Review: What does an Operating System do?

- Silberschatz and Gavin:
  "An OS is Similar to a government"
  - Begs the question: does a government do anything useful by itself?
- Coordinator and Traffic Cop:
  - Manages all resources
  - Settles conflicting requests for resources
  - Prevent errors and improper use of the computer
- Facilitator:
  - Provides facilities that everyone needs
  - Standard Libraries, Windowing systems
  - Make application programming easier, faster, less error-prone
- Some features reflect both tasks:
  - E.g. File system is needed by everyone (Facilitator)
  - But File system must be Protected (Traffic Cop)

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Review: Virtual Machine Abstraction

- Software Engineering Problem:
  - Turn hardware/software quirks ⇒ what programmers want/need
  - Optimize for convenience, utilization, security, reliability, etc...
- For Any OS area (e.g. file systems, virtual memory, networking, scheduling):
  - What’s the hardware interface? (physical reality)
  - What’s the application interface? (nicer abstraction)

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Review: Example of Address Translation

<table>
<thead>
<tr>
<th>Code</th>
<th>Data</th>
<th>Heap</th>
<th>Stack</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
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</table>

Translation Map 1

Translation Map 2

Physical Address Space

8/31/05 Kubiatowicz CS162 ©UCB Fall 2005 Lec 2.3

8/31/05 Kubiatowicz CS162 ©UCB Fall 2005 Lec 2.4
Review: Dual Mode Operation

- **Hardware** provides at least two modes:
  - "Kernel" mode (or "supervisor" or "protected")
  - "User" mode: Normal programs executed
- Some instructions/ops prohibited in user mode:
  - Example: cannot modify page tables in user mode
    - Attempt to modify ⇒ Exception generated
- Transitions from user mode to kernel mode:
  - System Calls, Interrupts, Other exceptions

Goals for Today

- History of Operating Systems
  - Really a history of resource-driven choices
- Operating Systems Structures
- Operating Systems Organizations

Note: Some slides and/or pictures in the following are adapted from slides ©2005 Silberschatz, Galvin, and Gagne

Moore's law change

<table>
<thead>
<tr>
<th></th>
<th>1981</th>
<th>2005</th>
<th>Factor</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>CPU MHz, Cycles/inst</strong></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>DRAM capacity</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Disk capacity</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Net bandwidth</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td># addr bits</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>#users/machine</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Price</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th></th>
<th>1981</th>
<th>2005</th>
<th>Factor</th>
</tr>
</thead>
<tbody>
<tr>
<td>CPU MHz, Cycles/inst</td>
<td>10</td>
<td>3800</td>
<td>380</td>
</tr>
<tr>
<td>DRAM capacity</td>
<td>128KB</td>
<td>46B</td>
<td>32,768</td>
</tr>
<tr>
<td>Disk capacity</td>
<td>10MB</td>
<td>1TB</td>
<td>100,000</td>
</tr>
<tr>
<td>Net bandwidth</td>
<td>9600 b/s</td>
<td>1 Gb/s</td>
<td>110,000</td>
</tr>
<tr>
<td># addr bits</td>
<td>16</td>
<td>32</td>
<td>2</td>
</tr>
<tr>
<td>#users/machine</td>
<td>10s</td>
<td>≤ 1</td>
<td>≤ 0.1</td>
</tr>
<tr>
<td>Price</td>
<td>$25,000</td>
<td>$4,000</td>
<td>0.2</td>
</tr>
</tbody>
</table>

Typical academic computer 1981 vs 2005

Dawn of time

ENIAC: (1945—1955)

- "The machine designed by Drs. Eckert and Mauchly was a monstrosity. When it was finished, the ENIAC filled an entire room, weighed thirty tons, and consumed two hundred kilowatts of power."
- http://ei.cs.vt.edu/~history/ENIAC.Richey.HTML
**History Phase 1 (1948–1970)**

**Hardware Expensive, Humans Cheap**

- When computers cost millions of $’s, optimize for more efficient use of the hardware!
  - Lack of interaction between user and computer

- User at console: one user at a time
- Batch monitor: load program, run, print

- Optimize to better use hardware
  - When user thinking at console, computer idle→BAD!
  - Feed computer batches and make users wait
  - Autograder for this course is similar
- **No protection**: what if batch program has bug?

