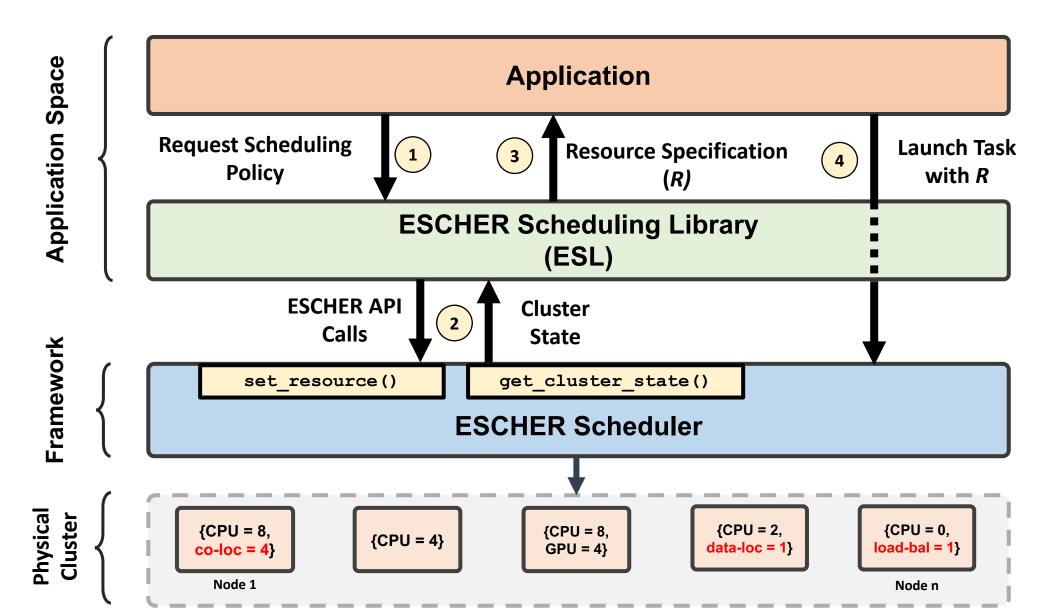
ESCHER Scheduling Libraries (ESLs)



- An app-level library of scheduling policies which encapsulate all state management for ephemeral resources
- Encourage code-reuse and simplify application code

```
def colocated_task():
    ...
def main():
    esl = CoLocationESL()
    coloc_res = esl.get_colocation_group("mygroup", res_req={gpu: 8})
    launch(colocated_tasks, res += coloc_res)
```

ESCHER Workflow



Implementation





kubernetes

Modified the Ray Scheduler to support online resource updates

No changes required in Kubernetes core – we reuse the extended resources API

Evaluation - AlphaZero



- AlphaZero trains an RL agent to play Go
- Training has two key processes:
 - **Board Generation**: CPU intensive generation of possible game states
 - **Board Evaluation**: A GPU agent predicts the "goodness" of the generated states and chooses an action
- These processes require both **co-location** and **load-balancing**

AlphaZero on Ray

0.10 **ESCHER** ESCHER is **1.5-2x faster** in Static Policy No Co-location 0.08 exploring board states than a locality-unaware scheduler Seconds Seconds 0.04 ESCHER performs comparably 0.02 with a static hard-coded policy with just **5 lines of code changes** 0.00 P50 P90 P99.9 P99 Latency Percentiles

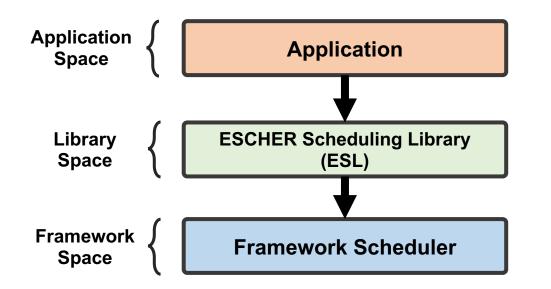
Board Exploration Latencies - 128 GPUs

AlphaZero on Ray

0.10 ESCHER ESCHER is **1.5-2x faster** in Static Policy No Co-location 0.08 exploring board states than a locality-unaware scheduler Sseconds 0.04 0.04 ESCHER performs comparably 0.02 Ν More results (MapReduce on 100 node K8s cluster, Hierarchical Fair W Sharing, Distributed Training, Microbenchmarks) in the paper!

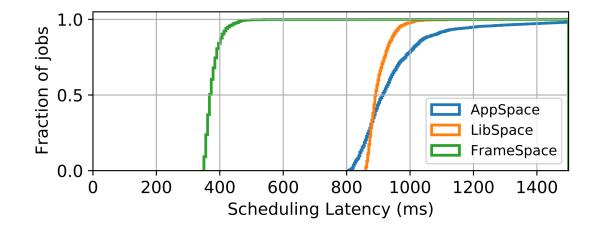
Board Exploration Latencies - 128 GPUs

ESCHER Overheads vs Evolvability



Gang Scheduling Implementation	Lines of Code	Median Scheduling Latency
AppSpace (ESCHER)		
LibSpace (ESLs in ESCHER)		
FrameSpace (Monolithic Scheduler)		

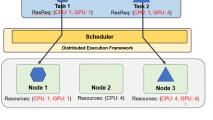
Using ESCHER adds latency for some policies such as gang scheduling, but significantly reduces the implementation burden.



23

ESCHER Summary

- Applications need fine-grained scheduling control without the complexity of implementing scheduling mechanisms.
- ESCHER presents an evolvable scheduler architecture with two key abstractions – a resource matching scheduler and set_resource API
- Ephemeral resources **are easily implemented in Ray and Kubernetes** and provide scheduling flexibility for a range of workloads with minimal overhead.



Application

