

CS 61B Lab 14
November 30–December 1, 2010

Goal: This lab will help you get comfortable with amortized analysis.

Consider a binary counter implemented as an array of booleans. Each digit of the counter is represented by a boolean in the array. The least significant bit is at index zero, and has a value of one. Each subsequent digit is worth twice the previous digit.

There is only one operation that changes the data structure: incrementing the counter. Therefore, the array increases like this.

```

index      ... 5 4 3 2 1 0
-----
digits     ... | 0 | 0 | 0 | 0 | 0 | 0 | ----\
          -----
          increment()
          <-----/
          ... | 0 | 0 | 0 | 0 | 0 | 1 | ----\
          -----
          increment()
          <-----/
          ... | 0 | 0 | 0 | 0 | 1 | 0 | ----\
          -----
          increment()
          <-----/
          ... | 0 | 0 | 0 | 0 | 1 | 1 | ----\
          -----
          increment()
          <-----/
          ... | 0 | 0 | 0 | 1 | 0 | 0 |
          -----
          .
          .
          .
          -----
          ... | 0 | 1 | 1 | 1 | 1 | 1 | ----\
          -----
          increment()
          <-----/
          ... | 1 | 0 | 0 | 0 | 0 | 0 |
          -----
    
```

Here's the code for incrementing the counter.

```

public class Counter {
    private boolean[] digits;

    public Counter(int size) {
        digits = new boolean[size];
    }

    public void increment() {
        int i = 0;
        while (digits[i]) {
            digits[i] = false;
            i++;
        }
        digits[i] = true;
    }
}
    
```

Obviously, the inner loop runs d times if the first d digits are ones, and d can be arbitrarily large. Your job is to prove that `increment()` runs in $O(1)$ amortized time.

Assume that:

- The cost (in time) of testing a bit, changing it, and incrementing i is at most one dollar (for all three computations together).
- The cost for the method call to `increment()`, plus initializing i , is exactly one dollar. This cost does not include the lines marked with asterisks (*).

Your job is to figure out:

- The amortized cost of an `increment()` operation, in dollars. Give the smallest amortized cost possible (given the information you have) that will ensure that your bank account never goes below zero.
- An explanation for why this analysis is correct. In other words, show that your bank balance will never go below zero.
- Suppose we add a `decrement()` operation, which decreases the counter by one. Why is the amortized cost of the counter operations no longer in $O(1)$?

Here's your big hint: simply by looking at the counter, you can tell very easily exactly how many dollars are in the bank.

If you're having trouble figuring it out, take a guess at the right dollar amount, then modify the code in `Counter.java` to keep track of the bank balance. See if you can see a connection between the counter and the bank balance.

Check-off

- 2 points: How many dollars are in the bank for any given counter value?
- 1 point: What is the amortized cost of `increment()`, in dollars? (Give the smallest answer that can be justified by the information above.)
- 1 point: Explain why adding a `decrement()` operation invalidates the analysis.