---

**Core Memories (1950s & 60s)**

- Core Memory stored data as magnetization in iron rings
  - Iron “cores” woven into a 2-dimensional mesh of wires
  - Origin of the term “Dump Core”
  - Rumor that IBM consulted Life Saver company
- See: [http://www.columbia.edu/acis/history/core.html](http://www.columbia.edu/acis/history/core.html)

---

**History Phase 1½ (late 60s/early 70s)**

- **Data channels, Interrupts**: overlap I/O and compute
  - DMA – Direct Memory Access for I/O devices
  - I/O can be completed asynchronously

- **Multiprogramming**: several programs run simultaneously
  - Small jobs not delayed by large jobs
  - More overlap between I/O and CPU
  - Need memory protection between programs and/or OS

- **Complexity gets out of hand**:
  - Multics: announced in 1963, ran in 1969
    - Turing award lecture from Fernando Corbató (key researcher): “On building systems that will fail”
  - OS 360: released with 1000 known bugs (APARs)
    - “Anomalous Program Activity Report”

- **OS finally becomes important science**:
  - How to deal with complexity???
- UNIX based on Multics but vastly simplified

---

**A Multics System (Circa 1976)**

- The 6180 at MIT IPC, skin doors open, circa 1976:
  - “We usually ran the machine with doors open so the operators could see the AQ register display, which gave you an idea of the machine load, and for convenient access to the EXECUTE button, which the operator would push to enter BOS if the machine crashed.”
- [http://www.multicians.org/multics-stories.html](http://www.multicians.org/multics-stories.html)
Early Disk History

1973: 1.7 Mbit/sq. in 140 MBytes
1979: 7.7 Mbit/sq. in 2,300 MBytes


History Phase 2 (1970 – 1985)
Hardware Cheaper, Humans Expensive

- Computers available for tens of thousands of dollars instead of millions
- OS Technology maturing/stabilizing
- Interactive timesharing:
  - Use cheap terminals (~$1000) to let multiple users interact with the system at the same time
  - Sacrifice CPU time to get better response time
  - Users do debugging, editing, and email online
- Problem: Thrashing
  - Performance very non-linear response with load
  - Thrashing caused by many factors including
    » Swapping, queueing

Administrivia: Time for Project Signup

- Project Signup:
  - The group signup page is now working
  - Only submit once per group!
  - Everyone in group must have logged into their cs162-xx accounts before you register the group
  - Make sure that you select at least 2 potential sections
  - Due date: Wednesday 9/7 by 11:59pm
- Next week, go to your pre-assigned section

<table>
<thead>
<tr>
<th>Section</th>
<th>Time</th>
<th>Location</th>
<th>TA</th>
</tr>
</thead>
<tbody>
<tr>
<td>101</td>
<td>Tu 1:00-2:00P</td>
<td>310 Hearst Mining</td>
<td>Dominic</td>
</tr>
<tr>
<td>102</td>
<td>W 10:00-11:00A</td>
<td>2 Evans</td>
<td>Rajesh</td>
</tr>
<tr>
<td>103</td>
<td>W 11:00-12:00P</td>
<td>85 Evans</td>
<td>Rajesh</td>
</tr>
<tr>
<td>104</td>
<td>W 1:00-2:00P</td>
<td>85 Evans</td>
<td>Chris</td>
</tr>
<tr>
<td>105</td>
<td>W 2:00-3:00P</td>
<td>85 Evans</td>
<td>Chris</td>
</tr>
</tbody>
</table>

Administrivia (2)

- Cs162-xx accounts:
  - Make sure you got an account form
  - If you haven’t logged in yet, you need to do so
- Clarification of late policy:
  - I need to enforce the midnight deadlines
- Nachos readers:
  - Available Friday from Northside Copy Central
  - Includes printouts of all of the code
- Book Status:
  - Campus bookstore is out of the books – new shipment scheduled next week
  - Ned’s has 19 of the Silbershatz books and one used copy of the Free BSD book
- Web cast archives available off lectures page
  - Just click on the title of a lecture for webcast
  - Only works for lectures that I have already given!
History Phase 3 (1981—)

Hardware Very Cheap, Humans Very Expensive

- Computer costs $1K, Programmer costs $100K/year
  - If you can make someone 1% more efficient by giving them a computer, it's worth it!
  - Use computers to make people more efficient
- Personal computing:
  - Computers cheap, so give everyone a PC
- Limited Hardware Resources Initially:
  - OS becomes a subroutine library
  - One application at a time (MSDOS, CP/M, ...)
- Eventually PCs become powerful:
  - OS regains all the complexity of a “big” OS
  - multiprogramming, memory protection, etc (NT, OS/2)
- Question: As hardware gets cheaper does need for OS go away?

