## CS61C
## Review of Cache/VM/TLB

**Lecture 27**

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## Outline

- Review Pipelining
- Review Interrupt/Polling Review slides
- Administrivia, "What's this Stuff Good for?"
- Conclusion



## Review 1/3: Cache/VM/TLB

- The Principle of Locality:
  - Program access a relatively small portion of the address space at any instant of time.
    - Temporal Locality: Locality in Time
    - Spatial Locality: Locality in Space
- 3 Major Categories of Cache Misses:
  - Compulsory Misses: sad facts of life. Example: cold start misses.
  - Capacity Misses: increase cache size
  - Conflict Misses: increase cache size and/or associativity.



## Review 2/3: Cache/VM/TLB

- Caches, TLBs, Virtual Memory all understood by examining how they deal with 4 questions:
  1) Where can block be placed?
  2) How is block found?
  3) What block is replaced on miss?
  4) How are writes handled?
- Page tables map virtual address to physical address
- TLBs are important for fast translation
- TLB misses are significant in processor performance



## Review 3/3: Cache/VM/TLB

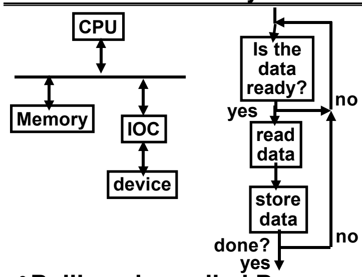
- Virtual memory was controversial at the time: can SW automatically manage 64KB across many programs?
  - 1000X DRAM growth removed controversy
- Today VM allows many processes to share single memory without having to swap all processes to disk;
  VM protection today is more important than memory hierarchy
- Today CPU time is a function of (ops, cache misses) vs. just f(ops): What does this mean to Compilers, Data structures, Algorithms?



## I/O Review Slide

- I/O gives computers their 5 senses
- I/O speed range is million to one
  - Mouse, keyboard, network, disk, display
- Processor speed means must synchronize with I/O devices before use

## Problem: How CPU Synch. with I/O device?




- Polling also called Programmed I/O
- Advantage: Simple - the processor is totally in control and does all the work



## Problems with Polling

- Polling overhead can consume a lot of CPU time when waiting for I/O device
  - busy wait loop not an efficient way to use the CPU unless the device is very fast!
- If not sure when need to do I/O, then lots of processor time spent when could be doing something else useful
- Solution: I/O Interrupt



## Why I/O Interrupt?

- Advantage: User program progress is only halted during actual transfer
- An I/O interrupt is like exception except:
  - An I/O interrupt is asynchronous
  - Further information needs to be conveyed
- An I/O interrupt is asynchronous with respect to instruction execution:
  - I/O interrupt is not associated with any instruction
  - I/O interrupt does not prevent any instruction from completion
    - CPU picks convenient point to take interrupt



## Example: Device Interrupt

External Interrupt

```
     ...
add  $r1,$r2,$r3
subi $r4,$r1,#4
slli $r4,$r4,#2
Hiccup(!)
lw   $r2,0($r4)
lw   $r3,4($r4)
add  $r2,$r2,$r3
sw   8($r4),$r2
     ...
```

PC saved

Restore PC

"Interrupt Handler"

```
Save registers
     ...
lw   $r1,20($r0)
lw   $r2,0($r1)
addi $r3,$r0,#5
sw   $r3,0($r1)
     ...
Restore registers
Clear current Int
```



## Review: Steps in Executing MIPS (Lec. 20)

1) Ifetch: Fetch Instruction, Increment PC
   - Page fault/Access fault on Instruction fetch?

2) Decode Instruction, Read Registers
   - Undefined Opcode?

3) Execute: Perform operation
   - Overflow?

4) Memory: read or write memory
   - Page fault/Access fault on Data access?

5) Write Back: Write Data to Register
   - I/O interrupts?



## Administrivia

- Everything but last 2 projects, last 2 homeworks on grade record is correct?
  - (We hope) graded last 2 homeworks, last 2 projects handed out in discussion this week; See Kelvin ASAP about disagreements
- Should have already filled out final survey to help future 61c; how many? haven't?
- Friday 61C Summary / Your Cal heritage / Cal v. Stanford CS education / HKN Evaluation
- Wed 5/12 Final 5-8PM in 1 Pimintel
  - Bring 2 sheets, both sides, #2 pencils
  - Sun 5/9 Final Review starting 2PM (1 Pimintel)

## What's it Good For: Sony Playstation 2000

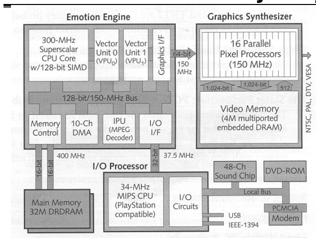



Figure 1. PlayStation 2000 employs an unprecedented level of parallelism to achieve workstation-class 3D performance.

