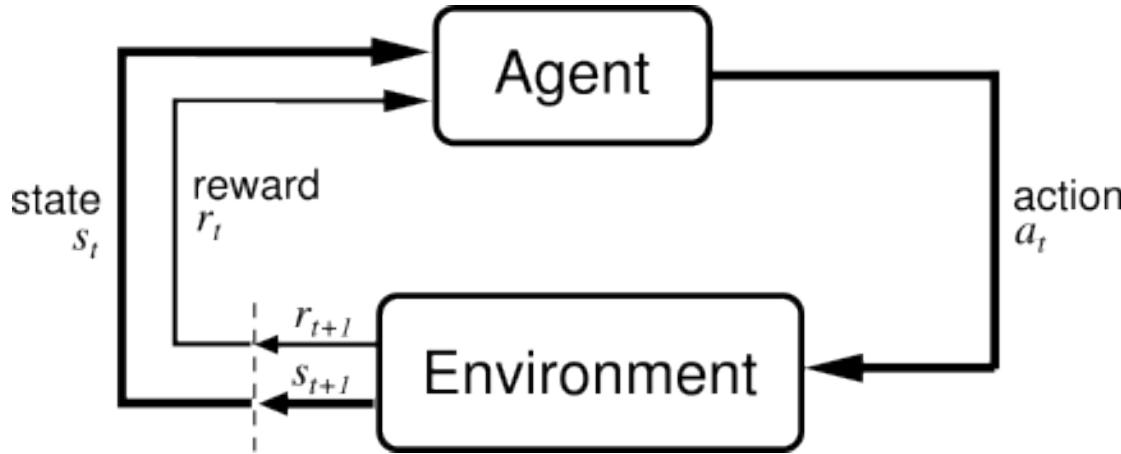


Solving Continuous MDPs with Discretization

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Markov Decision Process



Assumption: agent gets to observe the state

Markov Decision Process (S, A, T, R, γ, H)

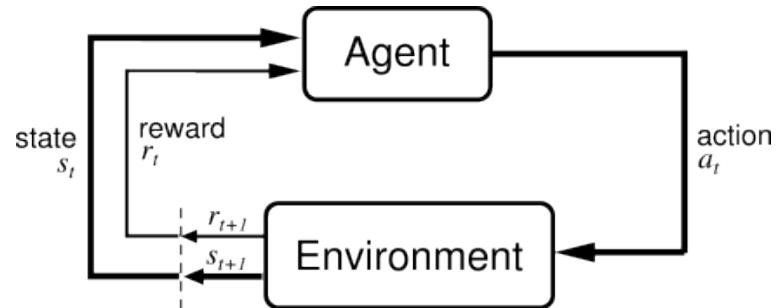
Given

- S : set of states
 - A : set of actions
 - $T: S \times A \times S \times \{0, 1, \dots, H\} \rightarrow [0, 1]$ $T_t(s, a, s') = P(s_{t+1} = s' \mid s_t = s, a_t = a)$
 - $R: S \times A \times S \times \{0, 1, \dots, H\} \rightarrow \mathbb{R}$ $R_t(s, a, s') = \text{reward for } (s_{t+1} = s', s_t = s, a_t = a)$
 - γ in $(0, 1]$: discount factor
- H: horizon over which the agent will act

Goal:

- Find $\pi^*: S \times \{0, 1, \dots, H\} \rightarrow A$ that maximizes expected sum of rewards, i.e.,

$$\pi^* = \arg \max_{\pi} E \left[\sum_{t=0}^H \gamma^t R_t(S_t, A_t, S_{t+1}) \mid \pi \right]$$



Value Iteration

Algorithm:

Start with $V_0^*(s) = 0$ for all s.

For $i = 1, \dots, H$

For all states s in S :

$$V_{i+1}^*(s) \leftarrow \max_a \sum_{s'} T(s, a, s') [R(s, a, s') + \gamma V_i^*(s')]$$

$$\pi_{i+1}^*(s) \leftarrow \arg \max_{a \in A} \sum_{s'} T(s, a, s') [R(s, a, s') + \gamma V_i^*(s')]$$

This is called a **value update** or **Bellman update/back-up**

$V_i^*(s)$ = expected sum of rewards accumulated starting from state s , acting optimally for i steps

$\pi_i^*(s)$ = optimal action when in state s and getting to act for i steps

Continuous State Spaces

- $S = \text{continuous set}$
- Value iteration becomes impractical as it requires to compute, for all states s in S :

$$V_{i+1}^*(s) \leftarrow \max_a \sum_{s'} T(s, a, s') [R(s, a, s') + V_i^*(s')]$$

