CS61C
Virtual Memory, 2nd try
+ Advice on Courses

Lecture 19

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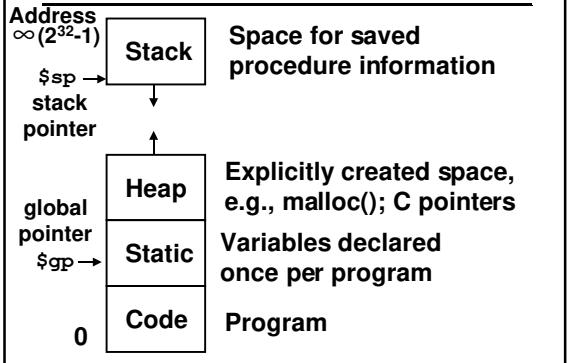
Review 1/1

- 2 Approaches to Writes
  - Write Through: Write twice, current level and next level down
  - Write Back: Write only to current level, write dirty block to next level down only on miss
- Reduce Miss Penalty?
  - Add a (L2) cache
- Manage memory-disk transfers?
  - Explain as cache: block = page, miss = page fault, write back, fully associative, LRU...
  - Included protection as bonus, now critical

Outline

- Sharing/Protecting Memory by Swapping
- Sharing /Protecting Memory by Base/Bounds
- Weakness of Base/Bounds Protection
- Administrivia, "Advice on courses"
- Virtual Memory/Paging
- Problems and solutions of Virtual Memory
- Conclusion

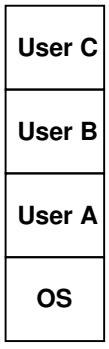
Review: C memory allocation



Memory Allocation in Reality

- Prior schemes assumed 1 process is running; reality is many processes
- Even a single person using:
  - Browser
  - Compiler
  - Mail
  - ...

(For now, ignore I/O split of address space)

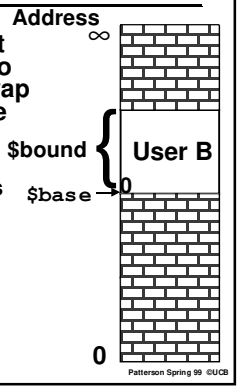


"Sharing" Solution #1 (earliest)

- Copy other Users (processes) to Disk when not in use
- When time to run a new process, swap new user and old user
- Protection: only 1 user at a time, so OK; (but OS not protected)
- Problem: slow to swap processes between memory and disk

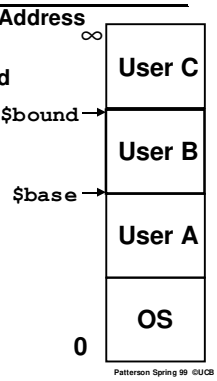
"Sharing" Solution #2: Base/Bounds

- Add Hardware to protect users from each other so that must not always swap to disk to share machine
- Add special registers:
  - Bound Register: address must fit inside this value
    - Also called "Protection Register"
  - Base Register: added to each memory address
    - Also called "Relocation Register"



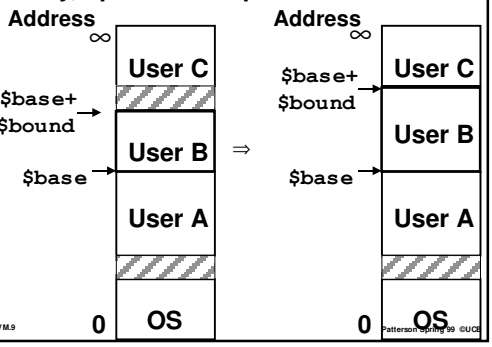
Base/Bounds mapping Function

- Real memory address:
  - if Program Address >= Bound ADDRESS ERROR
  - if Program Address < Bound Base + Program Address
- Note:
  - Allows safe sharing without always swapping to disk
  - $base, $bound are special registers, like $epc, NOT like $0, ..., $31



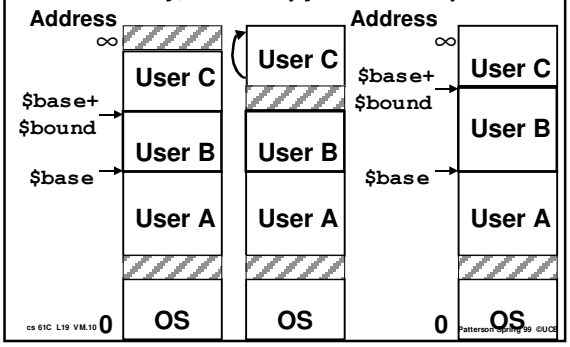
1: What if need to grow size of Process B?

- If lucky, space above process available



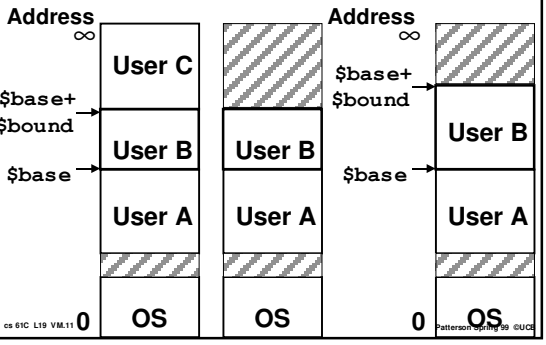
2: What if need to grow size of Process B?

