

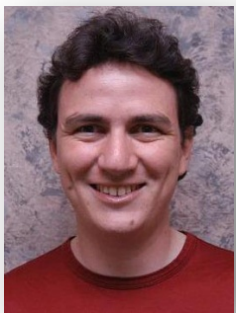
Thesis Defense

Parallel and Distributed Systems for Probabilistic Reasoning

Joseph E. Gonzalez



Thesis Committee:



Carlos Guestrin
University of
Washington & CMU



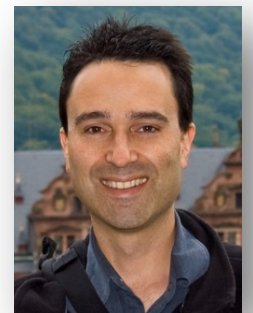
Guy Blelloch
CMU



David O'Hallaron
CMU



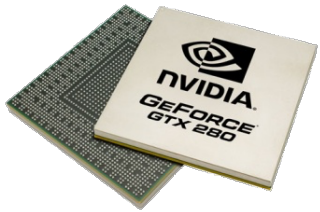
Alex Smola
CMU & Google



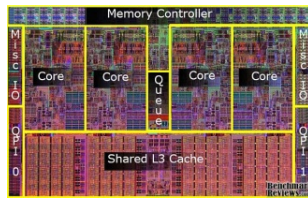
Jeff Bilmes
University of
Washington

*The foundations of **computation**
have changed ...*

New Parallel and Distributed Platforms



GPUs



Multicore



Clusters



Single Chip
Cloud Computers



Clouds

- New Opportunities
 - Increased processing and storage
- New Challenges
 - Parallel algorithm design and implementation

*The **scale** of
machine learning problems
is **exploding** ...*

The Age of Big Data



28 Million
Wikipedia Pages



1 Billion
Facebook Users



6 Billion
Flickr Photos



72 Hours a Minute
YouTube



WORLD U.S. N.Y. / REGION BUSINESS TEC

NEWS ANALYSIS

The Age of Big Data

By STEVE LOHR

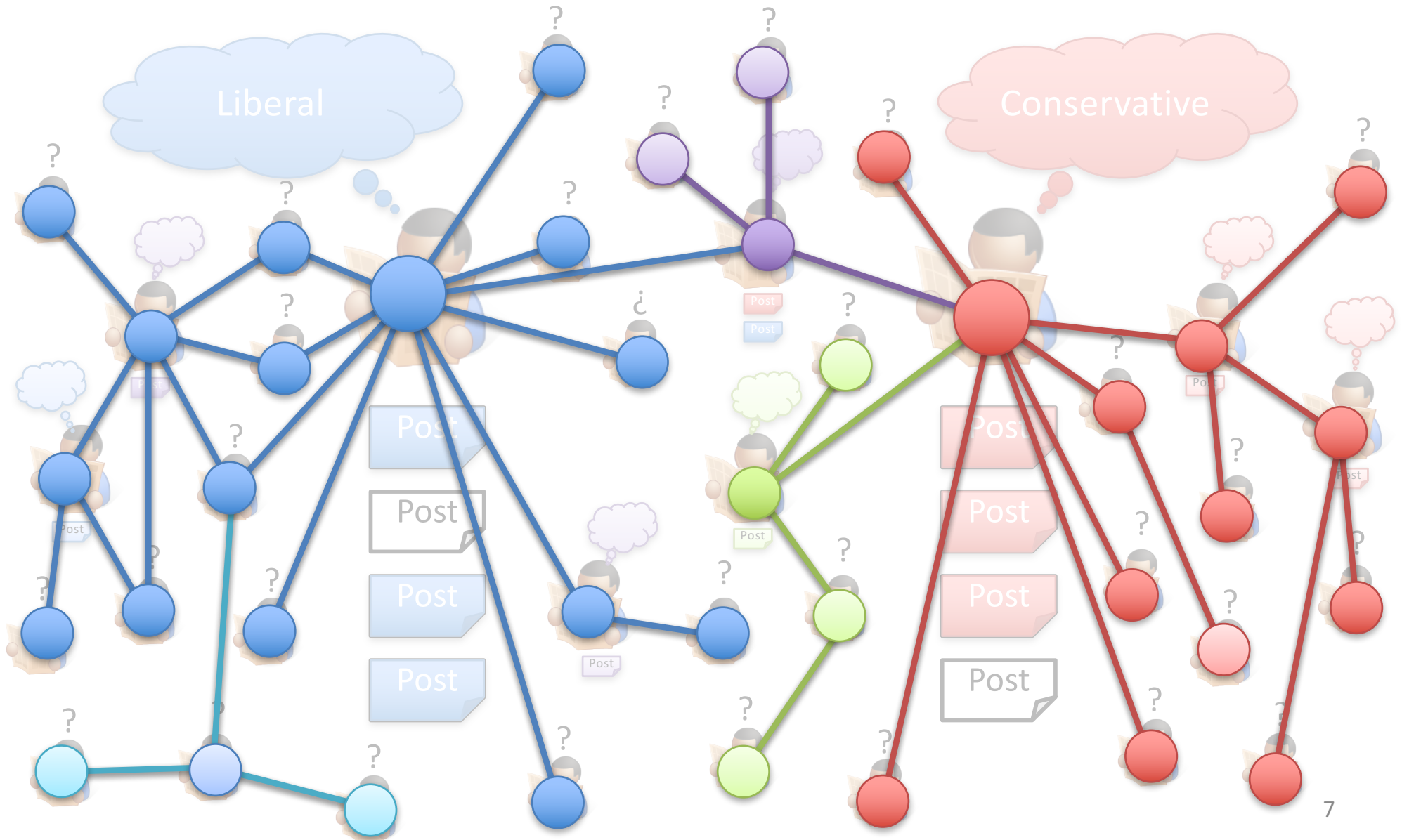
Published: February 11, 2012

“...growing at 50 percent a year...”

“... data a new class of economic asset,
like currency or gold.”

*Massive data provides
opportunities for
structured models...*

Example: *Estimate Political Bias*





facebook

flickr

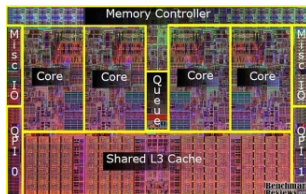
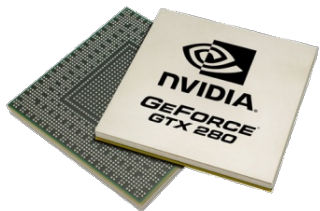
You Tube

Massive Structured Problems

Thesis:

*Parallel and Distributed Systems
for Probabilistic Reasoning*

Advances Parallel Hardware



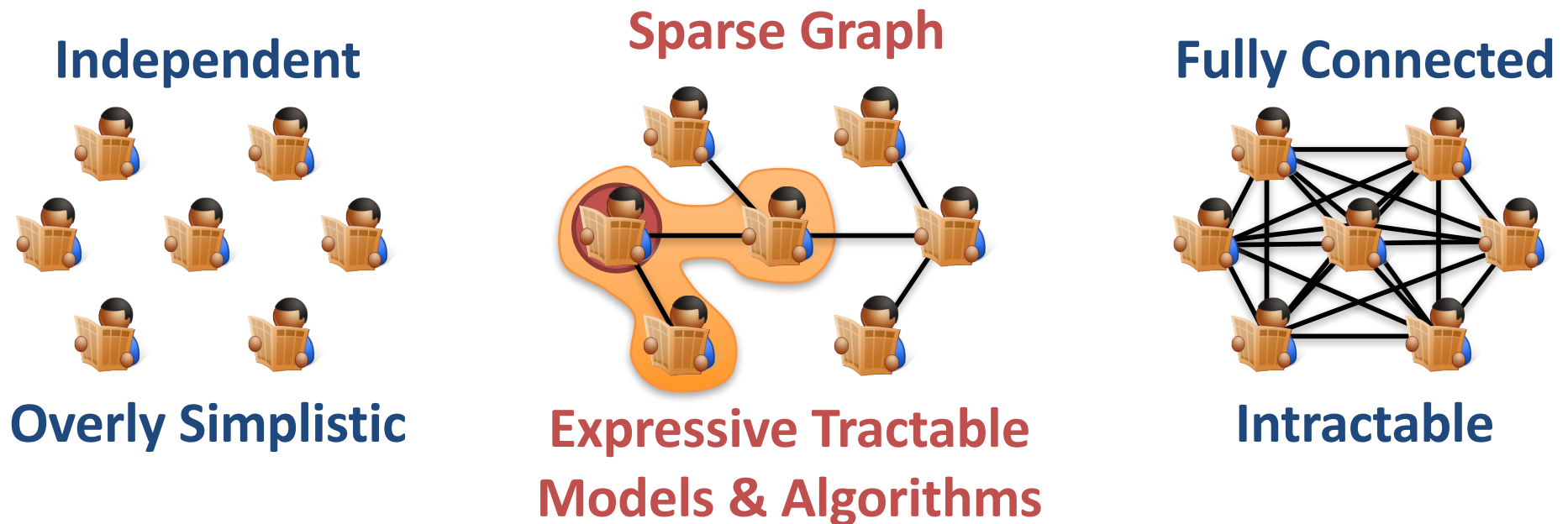
Thesis Statement: **GrAD** Methodology

*Efficient **parallel** and **distributed** systems for probabilistic reasoning:*

1. **Graphically** decompose *computational* and *statistical* dependencies
2. **Asynchronously** schedule computation
3. **Dynamically** identify and *prioritize* computation along critical paths

GrAD Methodology: Graphical

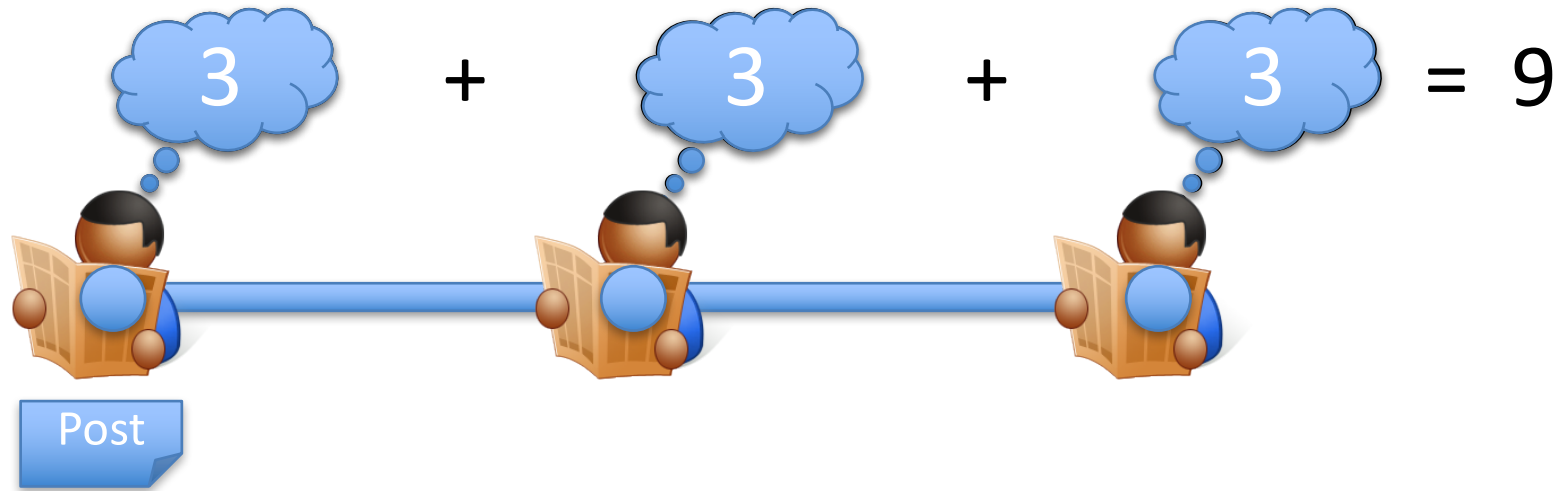
- Factor **statistical** and **computational** dependencies



- Improves **computational** and **statistical** efficiency
- Increases **parallelism**

Synchronous vs. Asynchronous

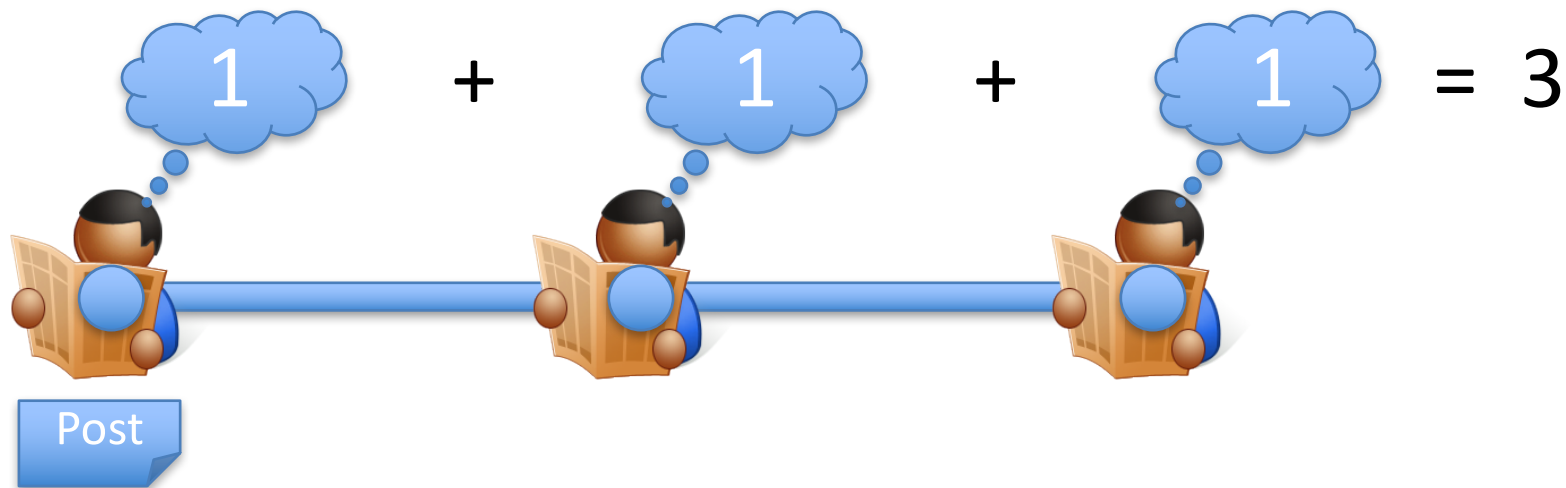
- **Synchronous:** *compute everything in parallel*



- Highly **parallel** – Maximum independent work
- Highly **inefficient** – Many wasted cycles