Markov chain approximation to continuous state space dynamics model ("discretization")

- Original MDP
 (S, A, T, R, γ, H)



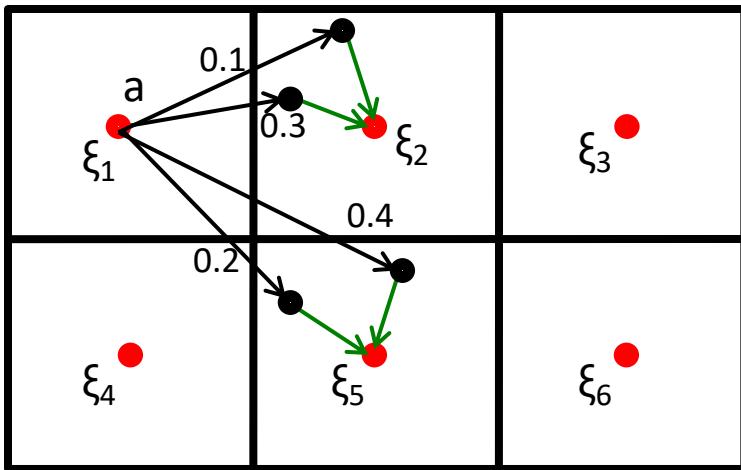
- Discretized MDP
 $(\bar{S}, \bar{A}, \bar{T}, \bar{R}, \gamma, H)$

- Grid the state-space: the vertices are the discrete states.
- Reduce the action space to a finite set.
 - Sometimes not needed:
 - When Bellman back-up can be computed exactly over the continuous action space
 - When we know only certain controls are part of the optimal policy (e.g., when we know the problem has a "bang-bang" optimal solution)
- Transition function: see next few slides.

Outline

- Discretization
- Lookahead policies
- Examples
- Guarantees
- Connection with function approximation

Discretization Approach 1: Snap onto nearest vertex



Discrete states: $\{ \xi_1, \dots, \xi_6 \}$

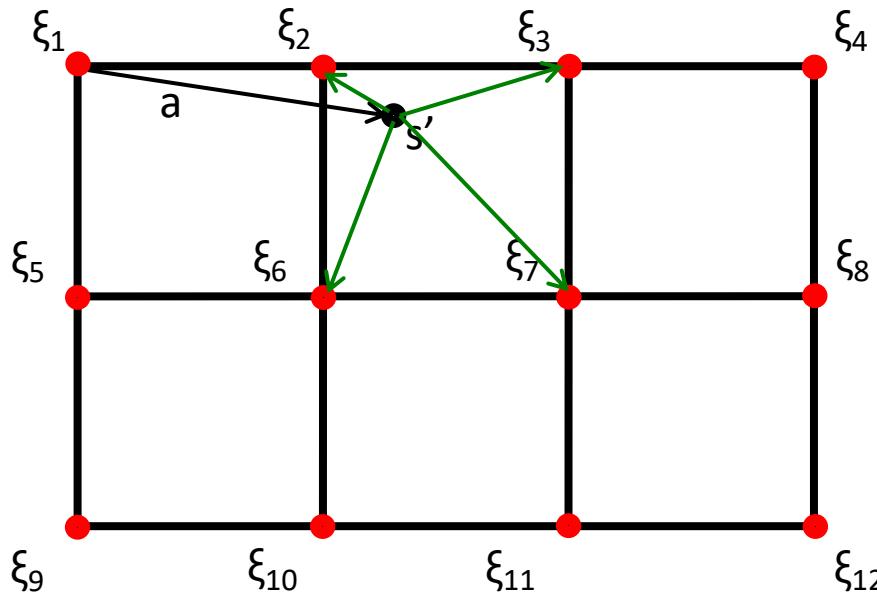
$$P(\xi_2|\xi_1, a) = 0.1 + 0.3 = 0.4;$$

$$P(\xi_5|\xi_1, a) = 0.4 + 0.2 = 0.6$$

Similarly define transition probabilities for all ξ_i

- Discrete MDP just over the states $\{\xi_1, \dots, \xi_6\}$, which we can solve with value iteration
- If a (state, action) pair can results in infinitely many (or very many) different next states: sample the next states from the next-state distribution