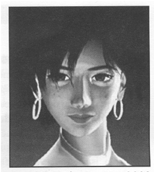


Figure 2. PlayStation 2000 screenshot. (Source: Namco)

- Emotion Engine: 6.2 GFLOPS, 75 million polygons per second (Microprocessor Report, 13:5)
  - Superscalar MIPS core + vector coprocessor
  - Claim: *Toy Story* realism brought to games!



## Problems with I/O Interrupts

- I/O interrupt is more complicated than exception:
  - Needs to convey the identity of the device generating the interrupt
- Special hardware is needed to:
  - Cause an interrupt (I/O device)
  - Detect an interrupt (processor)
  - Save the proper states to resume after the interrupt (processor)
- Where add special interrupt instructions, registers to instruction set?
- What prevents interrupt from occurring during interrupt handler?



## Review Coprocessor Registers

- Coprocessor 0 Registers:

| *name* | *number* | *usage* |
|---|---|---|
| BadVAddr | `$8` | Bad Virtual memory Address |
| Status | `$12` | Interrupt enable |
| Cause | `$13` | Exception type |
| EPC | `$14` | Instruction address |

  - Different registers from integer registers, just as Floating Point is another set of registers independent from integer registers



## Turn off interrupts? Interrupt Enable Bit

- Bit in Status Register determines whether or not interrupts enabled: Interrupt Enable bit (IE) (0 ⇒ off, 1 ⇒ on)
  - Also Kernel/User bit to support Virtual Memory modes

| (described later) | KU | IE |
|---|---|---|

Status Register



## Problems with Interrupt Enable

- Interrupt requests can have different urgencies
- Conventionally, from highest level to lowest level exception/interrupt levels:
  1) Bus error
  2) Illegal Instruction/Address trap
  3) High priority I/O Interrupt (fast response)
  4) Low priority I/O Interrupt (slow response)
- Alternative to blocking all interrupts?
  - Interrupt request needs to be prioritized



## Prioritizing Interrupts: Interrupt Mask

- Categorize interrupts and exceptions into levels, and allow selective interruption via Interrupt Mask(IM) in Status Register: 5 for HW interrupts
  - Interrupt only if IE==1 AND Mask bit == 1

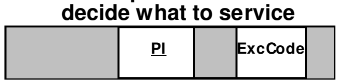
| | IM | | KU | IE |
|---|---|---|---|---|

Status Register

- How support interruption of lower priority interrupts?

## Interrupt levels

- Suppose there was an interrupt while the interrupt enable or mask bit is off: what should you do? (cannot ignore)
- Cause register has field--Pending Interrupts (PI)-- 5 bits wide (bits15:10) for each of the 5 HW interrupt levels
  - Bit becomes 1 when an interrupt at its level has occurred but not yet serviced
  - Interrupt routine checks pending interrupts ANDed with interrupt mask to decide what to service

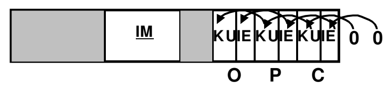



Cause Register



## Prioritizing Interrupts: Interrupt Mask

- To support interrupts of interrupts, have 3 deep stack in Status for IE,K/U bits:
  Current (1:0), Previous (3:2), Old (5:4)




Status Register

- How is MIPS software organized to take advantage of hardware priority scheme?



## Interrupt Levels in MIPS Software

- Conventionally, UNIX software system designed to have 4 to 6 Interrupt Priority Levels (IPL) that match the HW interrupt levels
- Processor always executing at one IPL, stored in a memory location and Status Register set accordingly
  - Processor at lowest IPL level, any interrupt accepted
  - Processor at highest IPL level, all interrupt ignored
  - Interrupt handlers and device drivers pick IPL to run at, faster response for some



## Handling Prioritized Interrupts

- OS convention to simplify software:
  - Process cannot be preempted by interrupt at same or lower level
  - Return to interrupted code as soon as no more interrupts at a higher level
  - Any piece of code is always run at same priority level
- How write interrupt routine so that it can be interrupted?



## Re-entrant Interrupt Routine?

- How allow interrupt of interrupts and safely save registers?
- Stack?
  - Resources consumed by each exception, so cannot tolerate arbitrary deep nesting of exceptions/interrupts
- With priority level system only interrupted by higher priority interrupt, so cannot be recursive
- ⇒ Only need one interrupt save area ("exception frame") per priority level



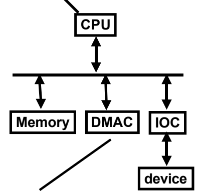
## Problems with CPU transferring data

- Typical I/O devices must transfer large amounts of data to memory of processor:
  - Disk must transfer complete block (4 KB? 16 KB?)
  - Large packets from network
  - Regions of frame buffer
- Can tie up processor depending on amount of I/O requests

## Delegating I/O Responsibility from CPU: DMA

CPU sends a starting address, direction, and length count to DMAC. Then issues "start".