GrAD Methodology: Asynchronous

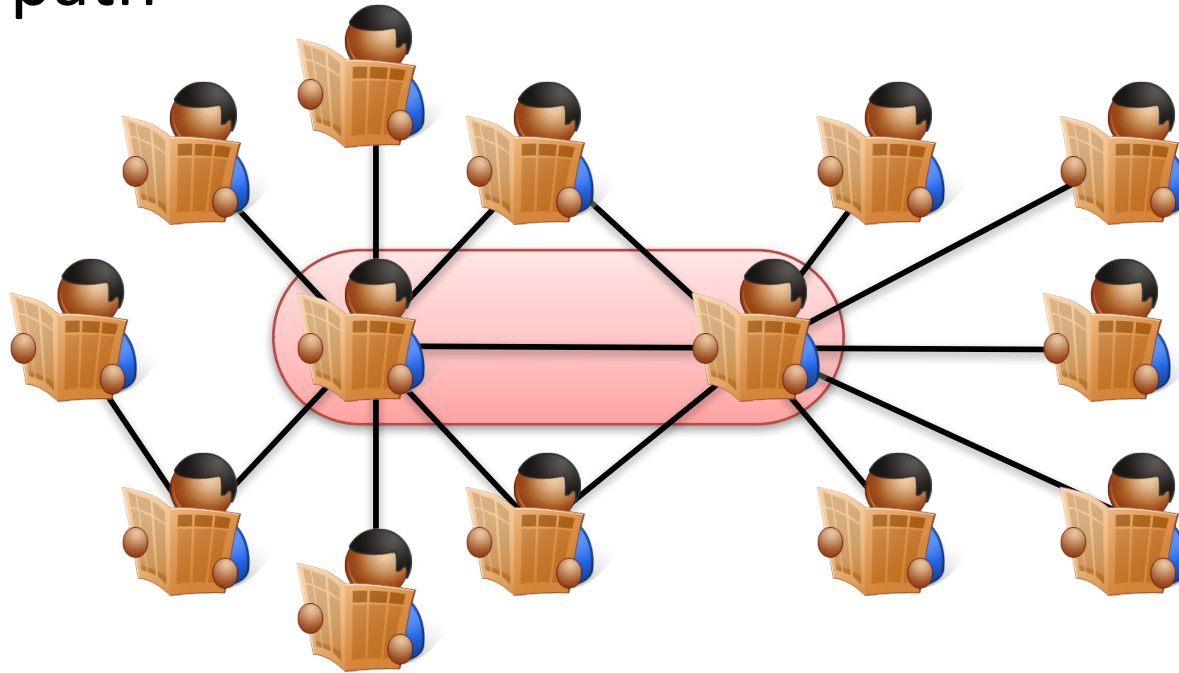
- Trigger computation as **new information arrives**



- Capture the flow of information:
 - More efficiently use network and processor resources
 - Guarantee algorithm correctness

GrAD Methodology: Dynamic

- Dynamically **identify** and **prioritize** computation along the critical path



- Focus computational resources where most effective:
 - Accelerated convergence
 - Increased work efficiency

We apply the GrAD methodology to

1. Probabilistic Graphical Models
2. **Parallel and Distributed Algorithms
for Probabilistic Inference**
3. **GraphLab & PowerGraph**



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You Tube

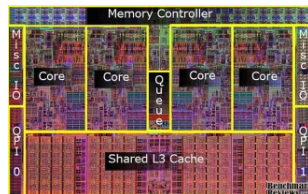
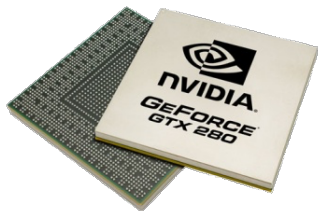
Massive Structured Problems

Probabilistic Graphical Models

Parallel and **Distributed** Algorithms
for Probabilistic **Inference**

GraphLab & PowerGraph

Advances Parallel Hardware



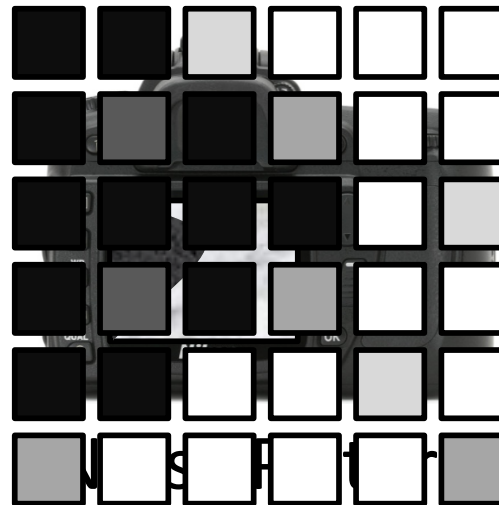
Probabilistic Graphical Models

Encode **Probabilistic** Structure



True Image

Noisy Pixels



Random Variables

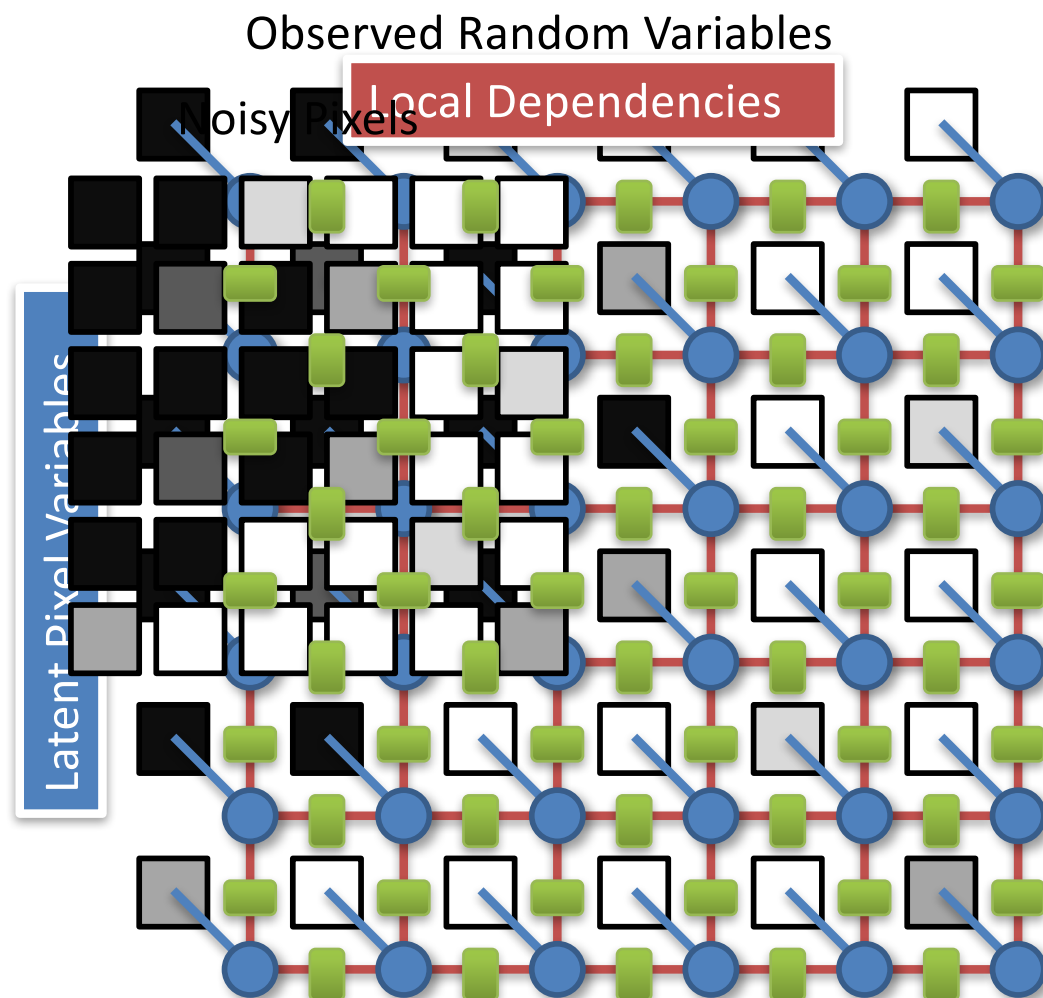
True *unobserved* values

Dependency Graph:

Represent dependencies

Parameters:

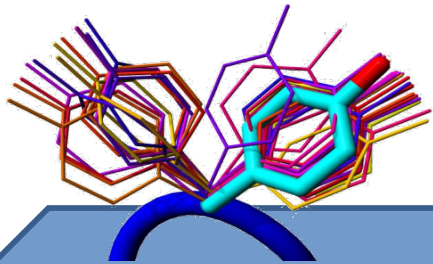
Characterize probabilities



$$\underbrace{P(X_1, \dots, X_n; \theta)}_{\text{Joint Probability}} \propto \prod_{\substack{(u,v) \in E \\ \text{Graph}}} \underbrace{f(X_u, X_v; \theta_{u,v})}_{\text{Factors}}$$

Graphical models provide a **common representation**

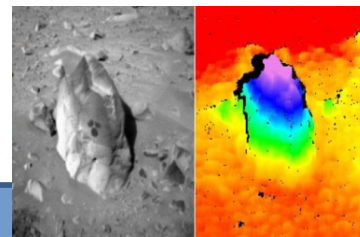
Protein Structure
Prediction



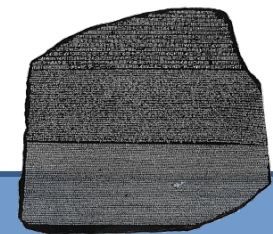
Movie
Recommendation



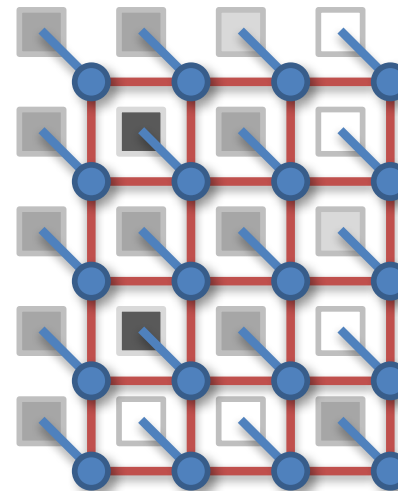
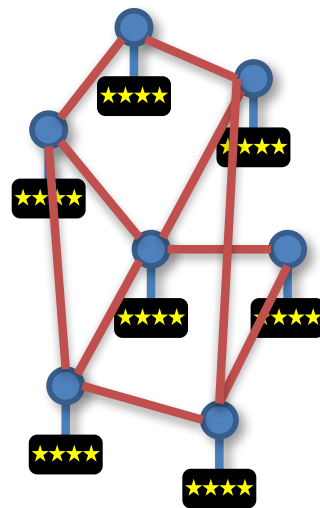
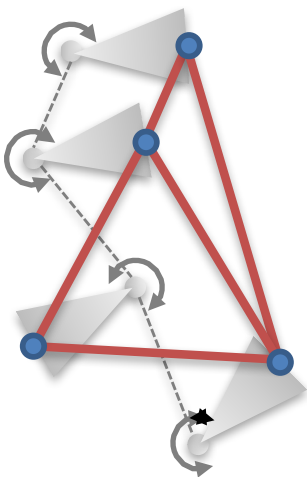
Computer
Vision



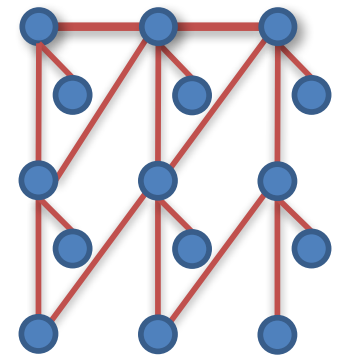
Machine
Translation



Probabilistic Graphical Models

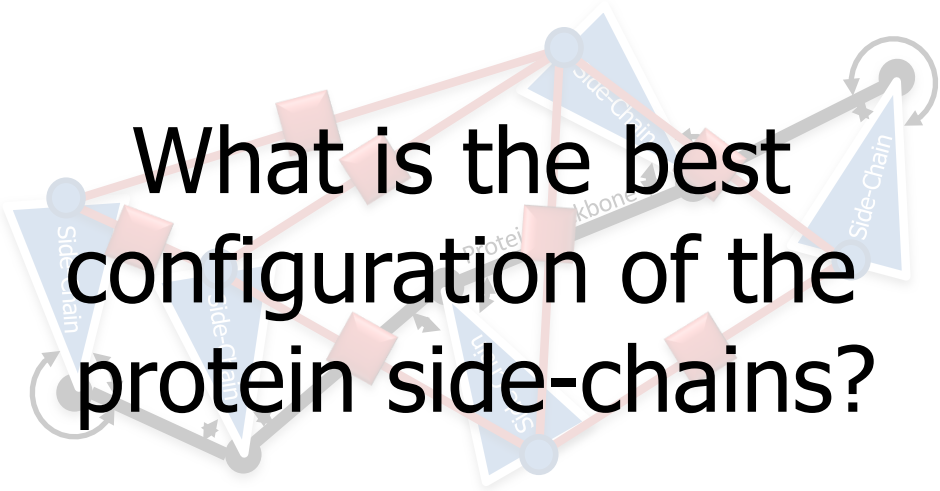


How are you?



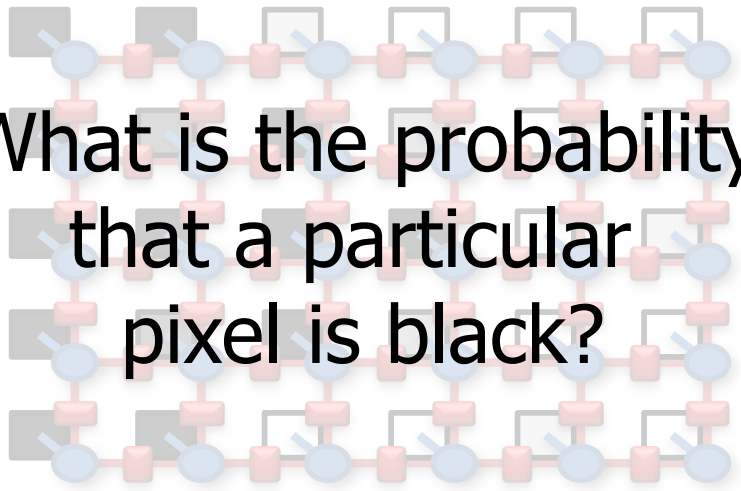
Probabilistic **Inference**

*Making **predictions** given the model structure and parameters*



What is the best configuration of the protein side-chains?

The diagram shows a protein backbone (grey line) with several side-chains (blue triangles) attached. Red arrows indicate the flexibility or movement of the side-chains. Labels include 'Side-Chain', 'Backbone', and 'Protein'.



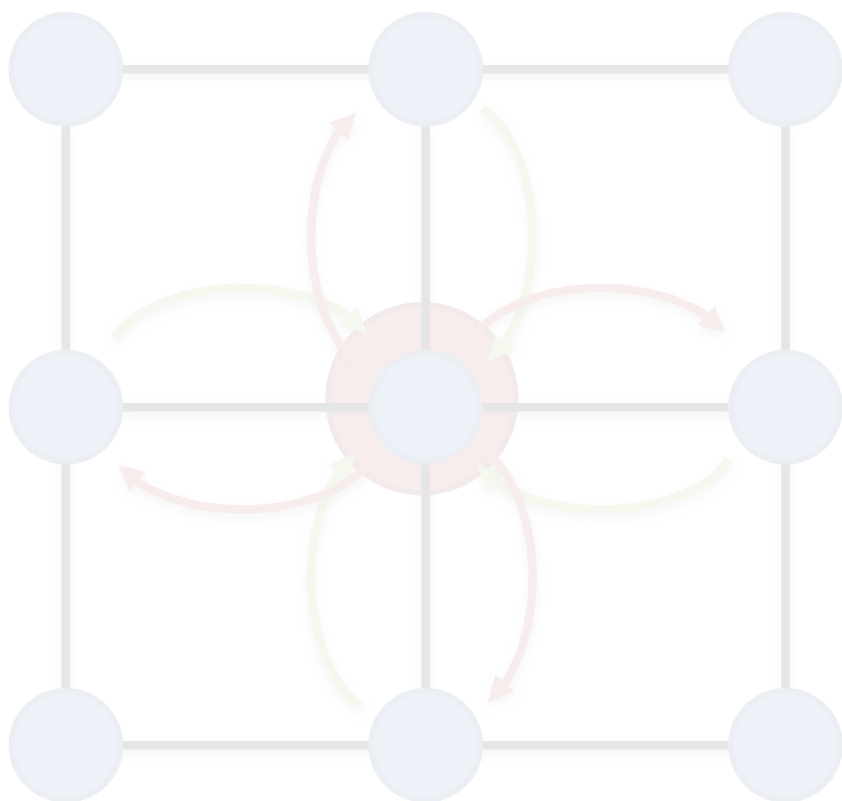
What is the probability that a particular pixel is black?