Discretization Approach 2: Stochastic Transition onto Neighboring Vertices



Discrete states: $\{\xi_1, \dots, \xi_{12}\}$

$$P(\xi_2 | \xi_1, a) = p_A$$

$$P(\xi_3 | \xi_1, a) = p_B$$

$$P(\xi_6 | \xi_1, a) = p_C$$

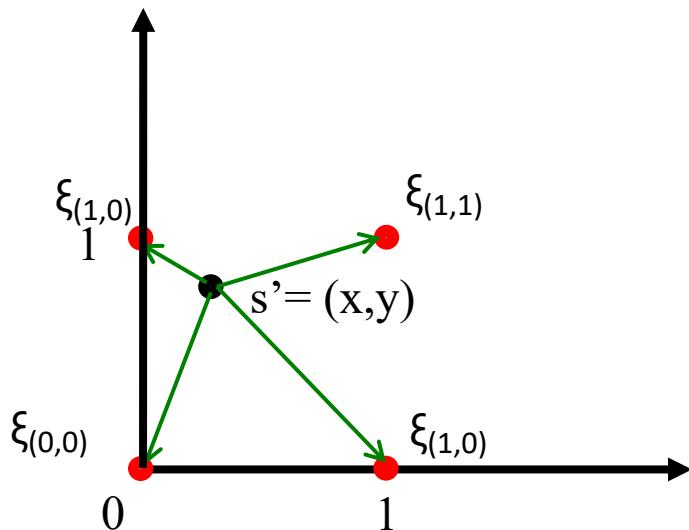
$$P(\xi_7 | \xi_1, a) = p_D$$

$$\text{s.t. } s' = p_A \xi_2 + p_B \xi_3 + p_C \xi_6 + p_D \xi_7$$

- If stochastic dynamics: Repeat procedure to account for all possible transitions and weight accordingly
- Many choices for p_A, p_B, p_C, p_D

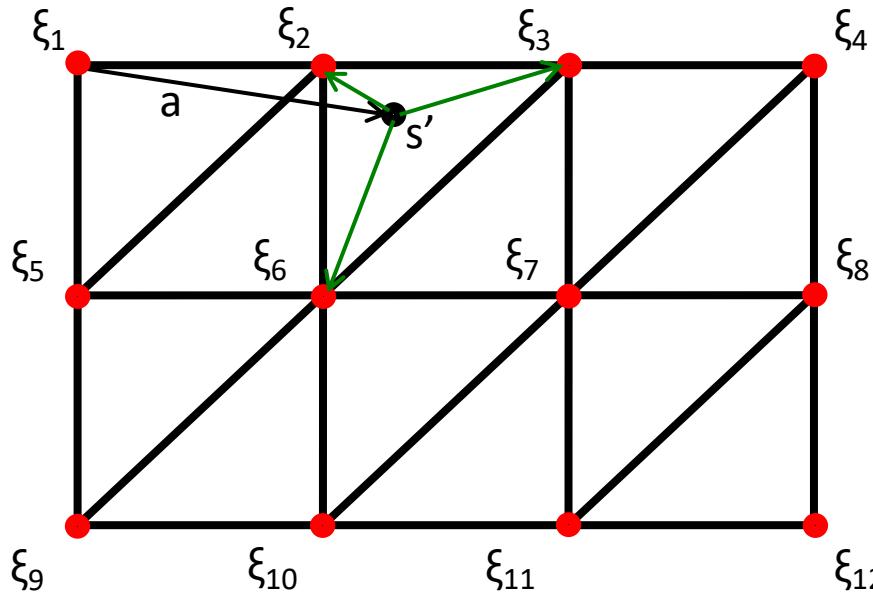
Discretization Approach 2: Stochastic Transition onto Neighboring Vertices

- One scheme to compute the weights: put in normalized coordinate system $[0,1] \times [0,1]$.



$$\begin{aligned} s' = (1 - x)(1 - y) & \quad \xi_{(0,0)} \\ +x(1 - y) & \quad \xi_{(0,1)} \\ +(1 - x)y & \quad \xi_{(1,0)} \\ +xy & \quad \xi_{(1,1)} \end{aligned}$$

Kuhn Triangulation**



Discrete states: $\{\xi_1, \dots, \xi_{12}\}$

$$P(\xi_2|\xi_1, a) = p_A;$$

$$P(\xi_3|\xi_1, a) = p_B;$$

$$P(\xi_6|\xi_1, a) = p_C;$$

$$\text{s.t. } s' = p_A\xi_2 + p_B\xi_3 + p_C\xi_6$$

Kuhn Triangulation**

- Allows efficient computation of the vertices participating in a point's barycentric coordinate system and of the convex interpolation weights (aka its barycentric coordinates)