- If unlucky, must copy to create space



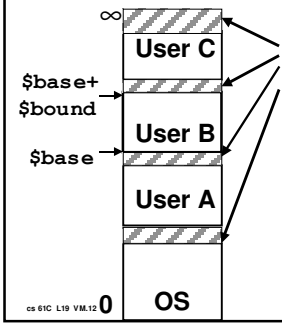
3: What if need to grow size of Process B?

- If unluckier, must swap to disk



Danger of Base/Bounds Sharing Scheme?

- Lots of free memory, but in small fragments ⇒ too small to be useful



Enough space for User D, but discontinuous so cannot use

- Called "Memory Fragmentation" Problem
  - Happens when processes quit too
- Base/bounds HW will not solve

Administrivia

- Project 5: Due 4/14: design and implement a cache (in software) and plug into instruction simulator
- Next Readings: 5.1 (skip clocking), 5.2, 4.5 (pages 230-240), 4.6 (pages 250-256, 264), 4.7 (pages 265-273)
  - How many lectures to cover: 2? 3?
- 9th homework: Due 4/14 or 4/16? 7PM
  - Exercises 7.35, 4.24
  - Vote on leaving on Wed vs. delaying to Friday (only when there is a project due on Wed)
  - Wed: TA office hours; Fri, spread 61C load

Administrivia: General Course Philosophy

- Take variety of undergrad courses now to get introduction to areas
  - Can learn advanced material on own later once know vocabulary
- Who knows what you will work on over a 40 year career?

Administrivia: Courses for Telebears

- General Philosophy
  - Take courses from great teachers!
  - HKN ratings; > 6 very good, < 5 not good
  - hkn.eecs/toplevel/coursesurveys.html
- Top Faculty / Course (may teach soon)

• CS 150 logic design	Katz	6.2	F92
• CS 152 computer	Patterson	6.7	S95
• CS 164 compilers	Rowe	6.1	S98
• CS 169 SW engin.	Brewer	6.2	S98
• CS 174 combinatorics	Sinclair	6.1	F97
• CS 186 data bases	Wang	6.2	S98

If many good teachers: My recommendations

- CS169 Software Engineering
  - Everyone writes programs, even HW designers
  - Often programs are written in groups ⇒ learn skill now in school (before it counts)
- EE122 Introduction to Communication Networks
  - World is getting connected; communications must play major role
- CS162 Operating Systems
  - All special-purpose HW will run a layer of SW that uses processes and concurrent programming; CS162 is the closest thing

If many good teachers: Courses to consider

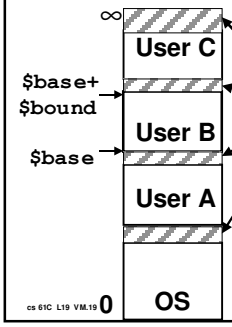
- E190 Technical Communication
  - Talent in writing and speaking critical for success
  - Now required for EECS majors
- CS 150 Lab Hardware Design
  - Hands on HW design
- CS 152 Design a Computer
- CS 188 Understand databases
  - Information more important now than computation?

Administrivia: Courses for Telebears

- Remember:
- Teacher quality more important to learning experience than official course content
- Take courses from great teachers!

Inspiration for Solution #3

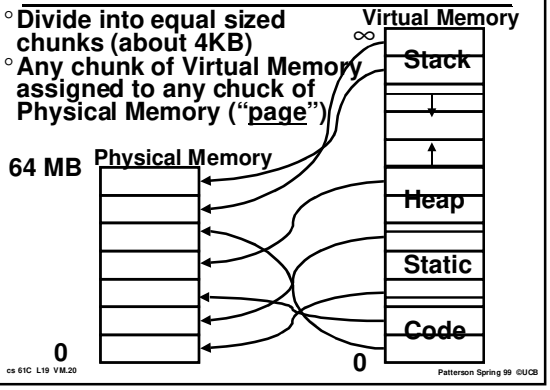
- Provide Hardware to allow discontinuous memory to look continuous to new process



- Enough space for User D, but discontinuous
- Solution #3 called Virtual Memory, or Paging
- Will specify new HW mapping function

Mapping Virtual Memory to Physical Memory

- Divide into equal sized chunks (about 4KB)
- Any chunk of Virtual Memory assigned to any chuck of Physical Memory ("page")



Virtual Memory Mapping Function

- Cannot have simple function to predict arbitrary mapping
- Use table lookup of mappings

Page Number	Offset

- Use table lookup ("Page Table") for mappings: Page number is index
- Virtual Memory Mapping Function
  - Physical Offset = Virtual Offset
  - Physical Page Number = PageTable[Virtual Page Number]

(P.P.N. also called "Page Frame")