The diagram shows a 5x5 grid of pixels. Each pixel is represented by a small square. Some pixels are black, some are white, and some are grey. Red arrows indicate the relationships between adjacent pixels.

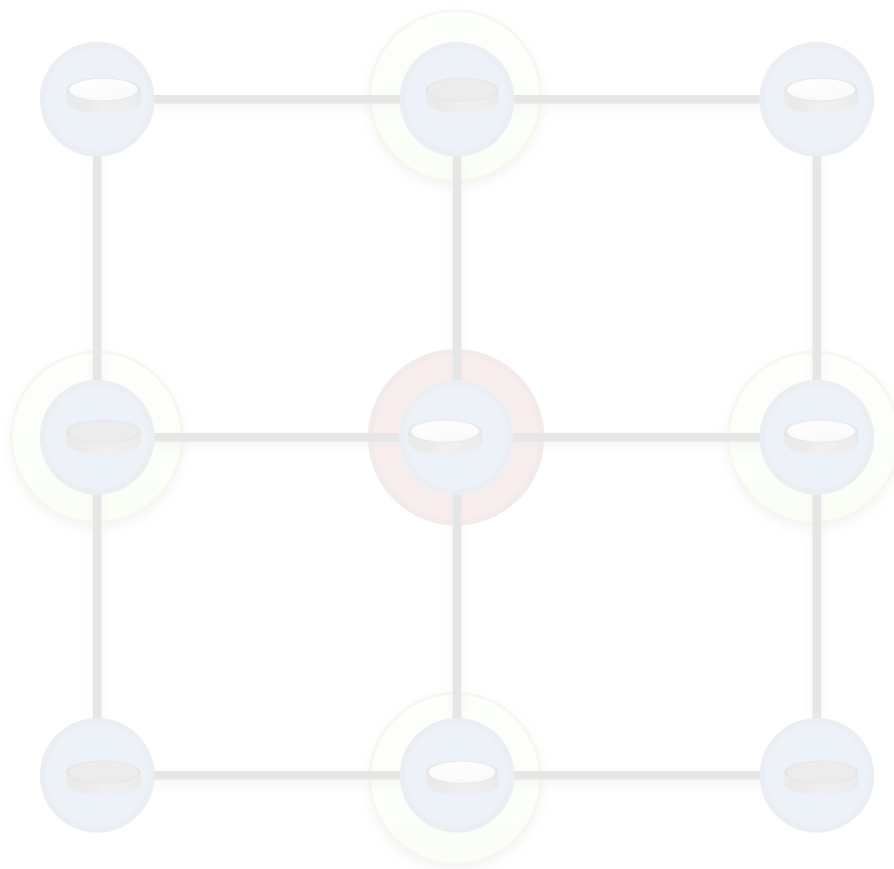
- **NP-complete** in general
 - Focus on *approximate* methods

Parallel and Distributed Algorithms for Probabilistic Inference

Belief Propagation



Gibbs Sampling



Parallel Belief Propagation

Joint Work With:

Yucheng Low

Carlos Guestrin

David O'Hallaron

Published Results

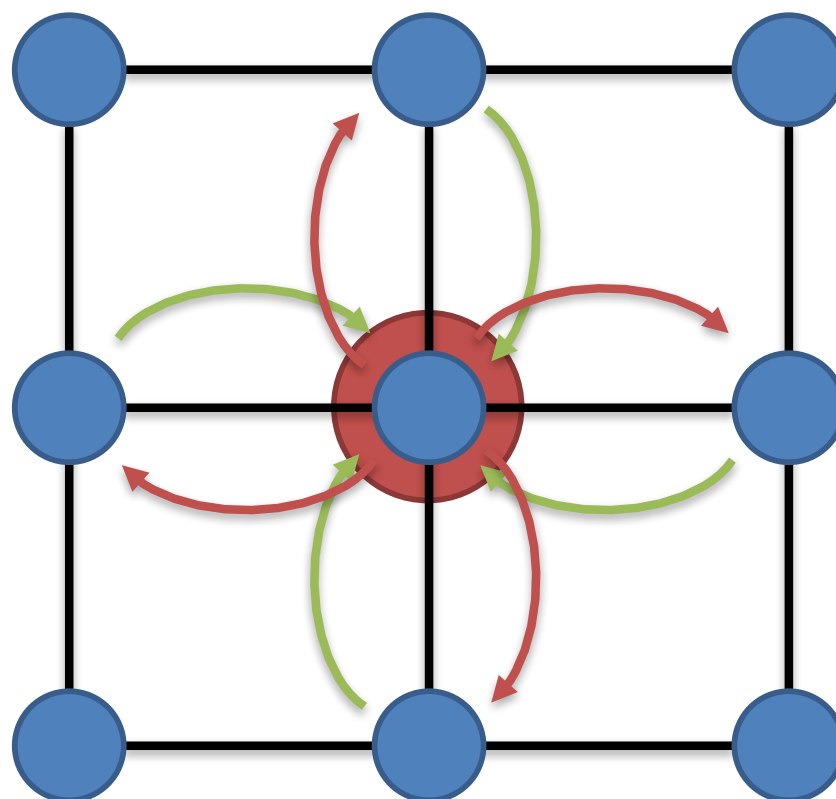
AISTATS'09

UAI'09

Chapter in SUML'10

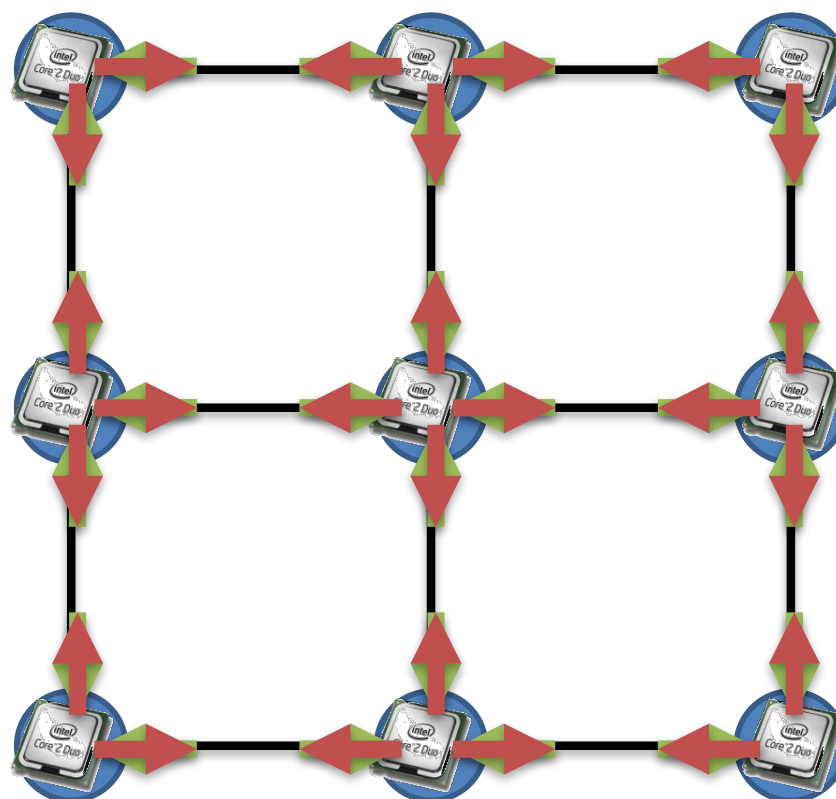
Loopy Belief Propagation (Loopy BP)

- Iteratively estimate the variable beliefs
 - Read **in messages**
 - Updates marginal estimate (**belief**)
 - Send updated **out messages**
- Repeat for all variables until convergence



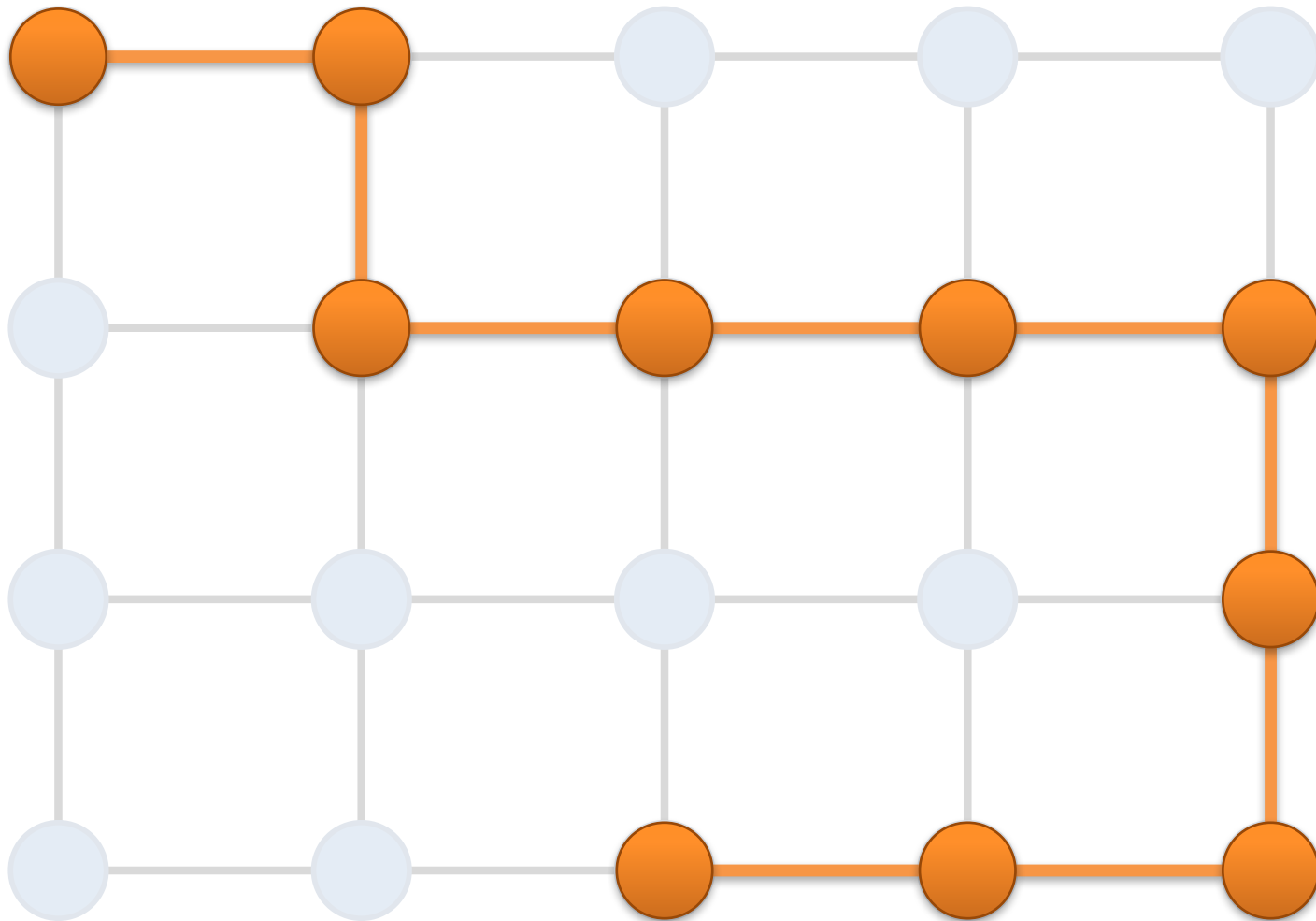
Synchronous Loopy BP

- Often considered embarrassingly parallel
 - Associate processor with each vertex
 - Receive all messages
 - Update all beliefs
 - Send all messages
- Proposed by:
 - Brunton et al. CRV'06
 - Mendiburu et al. GECC'07
 - Kang, et al. LDMTA'10
 - ...