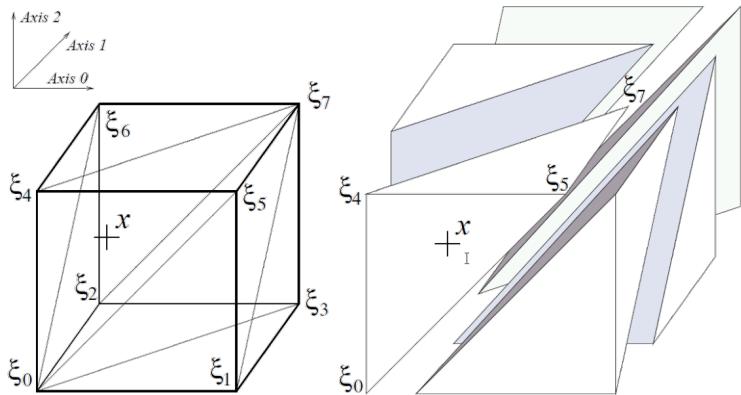


Figure 2. The Kuhn triangulation of a (3d) rectangle. The point x satisfying $1 \geq x_2 \geq x_0 \geq x_1 \geq 0$ is in the simplex $(\xi_0, \xi_4, \xi_5, \xi_7)$.

- See Munos and Moore, 2001 for further details.

Kuhn triangulation (from Munos and Moore)**

3.1. Computational issues

Although the number of simplexes inside a rectangle is factorial with the dimension d , the computation time for interpolating the value at any point inside a rectangle is only of order $(d \ln d)$, which corresponds to a sorting of the d relative coordinates (x_0, \dots, x_{d-1}) of the point inside the rectangle.

Assume we want to compute the indexes i_0, \dots, i_d of the $(d+1)$ vertices of the simplex containing a point defined by its relative coordinates (x_0, \dots, x_{d-1}) with respect to the rectangle in which it belongs to. Let $\{\xi_0, \dots, \xi_{2^d}\}$ be the corners of this d -rectangle. The indexes of the corners use the binary decomposition in dimension d , as illustrated in Figure 2. Computing these indexes is achieved by sorting the coordinates from the highest to the smallest: there exist indices j_0, \dots, j_{d-1} , permutation of $\{0, \dots, d-1\}$, such that $1 \geq x_{j_0} \geq x_{j_1} \geq \dots \geq x_{j_{d-1}} \geq 0$. Then the indices i_0, \dots, i_d of the $(d+1)$ vertices of the simplex containing the point are: $i_0 = 0$, $i_1 = i_0 + 2^{j_0}$, ..., $i_k = i_{k-1} + 2^{j_{k-1}}$, ..., $i_d = i_{d-1} + 2^{j_{d-1}} = 2^d - 1$. For example, if the coordinates satisfy: $1 \geq x_2 \geq x_0 \geq x_1 \geq 0$ (illustrated by the point x in Figure 2) then the vertices are: ξ_0 (every simplex contains this vertex, as well as $\xi_{2^d-1} = \xi_7$), ξ_4 (we added 2^2), ξ_5 (we added 2^0) and ξ_7 (we added 2^1).

Let us define the *barycentric coordinates* $\lambda_0, \dots, \lambda_d$ of the point x inside the simplex $\xi_{i_0}, \dots, \xi_{i_d}$ as the positive coefficients (uniquely) defined by: $\sum_{k=0}^d \lambda_k = 1$ and $\sum_{k=0}^d \lambda_k \xi_{i_k} = x$. Usually, these barycentric coordinates are expensive to compute; however, in the case of Kuhn triangulation these coefficients are simply: $\lambda_0 = 1 - x_{j_0}$, $\lambda_1 = x_{j_0} - x_{j_1}$, ..., $\lambda_k = x_{j_{k-1}} - x_{j_k}$, ..., $\lambda_d = x_{j_{d-1}} - 0 = x_{j_{d-1}}$. In the previous example, the barycentric coordinates are: $\lambda_0 = 1 - x_2$, $\lambda_1 = x_2 - x_0$, $\lambda_2 = x_0 - x_1$, $\lambda_3 = x_1$.

Discretization: Our Status

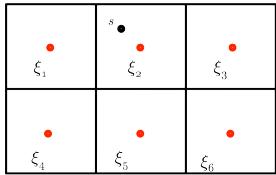
- Have seen two ways to turn a continuous state-space MDP into a discrete state-space MDP
- When we solve the discrete state-space MDP, we find:
 - Policy and value function for the discrete states
 - They are optimal for the discrete MDP, but typically not for the original MDP
- Remaining questions:
 - How to act when in a state that is not in the discrete states set?
 - How close to optimal are the obtained policy and value function?