° **Direct Memory Access (DMA):**
  • External to the CPU
  • Transfer blocks of data to or from memory without CPU intervention




DMA Controller (DMAC) provides signals for Peripheral Controller, and Memory Addresses and signals for Memory.



## Why DMA?

° DMA gives external device ability to write memory directly: much lower overhead than having processor request one word at a time



## Problems with DMA

° What if I/O devices write data that is currently in processor Cache?
  • The processor may never see new data!
  • Called “Cache coherence” problem

° Solutions:
  • Flush cache on every I/O operation (expensive)
  • Have hardware invalidate cache lines of potential address conflicts



## Problems with DMA

° Virtual Address or Physical Address?
  1) If virtual address, how do address translation, since memory uses physical addresses?
  2) If physical address, what happens if when cross a page boundary, as virtual memory may not be contiguous in physical memory?

° Solutions:
  1) Give DMA a small number of address translations, done by OS when start DMA
  2) Have a list of blocks, each no larger than a page, chained together



## Why use OS for I/O?

° The operating system acts as the interface between:
  • The I/O hardware and the program that requests I/O

° The Operating System must be able to prevent:
  • The user program from communicating with the I/O device directly

° If user programs could perform I/O directly:
  • Protection to the shared I/O resources could not be provided



## Responsibilities of the Operating System

° Three characteristics of the I/O systems:
  • The I/O system is shared by multiple program using the processor
  • I/O systems often use interrupts to communicate information about I/O operations.
    - Interrupts must be handled by the OS because they cause a transfer to supervisor mode
  • The low-level control of an I/O device is complex:
    - Managing a set of concurrent events
    - The requirements for correct device control are very detailed

## Operating System Requirements 1/2

- Provide protection to shared I/O resources
  - Guarantees that a user's program can only access the portions of an I/O device to which the user has rights
- Provides abstraction for accessing devices:
  - Supply routines that handle low-level device operation
- Handles the interrupts generated by I/O devices



## Operating System Requirements 2/2

- Provide equitable access to the shared I/O resources
  - All user programs must have equal access to the I/O resources
- Schedule accesses in order to enhance system throughput

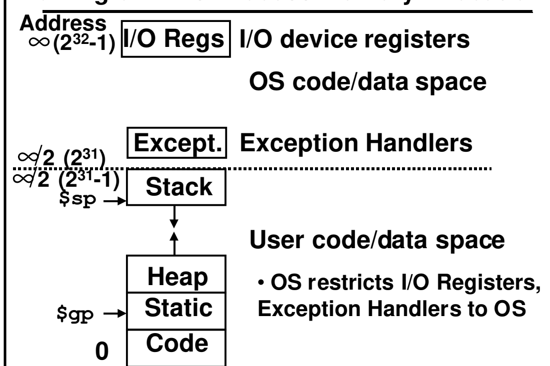


## How Protect I/O?

- MIPS memory maps I/O devices to allow load-store access to send commands, receive status and data
- To prevent user program from accessing data despite having a 32-bit virtual address, need protection
- (See above) MIPS CPU runs in 2 privilege levels: user mode and kernel mode
  - User mode: limited to bottom half of 32-bit virtual address
  - Kernel mode: can access full 32-bit virtual address; special areas to enable booting machine before TLB valid



## Drawing of MIPS Process Memory Allocation






## In More Depth: Actual MIPS address names

- Virtual address divided into 4 areas:

1) kuseg (low 2 GB) - for user mode, always translated via TLB and through cache

2) kseg0 (next 0.5 GB) - translated by striping off top 1 bit (kernel mode); maps to low 0.5GB of physical memory via caches

3) kseg1 (next 0.5 GB) - translated by striping off top 3 bits (kernel mode); maps to low 0.5GB of physical memory, not via caches

4) kseg2 (top 1 GB) - kernel mode, always translated via TLB and through cache



## How User safely invoke Operating System?

- 2 instructions
  - `break`: intended to implement break point debugging feature
  - `syscall`: intended to ask OS for specific services by passing argument in register

## Summary 1/2

- Wide range of devices
  - multimedia and high speed networking poise important challenges
- Delegating data transfer responsibility from the CPU: DMA
- I/O performance limited by weakest link in chain between OS and device
- Operating System started as shared I/O library



## Summary 2/2

- I/O device notifying the operating system:
  - Polling: it can waste a lot of processor time
  - I/O interrupt: similar to exception except it is asynchronous
- MIPS OS support / Interrupt control:
  - Interrupt Enable bit, stacked IE bits, Interrupt Priority Levels, Interrupt Mask
  - Support for OS abstraction: Kernel/User bit, stacked KU bits, `syscall`, `rfe`
  - MIPS follows coprocessor abstraction to add resources, instructions for OS
  - OS Re-entrant via restricting interrupt to higher priority