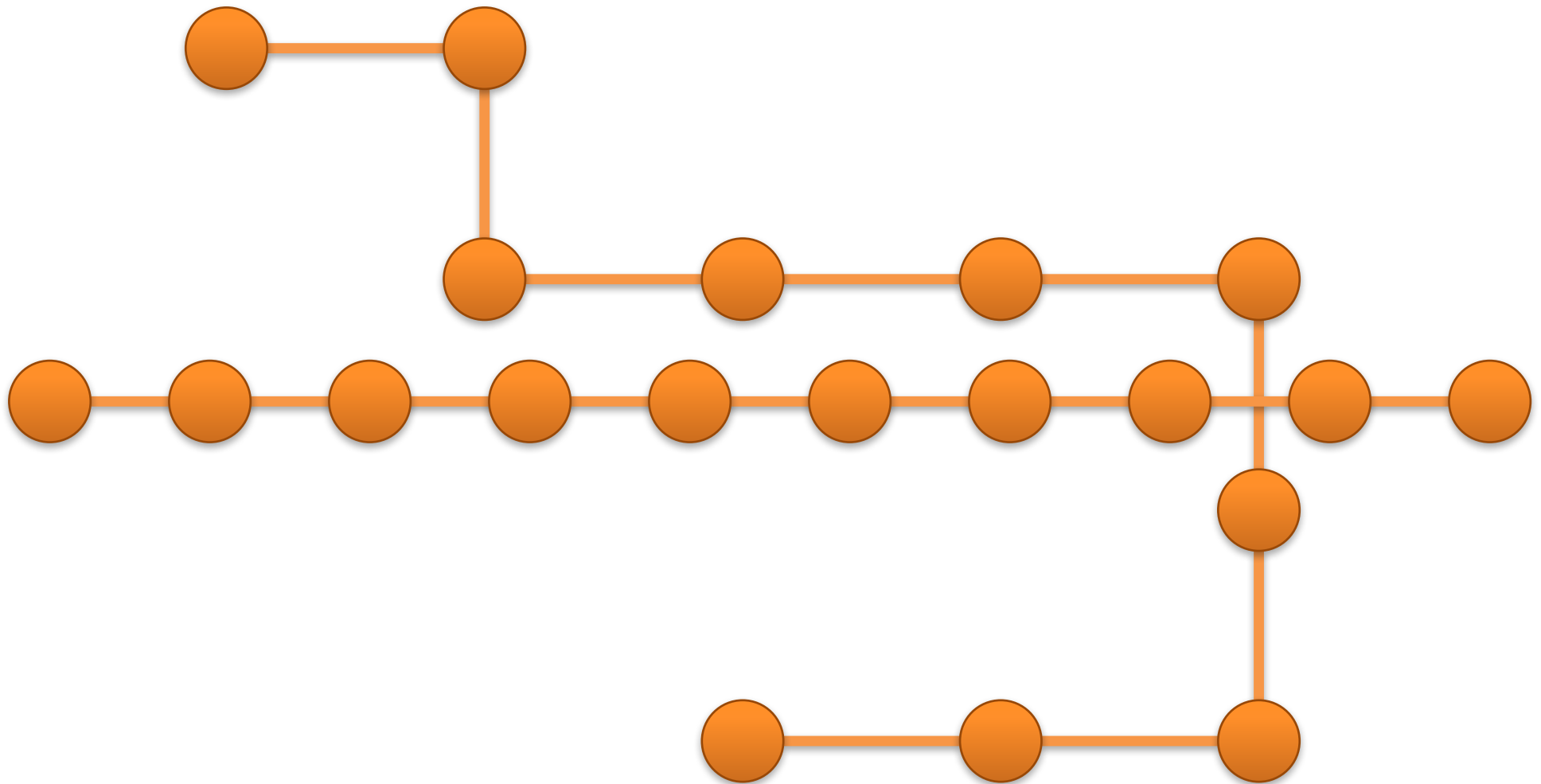


*Is Synchronous Loopy BP
an **efficient** parallel algorithm?*

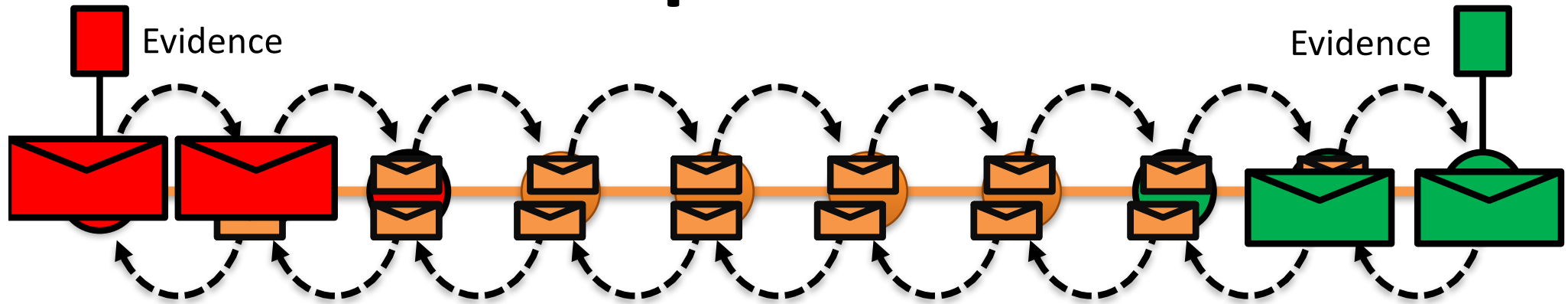
Sequential Computational Structure



Hidden Sequential Structure



Hidden Sequential Structure



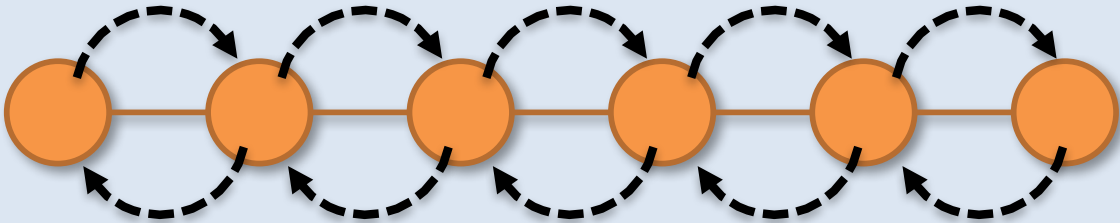
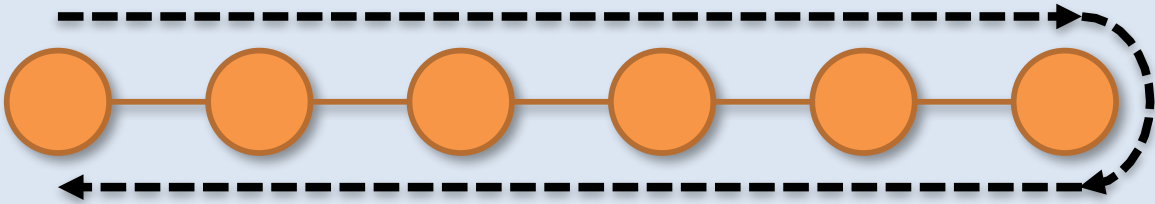
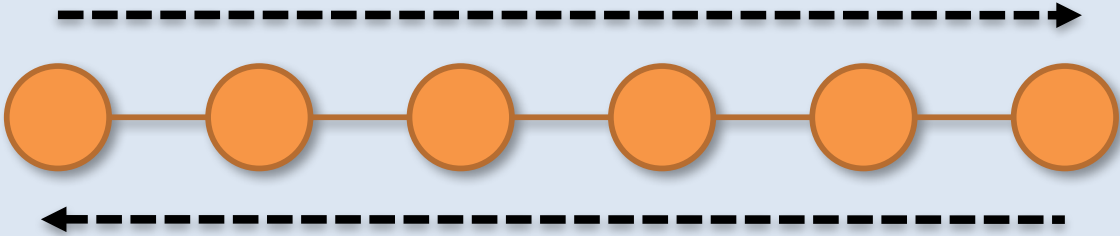
- Running Time:

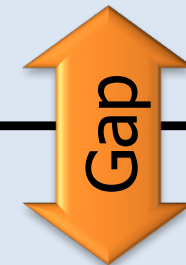
$$\frac{2n \text{ Messages Calculations}}{p \text{ Processors}} \times (n \text{ Iterations to Converge}) = \frac{2n^2}{p}$$

Time for a single
parallel iteration

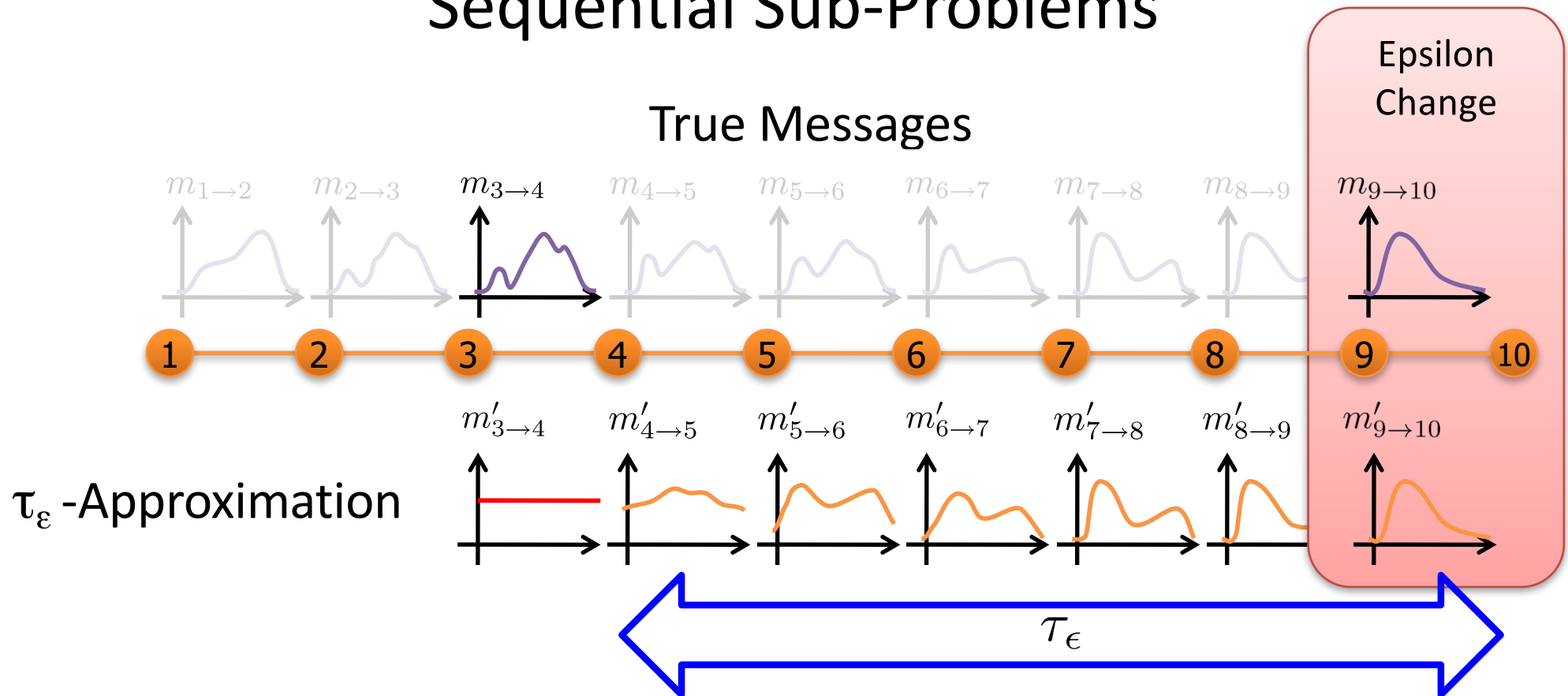
Number of Iterations

Optimal Sequential Algorithm

		Running Time
Naturally Parallel 		$2n^2/p$ $p \leq 2n$
Sequential (Fwd-Bkwd) 		$2n$ $p = 1$
Optimal Parallel 		n $p = 2$

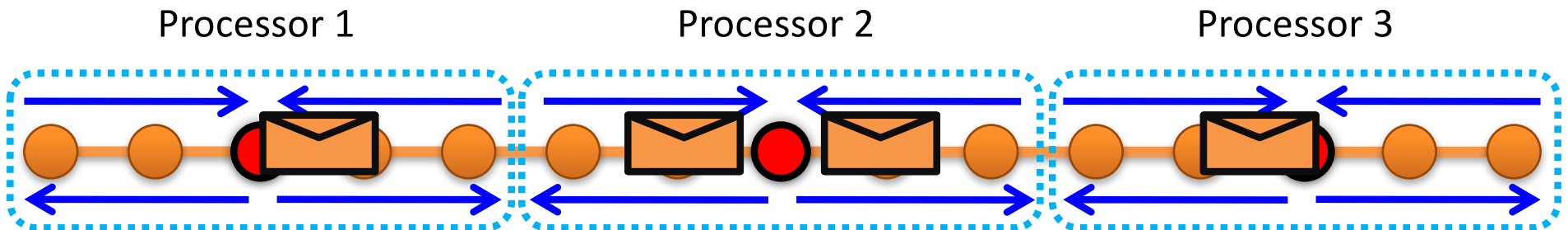


Role of model **Parameters** on Sequential Sub-Problems



- τ_ϵ represents the minimal sequential sub-problem
- Captures dependence on **model parameters**

Optimal Parallel Scheduling



Theorem:

Using p processors this algorithm achieves a τ_ϵ approximation in time:

$$O\left(\frac{n}{p} + \tau_\epsilon\right)$$

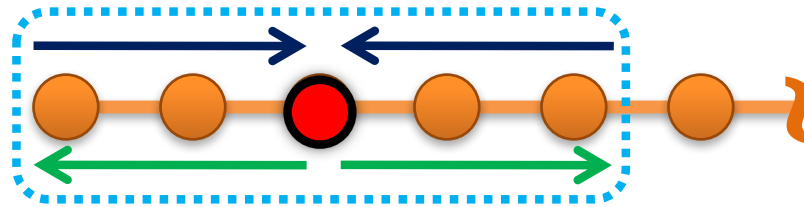
Parallel Component

Sequential Component

and is **optimal** for chain graphical models.

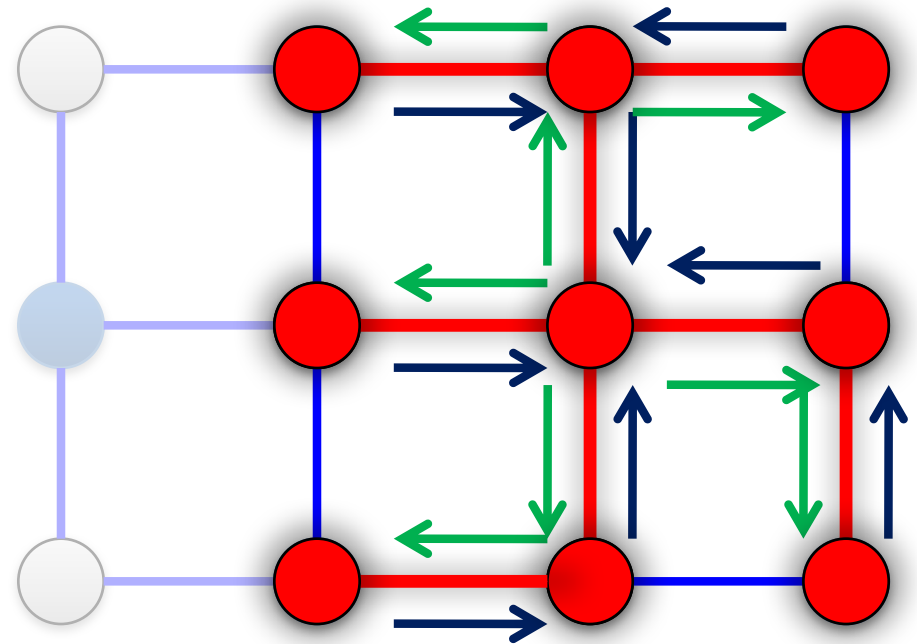
The Splash Operation

- Generalize the optimal chain algorithm:



to arbitrary cyclic graphs:

- 1) Grow a BFS Spanning tree with fixed size
- 2) Forward Pass computing all messages at each vertex
- 3) Backward Pass computing all messages at each vertex



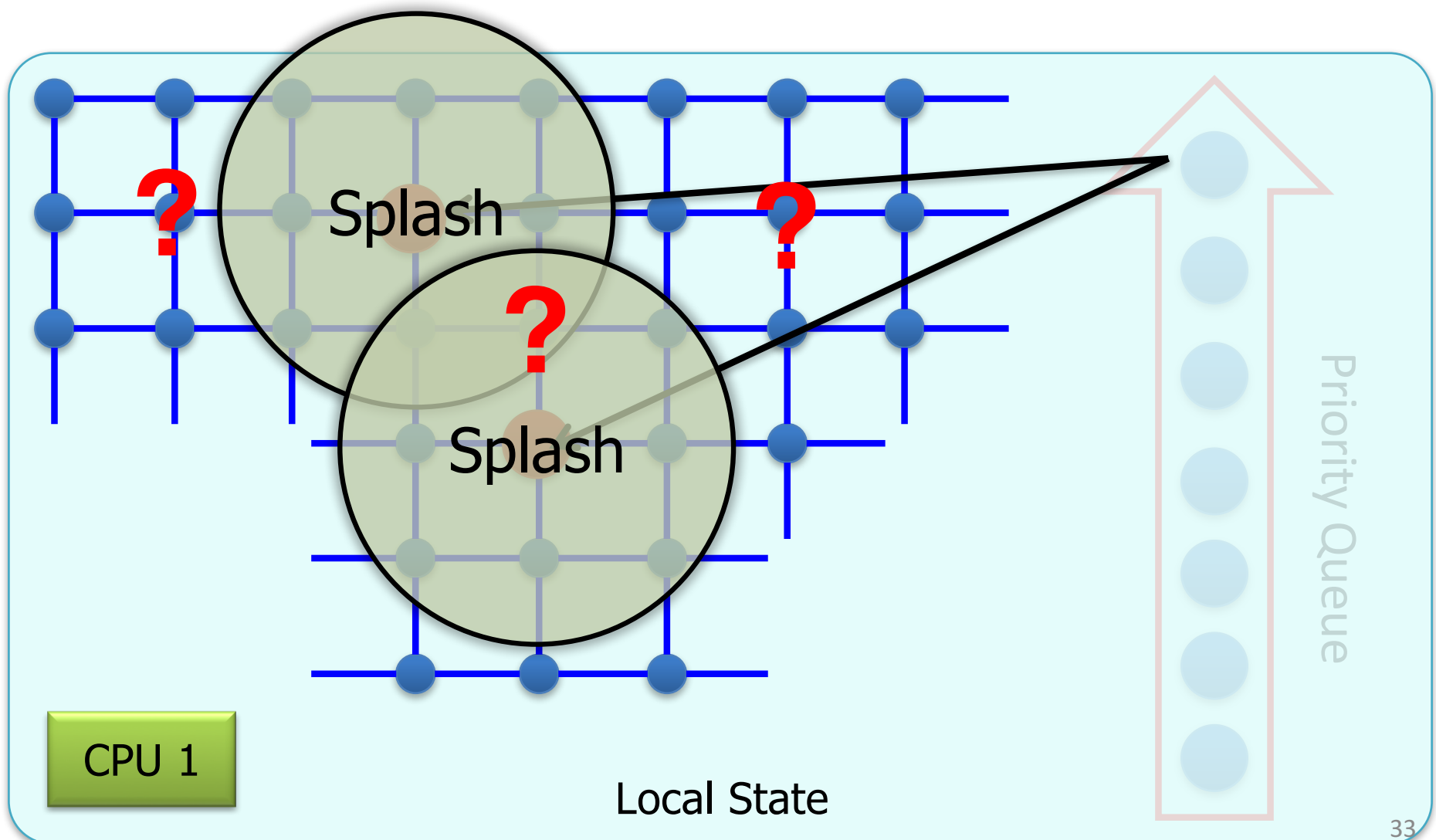
Running Parallel Splashes



- Partition the graph
- Schedule Splashes locally
- Transmit the messages along the boundary of the partition