How to Act (i): No Lookahead

- For state s not in discretization set choose action based on policy in nearby states

- Nearest Neighbor**

$$\pi(s) = \pi(\xi_i) \text{ for } \xi_i = \arg \min_{\xi \in \{\xi_1, \dots, \xi_N\}} \|s - \xi\|$$



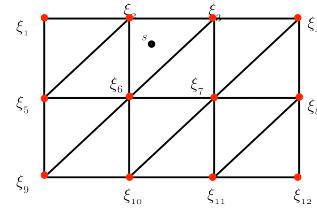
E.g., $\pi(s) = \pi(\xi_2)$

- Stochastic Interpolation:**

$$\text{Find } p_1, \dots, p_N \text{ s.t. } s = \sum_{i=1}^N p_i \xi_i$$

Choose $\pi(\xi_i)$ with probability p_i

For continuous actions, can also interpolate: $\sum_{i=1}^N p_i \pi(\xi_i)$



E.g., for $s = p_2\xi_2 + p_3\xi_3 + p_6\xi_6$, choose $\pi(\xi_2), \pi(\xi_3), \pi(\xi_6)$ with respective probabilities p_2, p_3, p_6

How to Act (ii): 1-step Lookahead

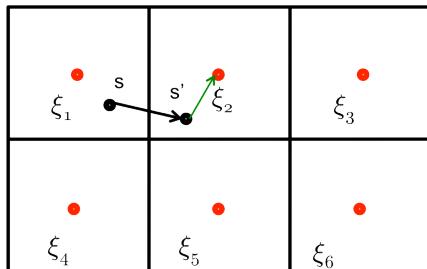
- Forward simulate for 1 step, calculate reward + value function at next state from discrete MDP

$$\max_{a_t} E \left[R(s_t, a_t) + \sum_i P(\xi_i; s_{t+1}) V(\xi_i) \right]$$

- if dynamics deterministic no expectation needed
- If dynamics stochastic, can approximate with samples

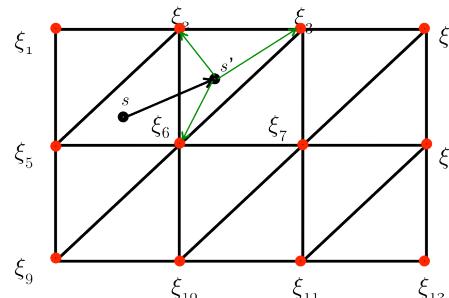
- Nearest Neighbor**

$$P(\xi_i; s') = \begin{cases} 1 & \text{if } \xi_i = \arg \min_{\xi \in \{\xi_1, \dots, \xi_N\}} \|s - \xi\| \\ 0 & \text{otherwise} \end{cases}$$



- Stochastic Interpolation**

$$P(\xi_i; s') \text{ such that } s' = \sum_{i=1}^N P(\xi_i; s') \xi_i$$



How to Act (iii): n-step Lookahead

$$\max_{a_t, a_{t+1}, \dots, a_{t+k-1}} E \left[R(s_t, a_t) + R(s_{t+1}, a_{t+1}) + \dots + R(s_{t+k-1}, a_{t+k-1}) + \sum_i P(\xi_i; s_{t+k}) V(\xi_i) \right]$$

■ What action space to maximize over, and how?

- Option 1: Enumerate sequences of discrete actions we ran value iteration with
- Option 2: Randomly sampled action sequences (“random shooting”)
- Option 3: Run optimization over the actions
 - Local gradient descent [see later lectures]
 - Cross-entropy method

Intermezzo: Cross-Entropy Method (CEM)

- CEM = black-box method for (approximately) solving:

$$\max_x f(x)$$

with $x \in \mathbb{R}^n$ and $f : \mathbb{R}^n \rightarrow \mathbb{R}$

Note: f need not be differentiable

Intermezzo: Cross-Entropy Method (CEM)

$$\max_x f(x)$$

CEM:

sample $\mu^{(0)} \sim \mathcal{N}(0, \sigma^2)$

for iter i = 1, 2, ...

 for e = 1, 2, ...

 sample $x^{(e)} \sim \mathcal{N}(\mu^{(i)}, \sigma^2)$

 compute $f(x^{(e)})$

 endfor

$\mu^{(i+1)} = \text{mean}\{x^{(e)} : f(x^{(e)}) \text{ in top } 10\%\}$

Intermezzo: Cross-Entropy Method (CEM)

