

## Lecture 28: Central Path

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In this lecture, we will go over the central path following method and the runtime complexity. We will also go over the preliminaries of Primal-Dual conic problems.

## 1 Path Following

The Path Following algorithm is used to solve standard constrained minimization problems of the following.

$$\begin{aligned} \min_x \langle c, x \rangle \\ \text{s.t. } x \in \mathcal{Q} \end{aligned} \quad (1)$$

Where  $\mathcal{Q}$  is a bounded closed convex set and we have a  $v$ -self-concordant barrier  $F$  on  $\mathcal{Q}$ . We will refer to the global minimum of (1) as  $c^*$ .

## 2 Central Path Method

We will solve (1) by tracing the "central path" as discussed in the previous lecture.

$$\begin{aligned} x^*(t) &= \arg \min_{x \in \text{dom} F} f(t; x) \\ f(t; x) &= t \langle c, x \rangle + F(x), \quad t \geq 0 \end{aligned} \quad (2)$$

We refer to  $x^*(t)$  as the "central path". If we take the gradient of  $f(t; x)$  and evaluate it at  $x^*(t)$ , it should be equal to 0, as per the first-order optimality conditions [1]. Thus, we get

$$\nabla_{x^*(t)} f(t; x) = tc + \nabla F(x^*(t)) = 0$$

As we can see from before,  $x^*(0) = \arg \min_{x \in \text{dom} F} F(x)$ . We will refer to  $x^*(0) = x_F^*$  as the "center" of the central path. In order to update/trace the central path (2), we will use the "Approximate Centering Condition" to do so. This is defined as

$$\lambda_{f(t, \cdot)}(x) = \|f'(t; x)\|_x^* = \|tc + \nabla F(x)\|_x^* \leq \beta \quad (3)$$

The  $\beta$  parameter is defined as the "centering parameter" as defined by Nesterov, and tends to be between  $[0, \frac{1}{2}]$ . Now we want to show that tracing the central path will result in a "close" approximation to  $c^*$ .

**Theorem 1.** 5.3.10 (Nesterov) [1]

For any  $t \geq 0$ ,

$$\langle c, x^*(t) \rangle - c^* \leq \frac{v}{t} \quad (4)$$

More generally, if  $x$  satisfies (3), then we have

$$\langle c, x \rangle - c^* \leq \frac{1}{t} \left( v + \frac{(\beta + \sqrt{v})\beta}{1 - \beta} \right) \quad (5)$$

When  $\beta = 0$ , (5) simplifies to (4).

An observation is that we often want  $F$  to have a small  $v$ -value. We also want to ensure that  $\beta$  does not change too much from one value of  $t$  to the next value of  $t$ .

Here we present two lemmas that presents how to increment  $t$  in an appropriate manner and establish that the growth in  $t$  is substantial (linear).

**Lemma 2.** 5.2.2 (Nesterov) *Let  $\tau \in [0, 1/2]$ . If  $x$  satisfies (3), then if*

$$|\gamma| \leq \tau - \tau^2 \left( 1 + \tau + \frac{\tau}{1 + \tau + \tau^2} \right)$$

then

$$\|t_t c + \nabla F(x_t)\|_{x_t}^* \leq \beta \tag{6}$$

Where  $t_t = t + \frac{\gamma}{\|c\|_x^*}$  and  $x_t = x - \frac{1}{1+\epsilon} [\nabla^2 F(x)]^{-1} (t_t c + \nabla F(x_t))$  and  $\epsilon = \frac{\lambda^2}{1+\lambda}$ ,  $\lambda = \|t_t c + \nabla F(x)\|_x^*$

**Lemma 3.** 5.3.2 (Nesterov)

$$\|c\|_x^* \leq \frac{1}{t} (\beta + \sqrt{v})$$

From these two lemmas we can derive the path following algorithm.

### 3 Algorithm

1. Set  $t_0 = 0$ . Choose  $\epsilon \geq 0, x_0 \in \text{dom}F$  such that  $\|\nabla F(x_0)\|_{x_0}^* \leq \beta$
2. For  $k \geq 0$

- $t_{k+1} = t_k + \frac{\gamma}{\|c\|_{x_k}^*}$
- $x_{k+1} = x_k - \frac{1}{1+\epsilon_k} [\nabla^2 F(x_k)]^{-1} (t_{k+1} c + \nabla F(x_k))$

Where  $\lambda_k = \|t_{k+1} c + \nabla F(x_k)\|_{x_k}^*$  and  $\epsilon_k = \frac{\lambda_k^2}{1+\lambda_k}$

3. Terminate when  $t_k \geq \frac{1}{\epsilon} \left( v + \frac{(\beta + \sqrt{v})\beta}{1-\beta} \right)$

It is not clear how to obtain  $x_0$ . We will go back to this later. Keep in mind that the termination condition is not dependent on the number of iterations  $k$ , but the value  $t_k$ , which doubles roughly (very fast!)