History Phase 3 (con't)

Graphical User Interfaces

- CS160 ⇒ All about GUIs
- Xerox Star: 1981
  - Originally a research project (Alto)
  - First “mice”, “windows”
- Apple Lisa/Machintosh: 1984
  - “Look and Feel” suit 1988
- Microsoft Windows:
  - Win 1.0 (1985)
  - Win 3.1 (1990)
  - Win 95 (1995)
  - Win NT (1993)
  - Win XP (2001)

History Phase 4 (1989—): Distributed Systems

- Networking (Local Area Networking)
  - Different machines share resources
  - Printers, File Servers, Web Servers
  - Client – Server Model
- Services
  - Computing
  - File Storage

History Phase 4 (con't): Distributed Systems

- Networking (Local Area Networking)
  - Different machines share resources
  - Printers, File Servers, Web Servers
  - Client – Server Model
- Services
  - Computing
  - File Storage

History Phase 5 (1995—): Mobile Systems

- Ubiquitous Mobile Devices
  - Laptops, PDAs, phones
  - Small, portable, and inexpensive
    - Recently twice as many smart phones as PDAs
    - Many computers/person!
  - Limited capabilities (memory, CPU, power, etc...)
- Wireless/Wide Area Networking
  - Leveraging the infrastructure
  - Huge distributed pool of resources extend devices
  - Traditional computers split into pieces. Wireless keyboards/mice, CPU distributed, storage remote
- Peer-to-peer systems
  - Many devices with equal responsibilities work together
  - Components of “Operating System” spread across globe
CITRIS's Model: A Societal Scale Information System

- Center for Information Technology Research in the Interest of Society
- The Network is the OS
  - Functionality spread throughout network
  - Scalable, Reliable, Secure Services

Moore's Law Reprise: My new X41 tablet

<table>
<thead>
<tr>
<th></th>
<th>1981</th>
<th>2005</th>
<th>2005 Laptop</th>
</tr>
</thead>
<tbody>
<tr>
<td>CPU MHz, Cycles/inst</td>
<td>10</td>
<td>3800</td>
<td>1500</td>
</tr>
<tr>
<td></td>
<td>3–10</td>
<td>0.25–0.5</td>
<td>0.25–0.5</td>
</tr>
<tr>
<td>DRAM capacity</td>
<td>128KB</td>
<td>4GB</td>
<td>1.5GB</td>
</tr>
<tr>
<td>Disk capacity</td>
<td>10MB</td>
<td>1TB</td>
<td>60GB</td>
</tr>
<tr>
<td>Net bandwidth</td>
<td>9600 b/s</td>
<td>1 Gb/s</td>
<td>1 Gb/s (wired)</td>
</tr>
<tr>
<td></td>
<td></td>
<td>54 Mb/s (wireless)</td>
<td>2 Mb/s (wide-area)</td>
</tr>
<tr>
<td># addr bits</td>
<td>16</td>
<td>32</td>
<td>32</td>
</tr>
<tr>
<td># users/machine</td>
<td>10s</td>
<td>≤ 1</td>
<td>≤ ½</td>
</tr>
<tr>
<td>Price</td>
<td>$25,000</td>
<td>$4,000</td>
<td>$2000</td>
</tr>
</tbody>
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Migration of Operating-System Concepts and Features

Microprocessors
Minicomputers
Mainframes
Supercomputers

Compare: Performance Trends (from CS152)
### History of OS: Summary

- Change is continuous and OSs should adapt
  - Not: look how stupid batch processing was
  - But: Made sense at the time
- Situation today is much like the late 60s
  - Small OS: 100K lines
  - Large OS: 10M lines (5M for the browser!)
  - 100-1000 people-years
- Complexity still reigns
  - NT under development from early 90's to late 90's
    » Still doesn't work very well
  - Jury still out on Windows 2000/XP
  - Windows "Longhorn" delayed many times
    » Latest release date of 2006+
    » Promised by removing some of the intended technology
- CS162: understand OSs to simplify them

### Operating Systems Components
**(What are the pieces of the OS)**

- Process Management
- Main-Memory Management
- I/O System management
- File Management
- Networking
- User Interfaces