CEM:

sample $\mu^{(0)} \sim \mathcal{N}(0, \sigma^2)$

for iter i = 1, 2, ...

 for e = 1, 2, ...

 sample $x^{(e)} \sim \mathcal{N}(\mu^{(i)}, \sigma^2)$

 compute $f(x^{(e)})$

 endfor

$\mu^{(i+1)} = \text{mean}\{x^{(e)} : f(x^{(e)}) \text{ in top 10\%}\}$

- sigma and 10% are hyperparameters
- can in principle also fit sigma to top 10% (or full covariance matrix if low-D)
- How about discrete action spaces?
 - Within top 10%, look at frequency of each discrete action in each time step, and use that as probability
 - Then sample from this distribution

Note: there are many variations, including a max-ent variation, which does a weighted mean based on $\exp(f(x))$

Outline

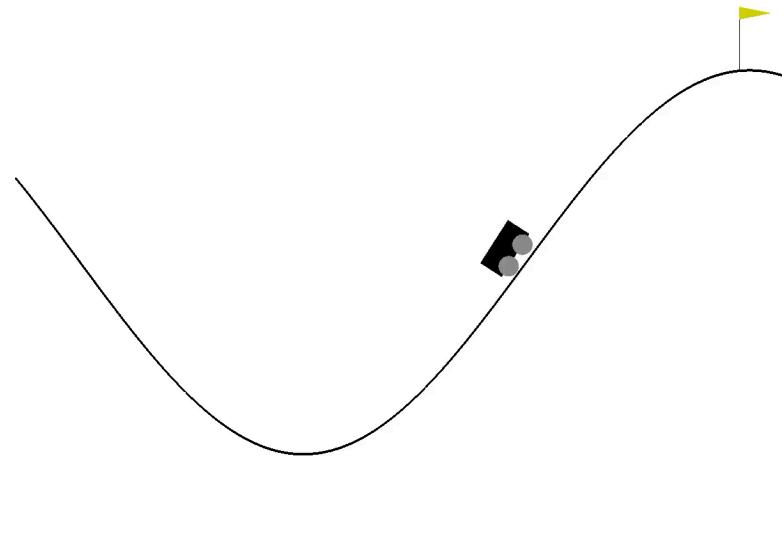
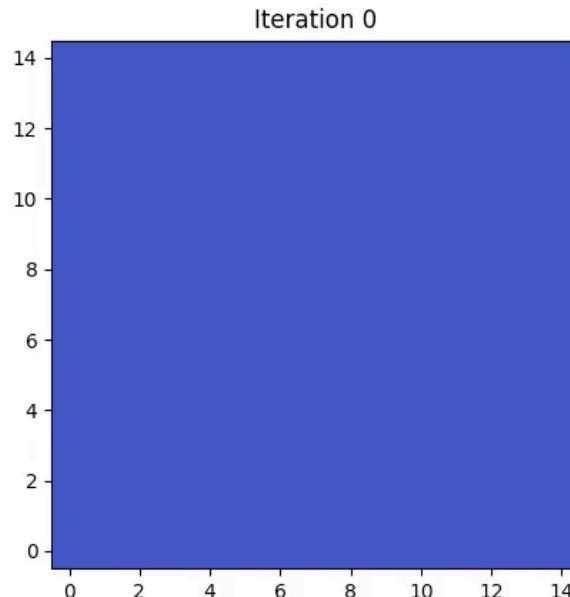
- Discretization
- Lookahead policies
- *Examples*
- Guarantees
- Connection with function approximation

Mountain Car

nearest neighbor

#discrete values per state dimension: 20

#discrete actions: 2 (as in original env)