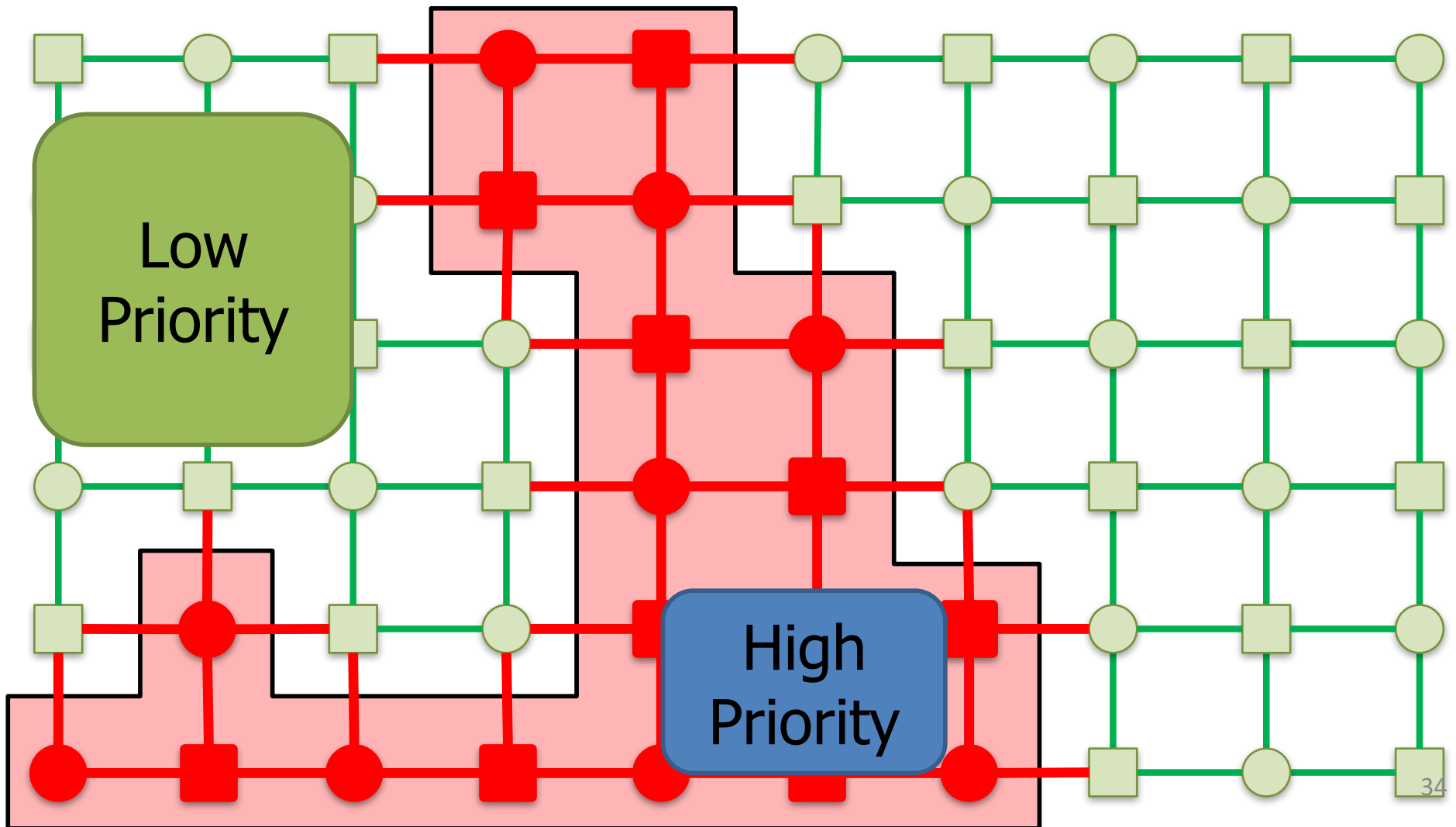
Priorities Determine the **Roots**

- Use a residual priority queue to select roots:

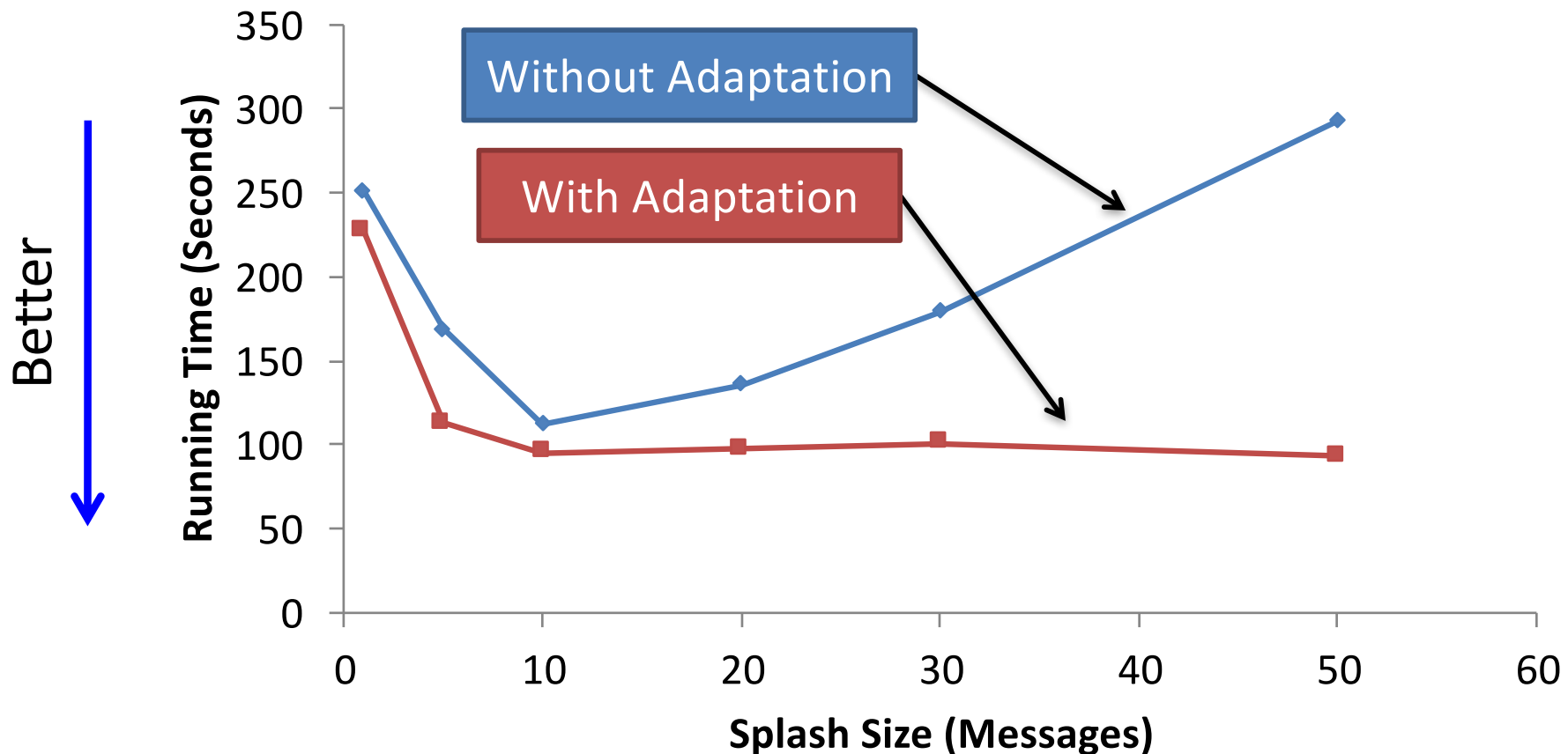


Dynamic Splashes

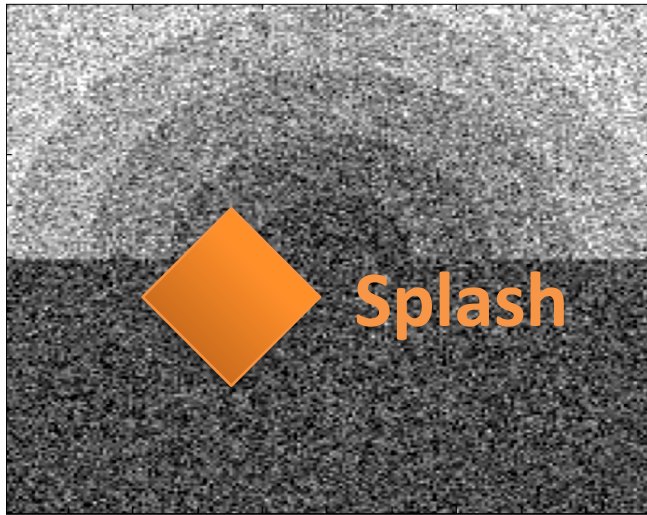
Priorities **adaptively** focus computation by determining the **shape** and **size** of each Splash



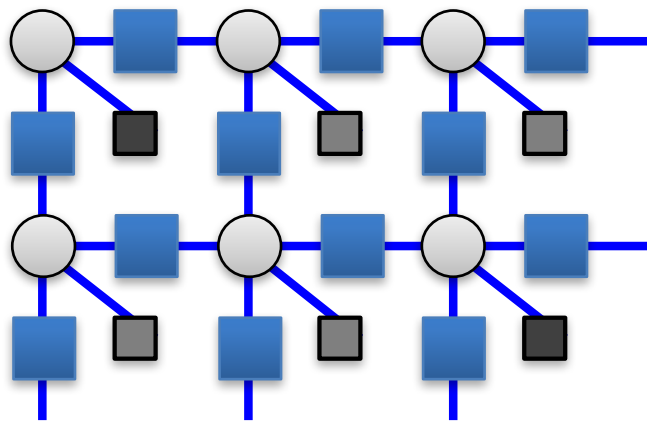
Dynamic Splashes automatically identify the **optimal** splash size



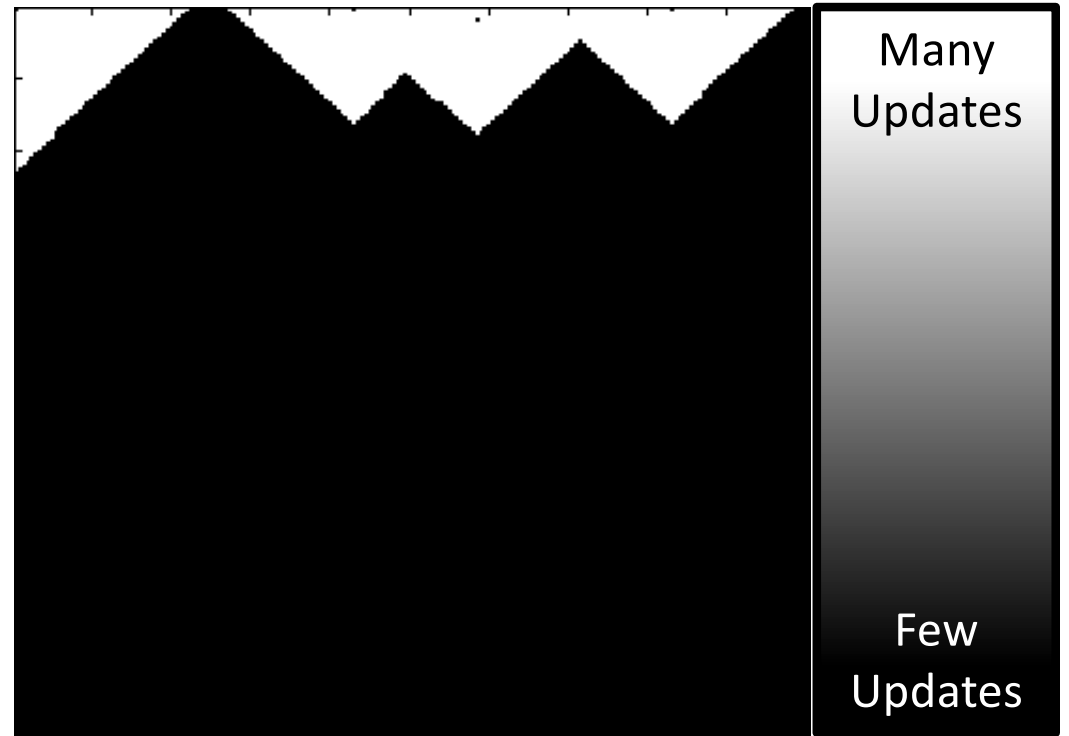
Splash Belief Propagation



Synthetic Noisy Image



Factor Graph



Vertex Updates

Algorithm identifies and focuses on hidden sequential structure

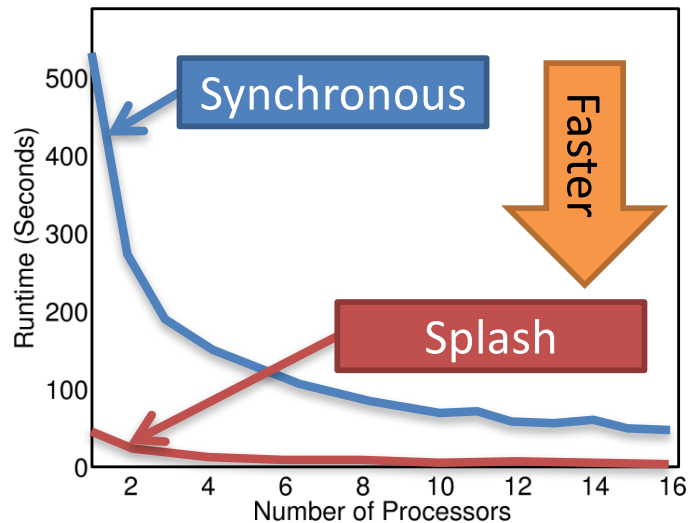
Evaluation

- System Design
 - Multicore and distributed implementations
 - Development was **time consuming**
- Evaluated on several real-world problems
 - Protein interaction and structure prediction
 - Markov Logic Networks
- Compared against several other variants
 - Faster, more efficient, more stable