Notes on Page Table

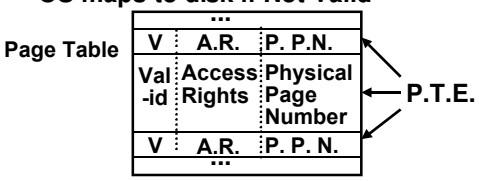
- Solves Fragmentation: all chunks same size, so all holes can be used
- To grow a process, ask Operating System
  - If unused pages, OS uses them first
  - If not, OS swaps some pages to disk
  - (Least Recently Used to pick pages to swap)
- OS must reserve "Swap Space" on disk for each process
- Each process has own Page Table
- Will add details, but Page Table is essence of Virtual Memory

Virtual Memory Problem #1

- Not enough physical memory!
  - Only, say, 64 MB of physical memory
  - N processes, each 4GB of virtual memory!
  - Could have 1K virtual pages/physical page!
- Spatial Locality to the rescue
  - Each page is 4 KB, lots of nearby references
  - No matter how big program is, at any time only accessing a few pages
  - "Working Set": recently used pages

Page Table Entry (PTE) Format

- Contains either Physical Page Number or indication not in Main Memory
- OS maps to disk if Not Valid



- If valid, also check if have permission to use page: Access Rights may be Read Only, Read/Write, Executable

Virtual Memory Problem #2

- Map every address ⇒ 1 extra memory accesses for every memory access
- Observation: since locality in pages of data, must be locality in virtual addresses of those pages
- Why not use a cache of virtual to physical address translations to make translation fast? (small is fast)
- For historical reasons, cache is called a Translation Lookaside Buffer, or TLB

Typical TLB Format

Virtual Address	Physical Address	Dirty	Ref	Valid	Access Rights

- TLB just a cache on the page table mappings
- TLB access time comparable to cache (much less than main memory access time)
- Ref: Used to help calculate LRU on replacement
- Dirty: since use write back, need to know whether or not to write page to disk when replaced

What if not in TLB?

- Option 1: Hardware checks page table and loads new Page Table Entry into TLB
- Option 2: Hardware traps to OS, up to OS to decide what to do
- MIPS follows Option 2: Hardware knows nothing about page table format

Virtual Memory Problem #3

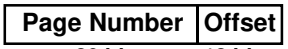
- Page Table too big!
  - 4GB Virtual Memory ÷ 4 KB page
    ⇒ ~ 1 million Page Table Entries
    ⇒ 4 MB just for Page Table
- Variety of solutions to tradeoff memory size of mapping function for slower when miss TLB
  - Make TLB large enough, highly associative so rarely miss on address translation
  - CS 162 will go over more options and in greater depth

Page Table Shrink #1:

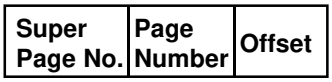
- Page Table Entry has Virtual Page Number and Physical Page Number (as in TLB entry)
- Keep only 1 Page Table Entry per Physical Page vs. per Virtual Page
- Find match on miss via hash function
- Size Reduction
  - e.g., 64 MB ÷ 4 KB ⇒ 16 K entries * 4 Bytes
  - 64 KB vs. 4096 KB for normal page table

Page Table Shrink #2:

- Single Page Table

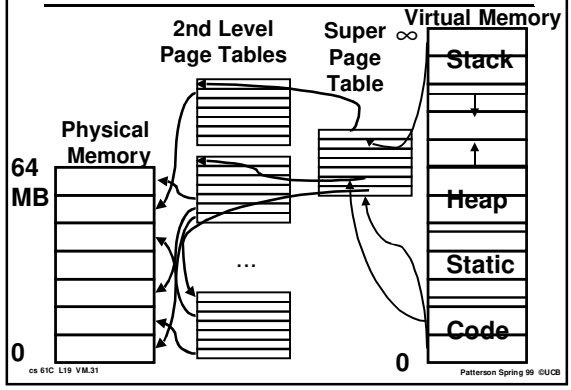


- Multilevel Page Table



- Only have second level page table for valid entries of super level page table

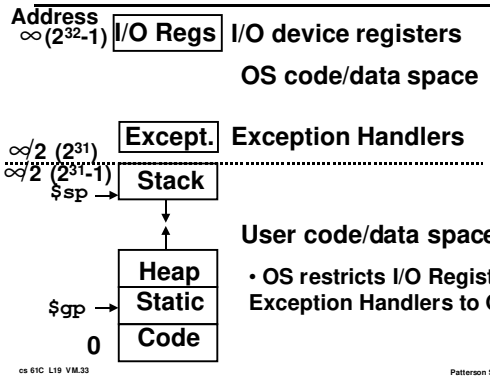
2-level Page Table



Space Savings for Multi-Level Page Table

- If only 10% of entries of Super Page Table have valid enties, then size is roughly 1/10-th of single level page table
  - Exercise 7.35 explores exact size

Note: Actual MIPS Process Memory Allocation



"And in Conclusion.." 1/1

- Virtual Memory allows protected sharing of memory between processes with less swapping to disk, less fragmentation than always swap or base/bound
- Spatial Locality means Working Set of Pages is all that must be in memory for process to run fairly well
- TLB to reduce performance cost of VM
- Need more compact representation to reduce memory size cost of simple 1-level page table (especially 32- ⇒ 64-bit address)
- Next: Introduction to processors design