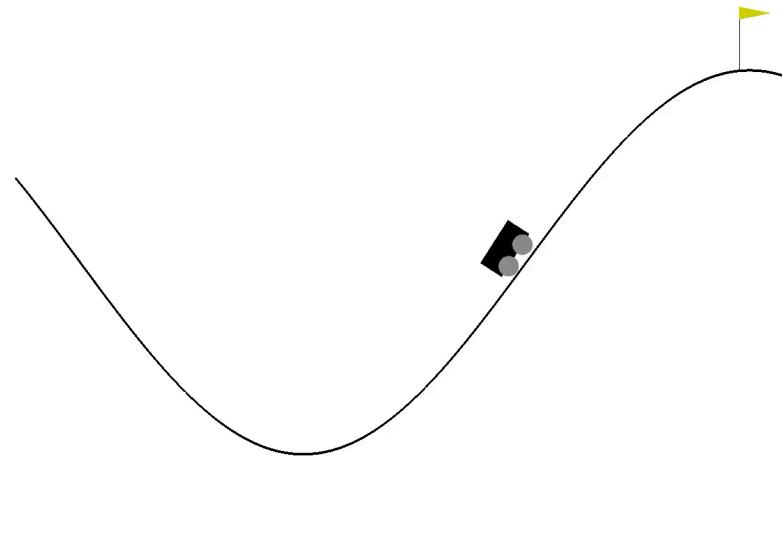
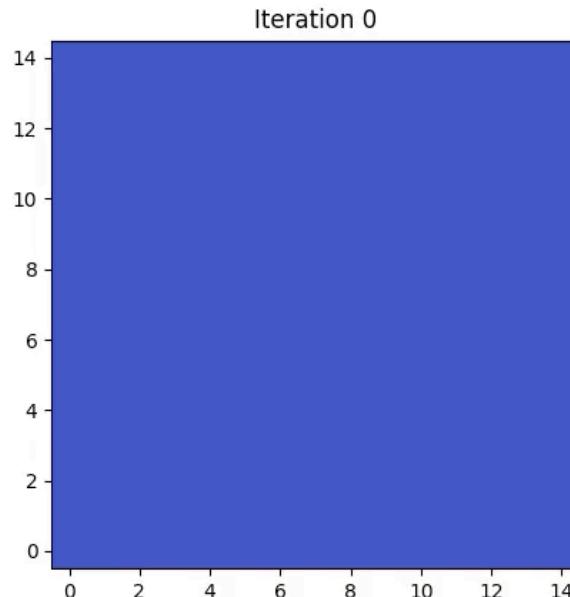


Mountain Car

nearest neighbor

#discrete values per state dimension: 150

#discrete actions: 2 (as in original env)

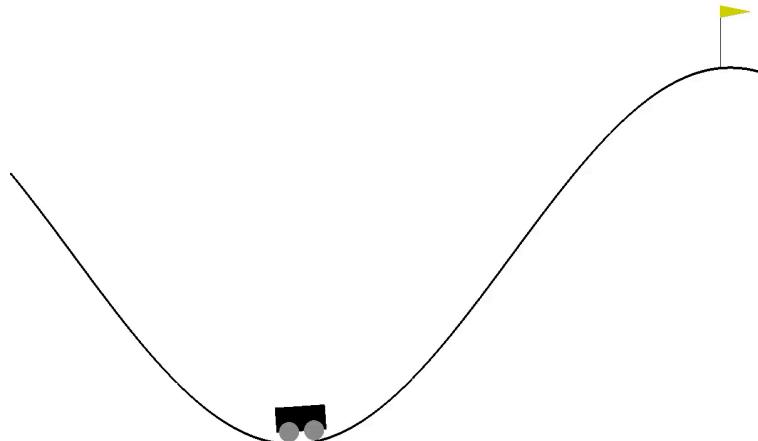
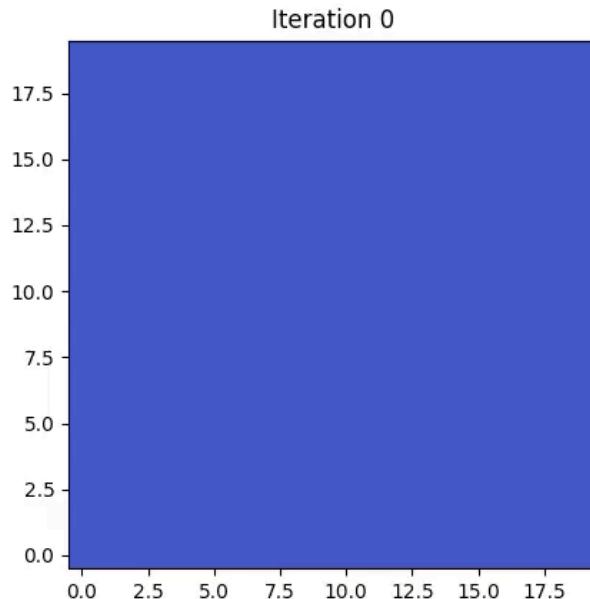


Mountain Car

linear

#discrete values per state dimension: 20

#discrete actions: 2 (as in original env)



Outline

- Discretization
- Lookahead policies
- Examples
- ***Guarantees***
- Connection with function approximation