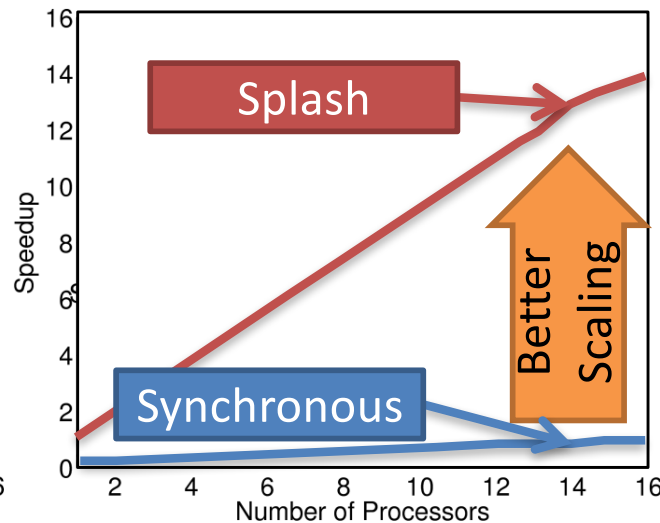
Representative Results

Protein Interaction Models: 14K Vertices, 21K Factors

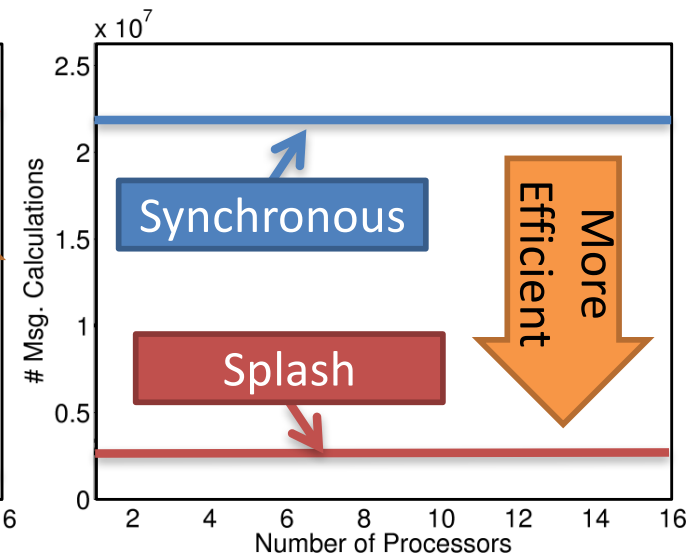
Runtime



Speedup



Total Work



- SplashBP converges more often
- Achieves better prediction accuracy

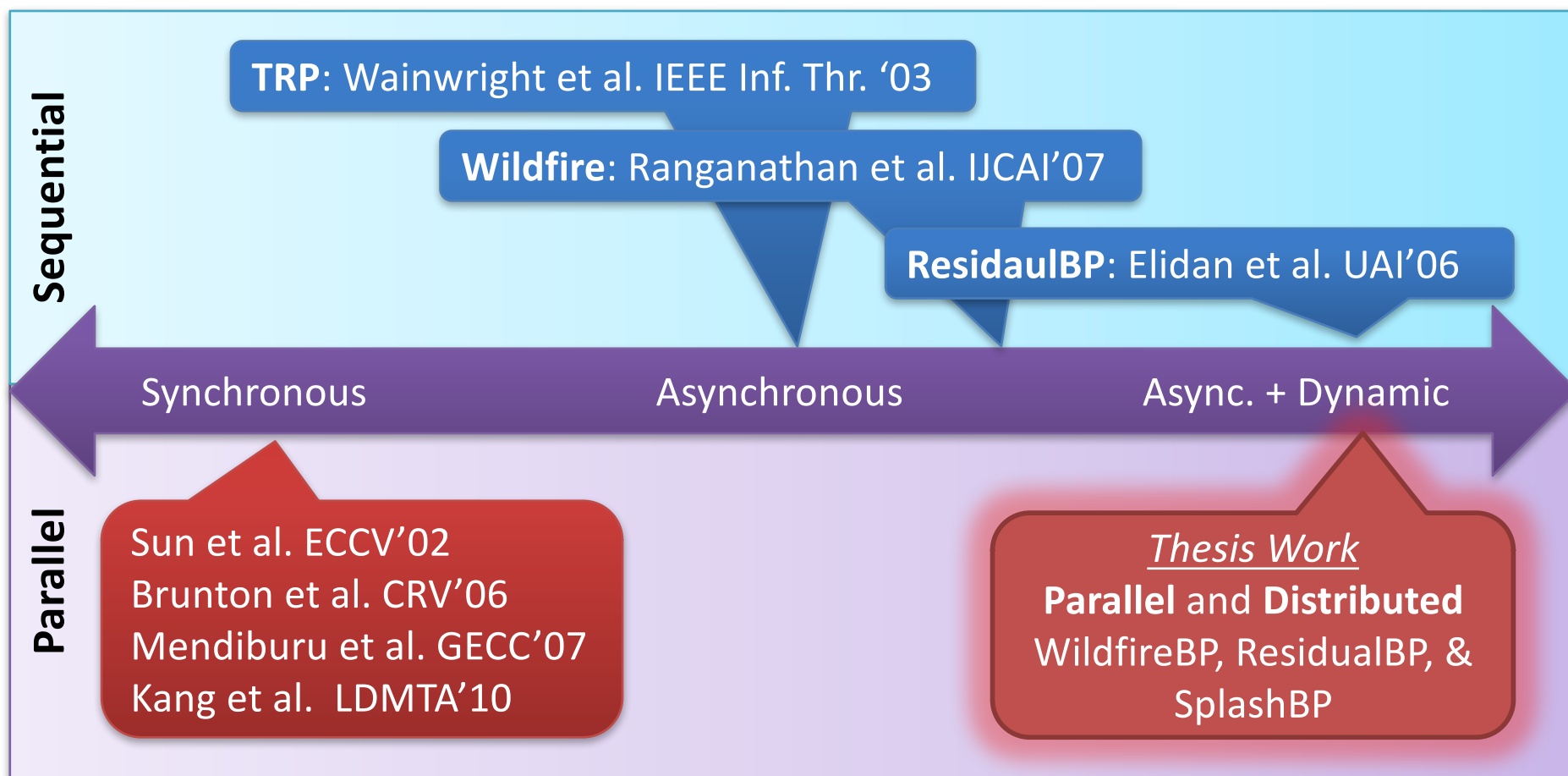
Summary: Belief Propagation

- **Asynchronous + Dynamic** → more efficient
 - *Theoretically and experimentally*
 - *Insight*: parallelize optimal **sequential** algorithm
 - *Tradeoff*: **Parallelism & Convergence**
- *Approximation* → Increased Parallelism
 - Exploit *weak* interactions (τ_ε – approximation)
- Key Contributions:
 - Demonstrate the importance of dynamic asynchronous scheduling in parallel inference
 - Theoretical analysis of work efficiency and relationship to model structure and parameters

GrAD Methodology

- **Graphical**
 - BP updates only depend on adjacent vertices
- **Asynchronous**
 - Compute messages sequentially within Splash
- **Dynamic**
 - Priority scheduling and adaptive Splashes

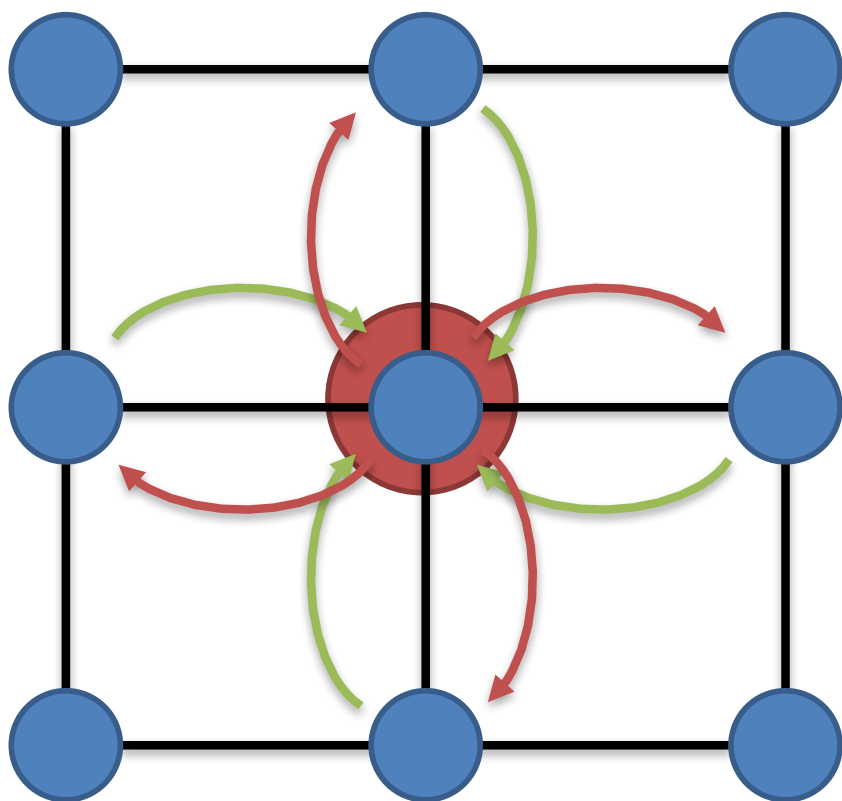
Additional Related Work



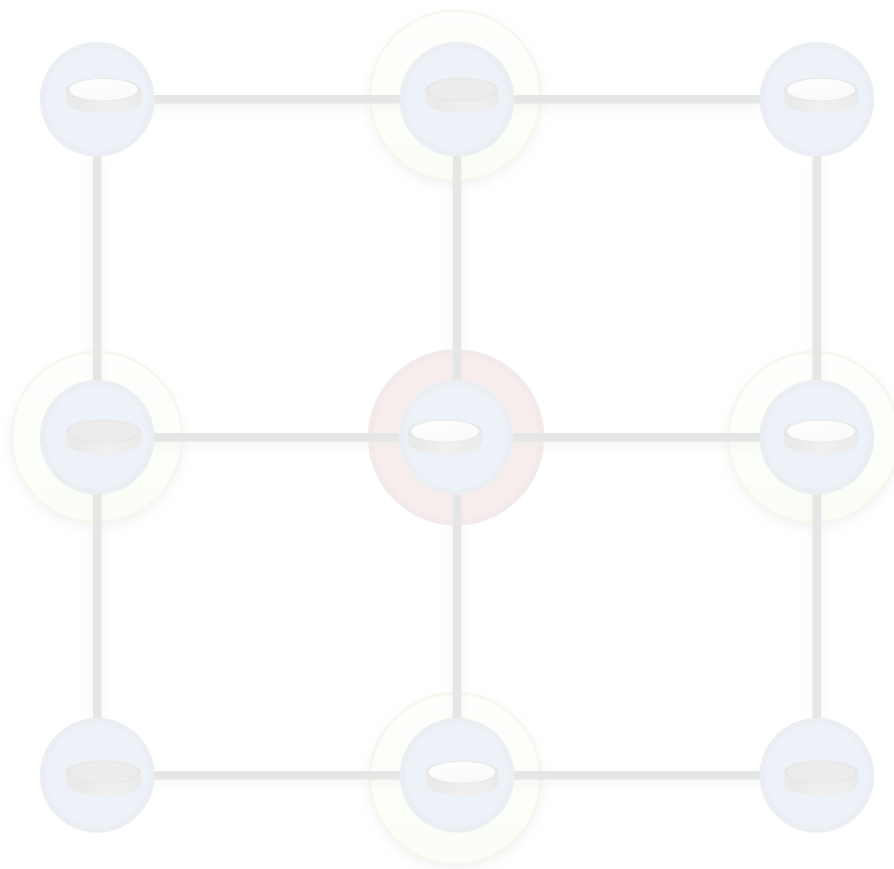
- **Parallel Exact Inference:** Pennock et al. UAI'98
- **Approximate Messages:** Ihler et al. JMLR'05

Parallel and Distributed Algorithms for Probabilistic Inference

Belief Propagation



Gibbs Sampling



Parallel Gibbs Sampling

*An **asynchronous** Gibbs Sampler that
dynamically addresses **strong dependencies**.*

Joint Work With

Yucheng Low

Arthur Gretton

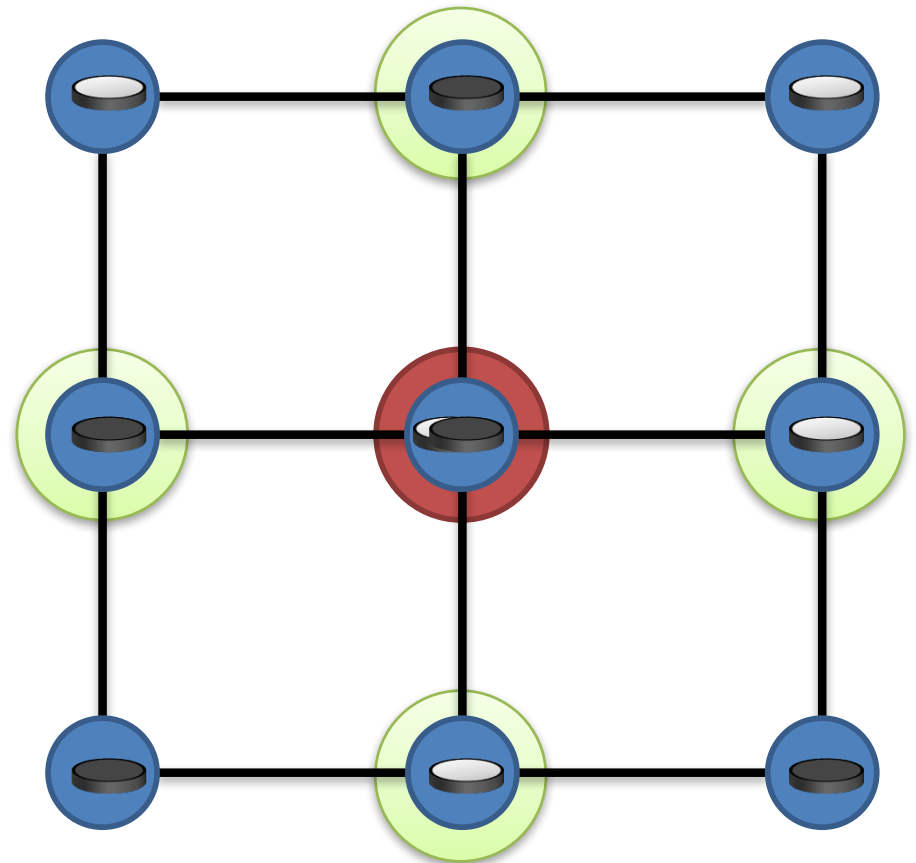
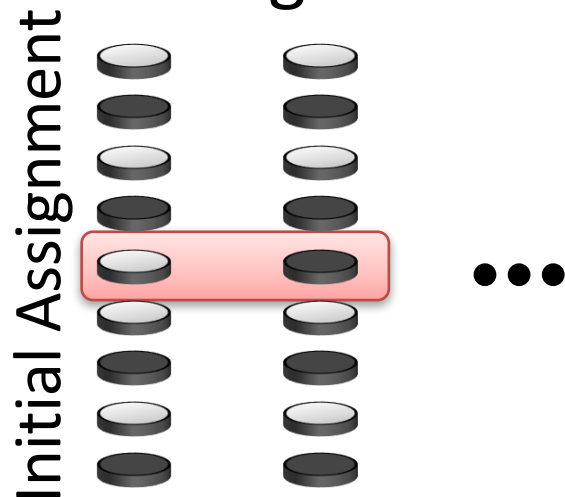
Carlos Guestrin

Published

AISTATS'11 (Related to work in WSDM'12)

Gibbs Sampling [Geman & Geman, 1984]

- **Sequentially** for each variable in the model
 - Select **variable**
 - Use **adjacent assignments** to construct a biased coin
 - Flip coin and update assignment to **variable**