This begs the question, how many iterations do we need?

**Theorem 4.** 5.3.11 (Nesterov) [1] *The path following algorithm terminates after  $N$  steps where*

$$N \geq \mathcal{O} \left( \sqrt{v} * \log \left( \frac{v \|c\|_{x_F}^*}{\epsilon} \right) \right)$$

**Proof** Let  $r_0 = \|x_0 - x_F^*\|_{x_0}^* \leq \frac{\beta}{1-\beta}$ . Using the same terminology from Lemma 2, we get

$$\frac{\gamma}{t_1} = \|c\|_{x_0}^* \leq \frac{1}{1-r_0} \|c\|_{x_F^*}^* \leq \frac{1-\beta}{1-2\beta} \|c\|_{x_F^*}^*$$

Thus

$$t_1 \geq \frac{\gamma(1-2\beta)}{(1-\beta)\|c\|_{x_F^*}^*}$$

More generally,

$$t_k \geq \frac{\gamma(1-2\beta)}{(1-\beta)\|c\|_{x_F^*}^*} \left(1 + \frac{\gamma}{\beta + \sqrt{v}}\right)^{k-1}$$

□

An important thing to note is that the first term within the  $\mathcal{O}$  is the dominant term while second term in the bound tends to be very small. This still begs the question: How do we get the initial point  $x_0$ ? We could use Dampened Newton's method to find  $x_0$  but it is too slow so we will use another "path following" method, but in reverse to find  $x_0$ .

## 4 Auxillary Path Following

Let us define  $y^*(t) = \arg \min_{y \in \text{dom}F} [-t\langle \nabla F(y_0), y \rangle + F(y)]$  We start with  $t_0 = 1$ , and then we decrease. If we take the gradient of  $y^*(t)$  and evaluate at  $y_0$ , we will get  $\nabla_y y^*(t)|_{y=y_0} = -\nabla F(y_0) + \nabla F(y_0) = 0$ . Thus it is trivial to get the first point  $y_0$  in this auxillary path following scheme. Our resulting algorithm will resemble the path following method in reverse to get the initial point  $x_0$  for the original problem.[1]

### 4.1 Scheme

1.  $y_0 \in \text{dom}F, t_0 = 1$
2. For  $k \geq 0$ 
  - $t_{k+1} = t_k + \frac{\gamma}{\|\nabla F(y_0)\|_{y_k}^*}$
  - $y_{k+1} = y_k - \frac{1}{1+\epsilon_k} [\nabla^2 F(y_k)]^{-1} (-t_{k+1} \nabla F(y_0) + \nabla F(y_k))$

Where  $\lambda_k = \|t_{k+1} \nabla F(y_0) - \nabla F(y_k)\|_{y_k}^*$  and  $\epsilon_k = \frac{\lambda_k^2}{1+\lambda_k}$

3. Terminate when  $\|\nabla F(y_k)\|_{y_k}^* \leq \tau$ . Set  $\epsilon_k = \frac{\lambda_F(y_k)^2}{1+\lambda_F(y_k)}$  and  $\bar{x} = y_k - \frac{1}{1+\epsilon_k} [\nabla^2 F(y_k)]^{-1} (\nabla F(y_k))$

## 5 Primal-Dual Conic Problems

**Definition 5.** A set  $K \subset \mathbb{R}^n$  is a convex cone if  $K$  is a nonempty convex set with the property that  $x \in K, t \geq 0 \Rightarrow tx \in K$

**Definition 6.** A convex cone  $K$  is pointed if  $K$  does not contain lines (only rays)

### 5.1 Dual Cones

**Definition 7.**  $K^* = \{s \in \mathbb{R}^n \mid \langle s, x \rangle \geq 0 \forall x \in K\}$

Let us go through an example of finding the dual cone of  $K = \{x \in \mathbb{R}^n \mid x > 0\}$  (the nonnegative orthant cone). We can see that by simply plugging in  $e_1, e_2, \dots, e_n$ , where  $e_i$  is the  $i^{\text{th}}$  basis vector, into the definition of the dual cone of  $K$ , we will get  $K^* = \{x \in \mathbb{R}^n \mid x > 0\} = K$ . Because  $K = K^*$ , we call  $K$  self-dual.

**Theorem 8.** *Let  $K$  be a closed convex cone. Then*

1.  $K^*$  is a closed convex cone. Moreover,  $(K^*)^* = K$
2.  $K$  is pointed iff  $K^*$  has a nonempty interior (and vice versa)
3.  $s \in K^*$  is strictly positive on  $K$  (besides the 0 vector)  $\iff K(s) = \{x \in K \mid \langle s, x \rangle \leq 1\}$

## References

- [1] Y. Nesterov, *Lectures on Convex Optimization*. Springer Optimization and Its Applications.