Discretization Quality Guarantees

- Typical guarantees:
 - Assume: smoothness of cost function, transition model
 - For $h \rightarrow 0$, the discretized value function will approach the true value function
- To obtain guarantee about resulting policy, combine above with a general result about MDP's:
 - One-step lookahead policy based on value function V which is close to V^* is a policy that attains value close to V^*

Quality of Value Function Obtained from Discrete MDP: Proof Techniques

- Chow and Tsitsiklis, 1991:
 - Show that one discretized back-up is close to one “complete” back-up + then show sequence of back-ups is also close
- Kushner and Dupuis, 2001:
 - Show that sample paths in discrete stochastic MDP approach sample paths in continuous (deterministic) MDP [also proofs for stochastic continuous, bit more complex]
- Function approximation based proof (see later slides for what is meant with “function approximation”)
 - Great descriptions: Gordon, 1995; Tsitsiklis and Van Roy, 1996

Example result (Chow and Tsitsiklis,1991)**

A.1: $|g(x, u) - g(x', u')| \leq K \|(x, u) - (x', u')\|_\infty$,
for all $x, x' \in S$ and $u, u' \in C$;

A.2: $|P(y | x, u) - P(y' | x', u')| \leq K \|(y, x, u) - (y', x', u')\|_\infty$, for all $x, x', y, y' \in S$ and $u, u' \in C$;

A.3: for any $x, x' \in S$ and any $u' \in U(x')$, there exists some $u \in U(x)$ such that $\|u - u'\|_\infty \leq K \|x - x'\|_\infty$;

A.4: $0 \leq P(y | x, u) \leq K$ and $\int_S P(y | x, u) dy = 1$,
for all $x, y \in S$ and $u \in C$.

Theorem 3.1: There exist constants K_1 and K_2 (depending only on the constant K of assumptions A.1–A.4) such that for all $h \in (0, 1/2K]$ and all $J \in \mathcal{B}(S)$

$$\|TJ - \tilde{T}_h J\|_\infty \leq (K_1 + \alpha K_2 \|J\|_S) h. \quad (3.6)$$

Furthermore,

$$\|J^* - \tilde{J}_h^*\|_\infty \leq \frac{1}{1 - \alpha} (K_1 + \alpha K_2 \|J^*\|_S) h. \quad (3.7)$$

Outline

- Discretization
- Lookahead policies
- Examples
- Guarantees
- *Connection with function approximation*

Value Iteration with Function Approximation

Alternative interpretation of the discretization methods:

Start with $V_0^*(s) = 0$ for all s .

For $i = 0, 1, \dots, H-1$

for all states $s \in \bar{S}$,
(\bar{S} is the discrete state set)

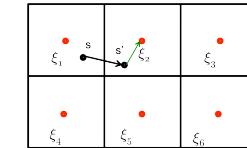
$$V_{i+1}^*(s) \leftarrow \max_a \sum_{s'} T(s, a, s') [R(s, a, s') + \hat{V}_i^*(s')]$$

with:

$$\hat{V}_i^*(s') = \sum_j P(\xi_j; s') V_i^*(\xi_j)$$

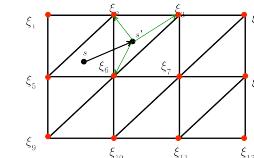
0'th Order Function Approximation

$$P(\xi_i; s') = \begin{cases} 1 & \text{if } \xi_i = \arg \min_{\xi \in \{\xi_1, \dots, \xi_N\}} \|s - \xi\| \\ 0 & \text{otherwise} \end{cases}$$



1st Order Function Approximation

$$P(\xi_i; s') \text{ such that } s' = \sum_{i=1}^N P(\xi_i; s') \xi_i$$



Discretization as Function Approximation

- Nearest neighbor discretization:
 - builds piecewise constant approximation of value function
- Stochastic transition onto nearest neighbors:
 - n-linear function approximation
 - Kuhn: piecewise (over “triangles”) linear approximation of value function

Continuous time**

- One might want to discretize time in a variable way such that one discrete time transition roughly corresponds to a transition into neighboring grid points/regions
- Discounting: $\exp(-\beta\delta t)$
 δt depends on the state and action

See, e.g., Munos and Moore, 2001 for details.

Note: Numerical methods research refers to this connection between time and space as the CFL (Courant Friedrichs Levy) condition. Googling for this term will give you more background info.

!! 1 nearest neighbor tends to be especially sensitive to having the correct match [Indeed, with a mismatch between time and space 1 nearest neighbor might end up mapping many states to only transition to themselves no matter which action is taken.]