### Operating System Services
**(What things does the OS do?)**

- Services that (more-or-less) map onto components
  - Program execution
    » How do you execute concurrent sequences of instructions?
  - I/O operations
    » Standardized interfaces to extremely diverse devices
  - File system manipulation
    » How do you read/write/preserve files?
    » Looming concern: How do you even find files???
  - Communications
    » Networking protocols/Interface with CyberSpace?
- Cross-cutting capabilities
  - Error detection & recovery
  - Resource allocation
  - Accounting
  - Protection
System Calls (What is the API)

- See Chapter 2 of 7th edition or Chapter 3 of 6th

Operating Systems Structure (What is the organizational Principle?)

- Simple
  - Only one or two levels of code
- Layered
  - Lower levels independent of upper levels
- Microkernel
  - OS built from many user-level processes
- Modular
  - Core kernel with Dynamically loadable modules

Simple Structure

- MS-DOS - written to provide the most functionality in the least space
  - Not divided into modules
  - Interfaces and levels of functionality not well separated

UNIX: Also "Simple" Structure

- UNIX - limited by hardware functionality
- Original UNIX operating system consists of two separable parts:
  - Systems programs
  - The kernel
    » Consists of everything below the system-call interface and above the physical hardware
    » Provides the file system, CPU scheduling, memory management, and other operating-system functions;
    » Many interacting functions for one level
### UNIX System Structure

- **User Mode**
  - Applications: (the users)
  - Standard Libs: shells and commands, compilers and interpreters, system libraries

- **Kernel Mode**
  - Kernel: system-call interface to the kernel
    - Signals: terminal handling
    - Character I/O
    - System file
    - Disk and tape drivers
    - CPU scheduling
    - Page replacement
    - Demand paging
    - Virtual memory

- **Hardware**
  - Terminal controllers, terminals
  - Device controllers, disks and tapes
  - Memory controllers, physical memory

### Layered Structure

- Operating system is divided many layers (levels)
  - Each built on top of lower layers
  - Bottom layer (layer 0) is hardware
  - Highest layer (layer N) is the user interface

- Each layer uses functions (operations) and services of only lower-level layers
  - Advantage: modularity ⇒ Easier debugging/Maintenance
  - Not always possible: Does process scheduler lie above or below virtual memory layer?
    - Need to reschedule processor while waiting for paging
    - May need to page in information about tasks

- Important: Machine-dependent vs independent layers
  - Easier migration between platforms
  - Easier evolution of hardware platform
  - Good idea for you as well!

### Layered Operating System

- Moves as much from the kernel into "user" space
  - Small core OS running at kernel level
  - OS Services built from many independent user-level processes

- Communication between modules with message passing

- Benefits:
  - Easier to extend a microkernel
  - Easier to port OS to new architectures
  - More reliable (less code is running in kernel mode)
  - Fault Isolation (parts of kernel protected from other parts)
  - More secure

- Detriments:
  - Performance overhead severe for naïve implementation
**Modules-based Structure**

- Most modern operating systems implement modules
  - Uses object-oriented approach
  - Each core component is separate
  - Each talks to the others over known interfaces
  - Each is loadable as needed within the kernel
- Overall, similar to layers but with more flexible

**Operating System Design Goals**
(What is this OS trying to achieve?)

- $2000$ price point?
- Fault tolerance/Fast failover/High Availability?
- High Performance?
- Real Time Capable?

---

**Implementation Issues**
(How is the OS implemented?)

- Policy vs. Mechanism
  - Policy: *What* do you want to do?
  - Mechanism: *How* are you going to do it?
  - Should be separated, since policies change
- Algorithms used
  - Linear, Tree-based, Log Structured, etc...
- Event models used
  - threads vs event loops
- Backward compatibility issues
  - Very important for Windows 2000
- System generation/configuration
  - How to make generic OS fit on specific hardware

**Conclusion**

- Rapid Change in Hardware Leads to changing OS
  - Batch ⇒ Multiprogramming ⇒ Timeshare ⇒ Graphical UI ⇒ Ubiquitous Devices ⇒ Cyberspace/Metaverse/??
- OS features migrated from mainframes ⇒ PCs
- Standard Components and Services
  - Process Control
  - Main Memory
  - I/O
  - File System
  - UI
- Policy vs Mechanism
  - Crucial division: not always properly separated!
- Complexity is always out of control
  - However, “Resistance is NOT Useless!”