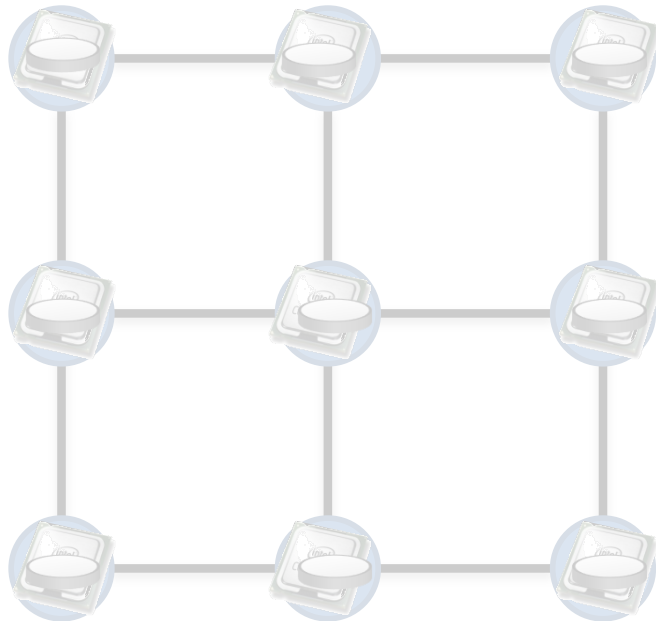


*Can we sample multiple
variables in **parallel**?*

From the original paper on Gibbs Sampling:

*“...the MRF can be divided into collections of [variables] with each collection assigned to an **independently** running **asynchronous processor**.”*

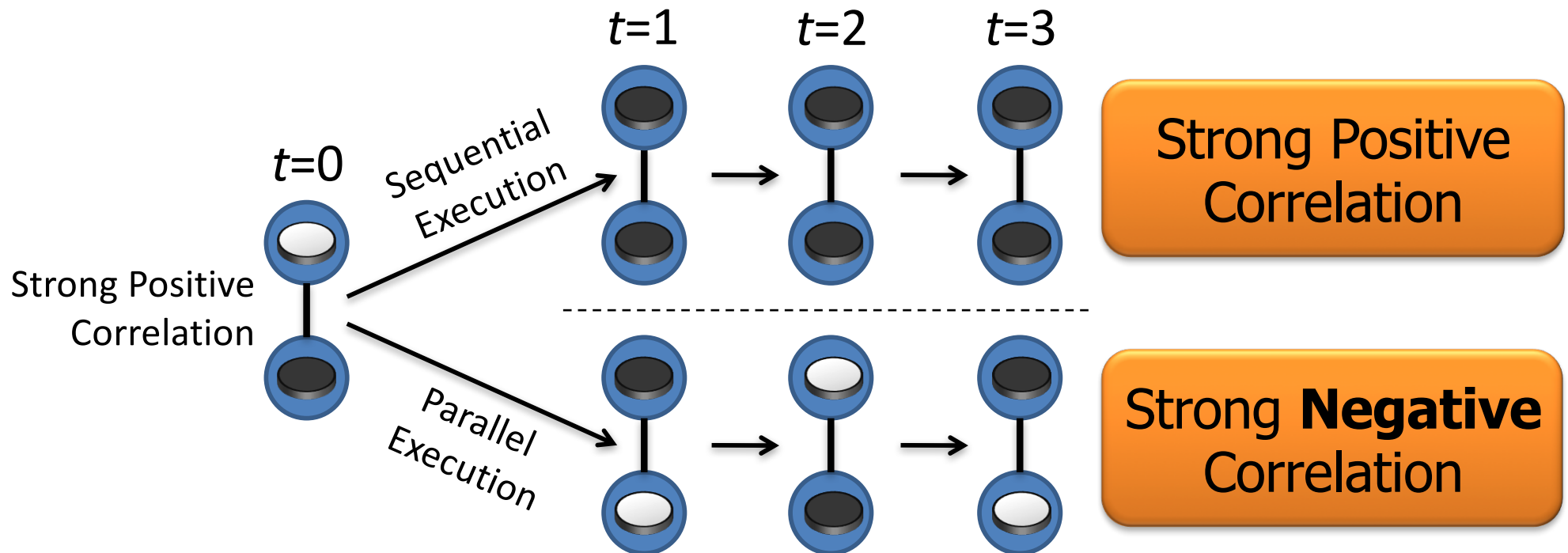
-- Stuart and Donald Geman, 1984.



Embarrassingly
Parallel!

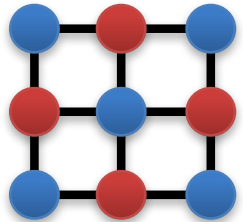
Converges to the
wrong distribution!

The problem with Synchronous Gibbs sampling

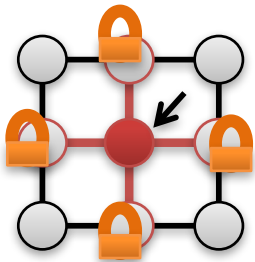


- *Adjacent variables **cannot** be sampled simultaneously.*

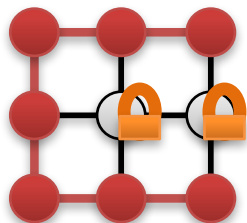
Introduced Three Convergent Samplers



Chromatic: Use graph coloring to synchronously sample independent sets



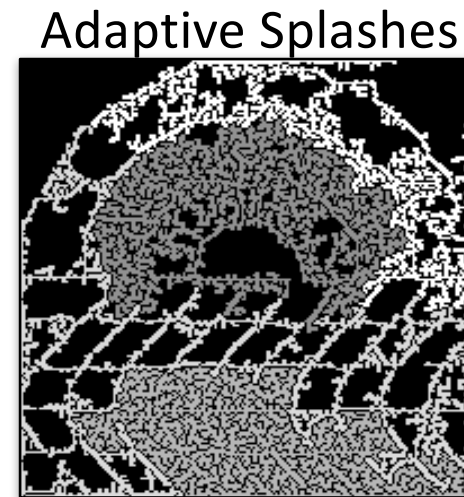
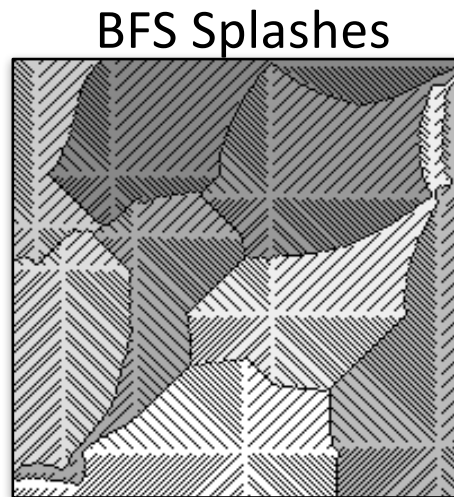
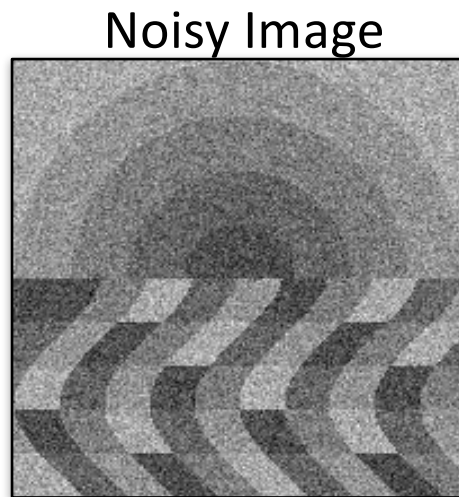
Asynchronous: Markov Blanket Locks ensure serializable execution



Splash: Adaptively constructs thin junction tree blocks

Dynamically Prioritized Sampling

- Prioritize Gibbs updates
- Adapt the **shape** of the Splash to span strongly coupled variables:



Theorem: *Chromatic* Sampler

- **Ergodic:** converges to the correct distribution
 - Based on graph coloring of the Markov Random Field
- **Quantifiable** acceleration in **mixing**

Time to update
all variables once

$$O \left(\frac{n}{p} + k \right)$$

The diagram shows three labels on the right with blue arrows pointing to parts of the formula $O \left(\frac{n}{p} + k \right)$:

- An arrow from "# Variables" points to n .
- An arrow from "# Colors" points to k .
- An arrow from "# Processors" points to p .

Theorem

Asynchronous and *Splash Gibbs* Sampler

- **Ergodic:** converges to the correct distribution
 - Requires vanishing adaptation
 - Corrected an error in a result by Levin & Casella *J. Multivar. Anal.* '06
- **Expected Parallelism:**

$\mathbf{E}(\# \text{active processors})$

$$\geq 1 + (p - 1) \left(1 - (p - 1) \left(\frac{d + 1}{n} \right) \right)$$

Processors

Variables

Max Degree

Evaluation

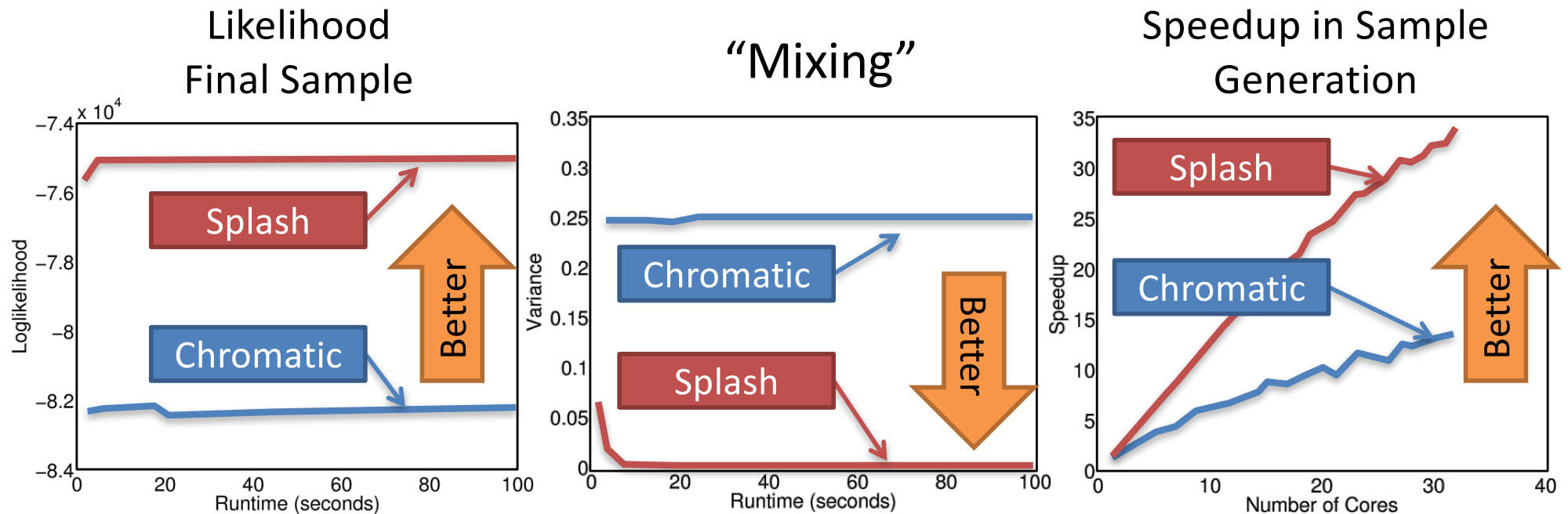
- Implemented multicore version:
 - Built using a GraphLab prototype
 - Substantially shorter development time
 - Novel **junction tree** construction algorithm
 - Markov blanket locking protocol
- Evaluated on large real-world problems

Experimental Results

- Markov logic network with **strong dependencies**

10K Variables

28K Factors



- The *Splash* sampler outperforms the *Chromatic* sampler on models with **strong dependencies**

Contributions: Gibbs Sampling

- Proposed **three** *convergent* Gibbs samplers
 - Chromatic, Asynchronous, Splash
 - Spectrum partially synchronous to asynchronous
 - New algorithms for junction tree construction
- Theoretical analysis of parallel Gibbs sampling
 - Convergence of asynchronous blocking
 - Relate parallelism to model structure
 - Stationary distribution of synchronous sampler
- Experimental analysis on real-world problems and systems

GrAD Methodology

- **Graphical**
 - Gibbs updates depend only on neighbors in MRF
- **Asynchronous**
 - Graph *Coloring* and *Markov Blanket Locks*
- **Dynamic**
 - Prioritized updates and adaptive Splash

Related Work

Ergodic (Convergent)

- Geman & Geman. Pami '84
- **Trees:** Hamze et al. UAI'04
- **Dynamic Blocking:** Barbu et al. IEEE Trans Pattern Analysis '05

Thesis

Chromatic,
Asynchronous,
and
Splash Gibbs

Parallel & Distributed

LDA & Bayesian Networks

- Newman et al. NIPS'07
- Asuncion et al. NIPS'08
- Yan et al. NIPS'09

Amr et al. WSDM'12

- Asynchronous approximations empirically perform well



facebook

flickr

You Tube

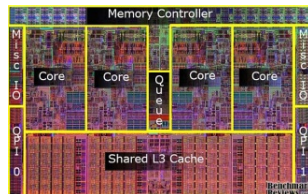
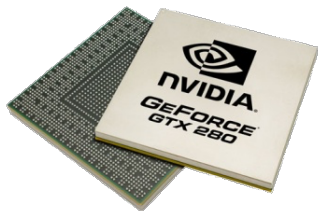
Massive Structured Problems

Probabilistic Graphical Models

Parallel and **Distributed** Algorithms
for Probabilistic **Inference**

GraphLab & PowerGraph

Advances Parallel Hardware



Parallel Algorithms for Probabilistic Inference

GraphLab & PowerGraph

Parallel Hardware

Joint Work With

Yucheng Low Aapo Kyrola Haijie Gu Danny Bickson
Carlos Guestrin Joe Hellerstein Guy Blelloch David O'Hallaron

Published Results

UAI'10 VLDB'12

How do we **design** and **implement** **GrAD Algorithms**

We could:

- *design* and *implement* for each architecture?
 - **Time consuming**
 - Repeatedly solving the same system problems
- use high-level abstractions like **MapReduce**?
 - Unable to express:

- **Graphical**
- **Asynchronous**
- **Dynamic**

GrAD Methodology

Solution: **GraphLab**

GraphLab is a **Graph-Parallel** Abstraction



Map Reduce

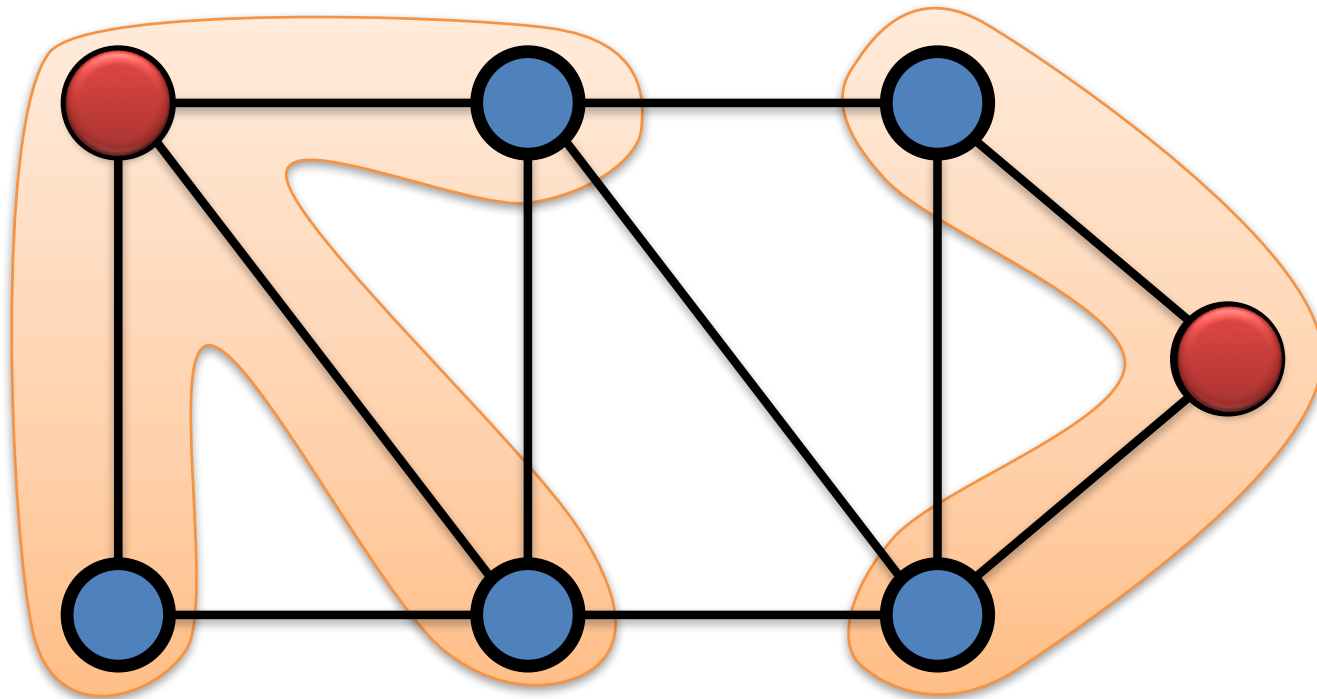
- *Independent Data*
- *Single Pass*
- *Synchronous*

GraphLab

- Graph Structured Data
- Iterative Computation
- Dynamic + Asynchronous

The GraphLab Abstraction

- A user-defined **Vertex Program** runs on each vertex
- **Graph** constrains **interaction** along edges
 - Directly **read** and **modify** the state of adjacent vertices and edges
- **Parallelism**: run multiple vertex programs simultaneously



The GraphLab Vertex Program

Vertex Programs directly **access** adjacent vertices and edges

```
GraphLab_PageRank(i)
```

```
// Compute sum over neighbors
```

```
total = 0
```

```
foreach( j in neighbors(i)):
```

```
    total = total + R[j] * wji
```

```
// Update the PageRank
```

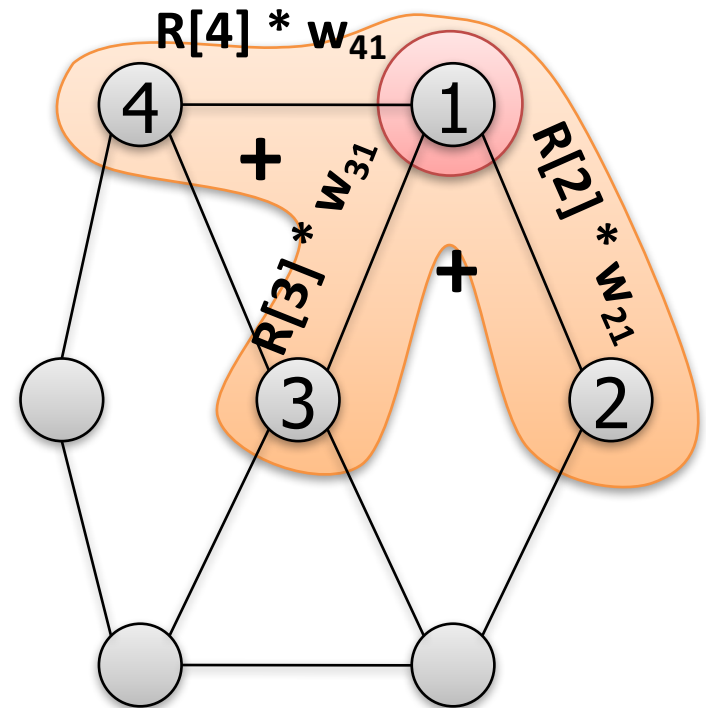
```
R[i] = total
```

```
// Trigger neighbors to run again
```

```
priority = |R[i] - oldR[i]|
```

```
if R[i] not converged then
```

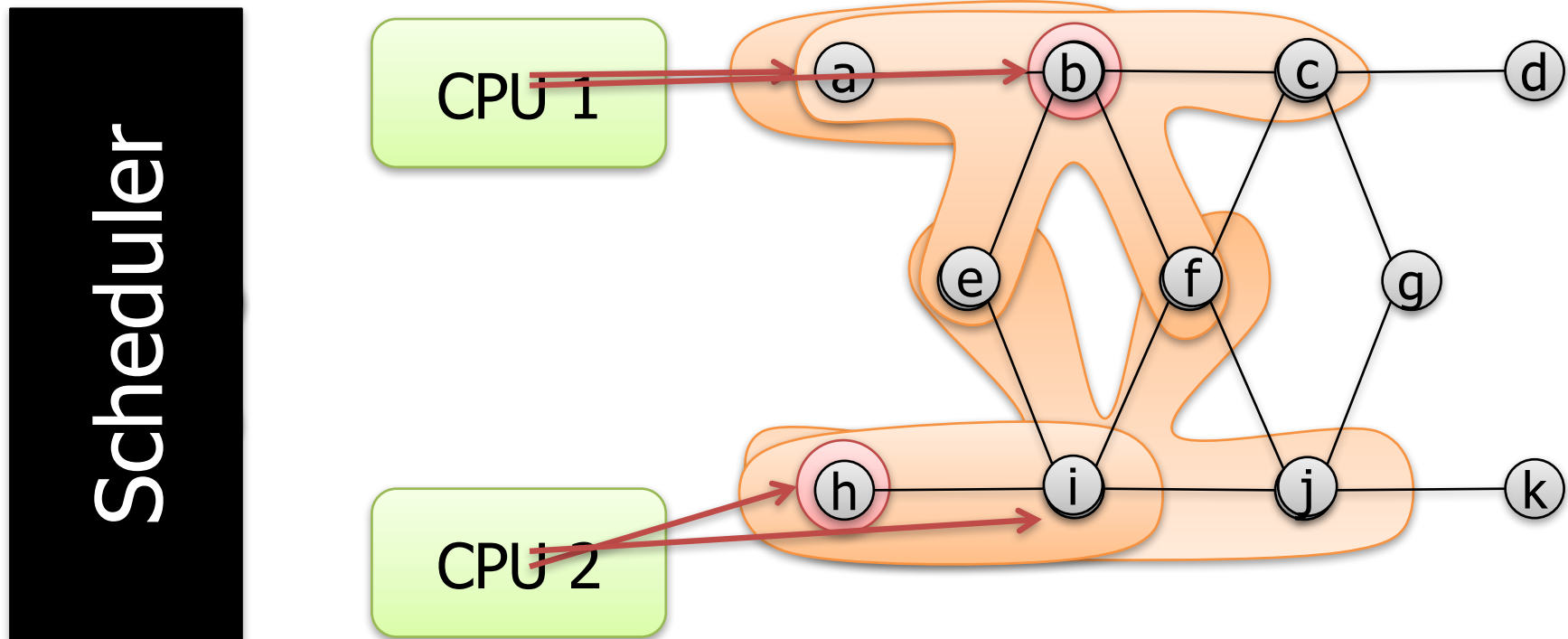
```
    signal neighbors(i) with priority
```



Dynamics

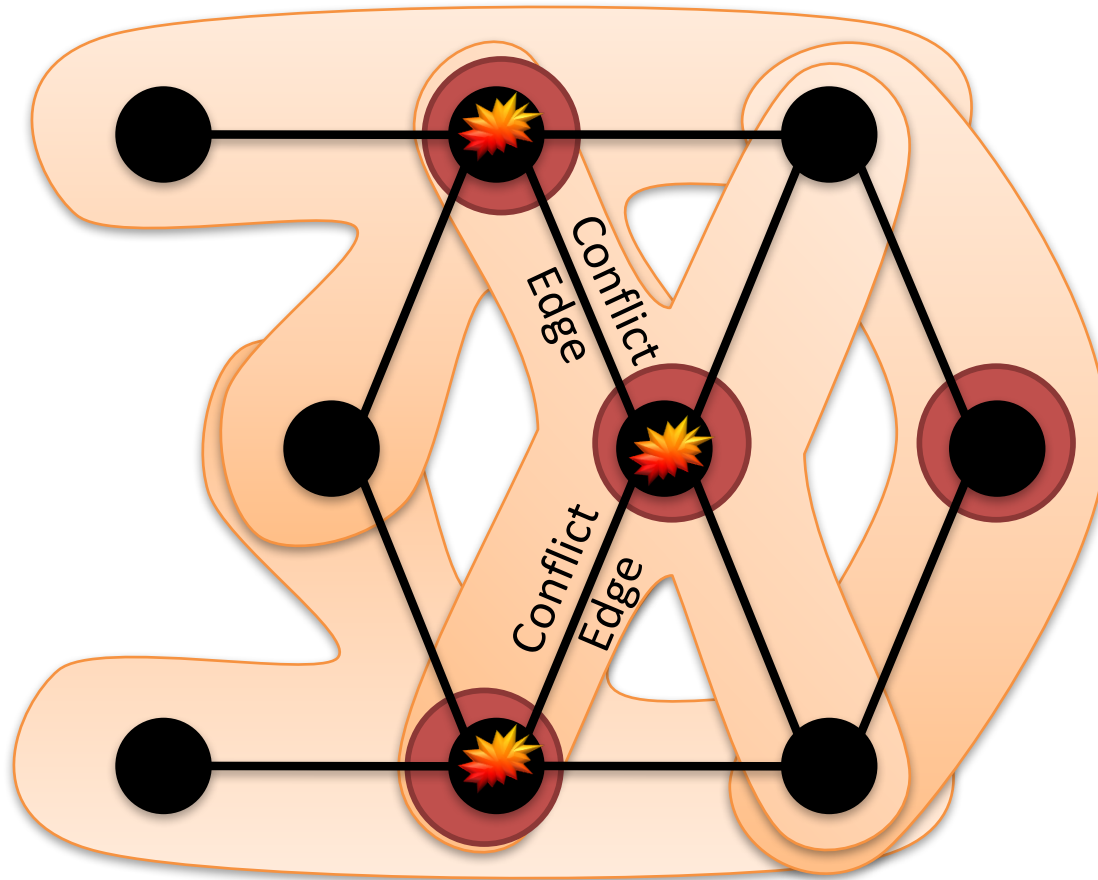
GraphLab is **Asynchronous**

The **scheduler** determines the order that vertices are executed



Scheduler can **prioritize** vertices.

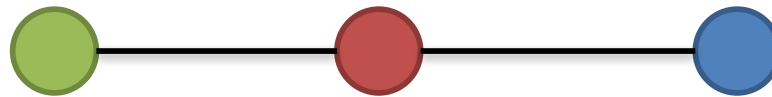
GraphLab is **Serializable**



- Automatically ensures **serializable** executions

Serializable Execution

For **each parallel execution**, there exists a **sequential execution** of vertex-programs which produces the same result.



Parallel

CPU 1

CPU 2

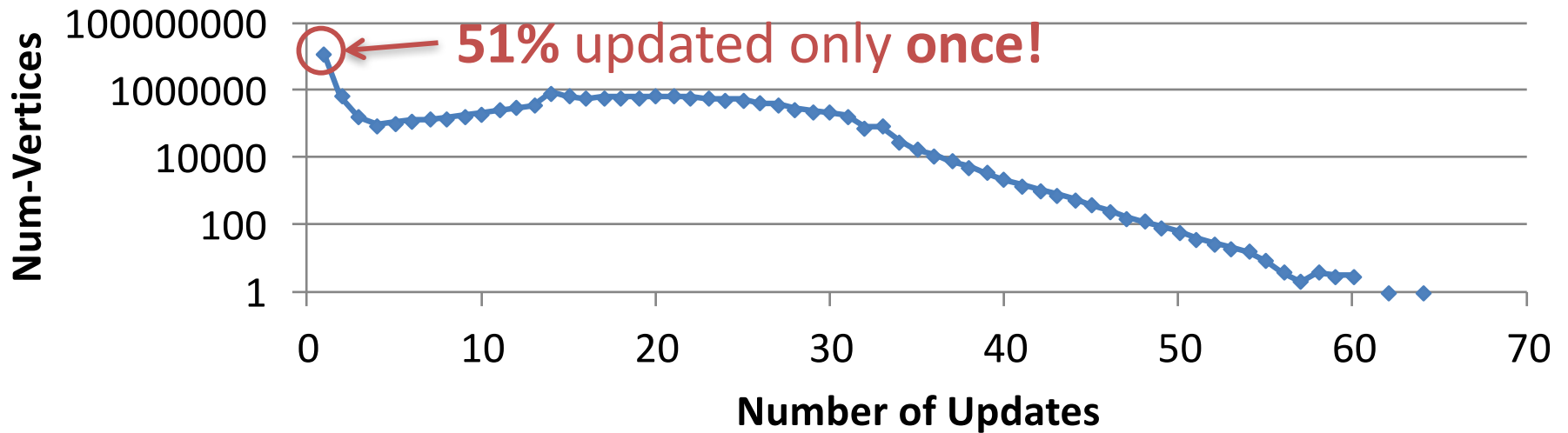
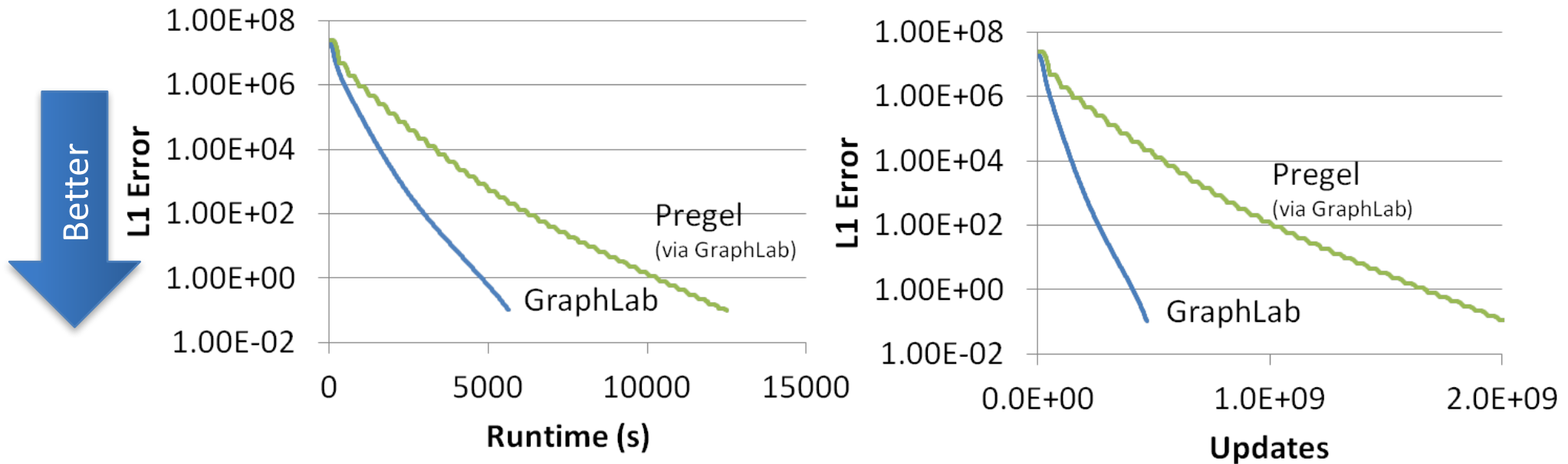
Sequential

Single
CPU

The GraphLab System

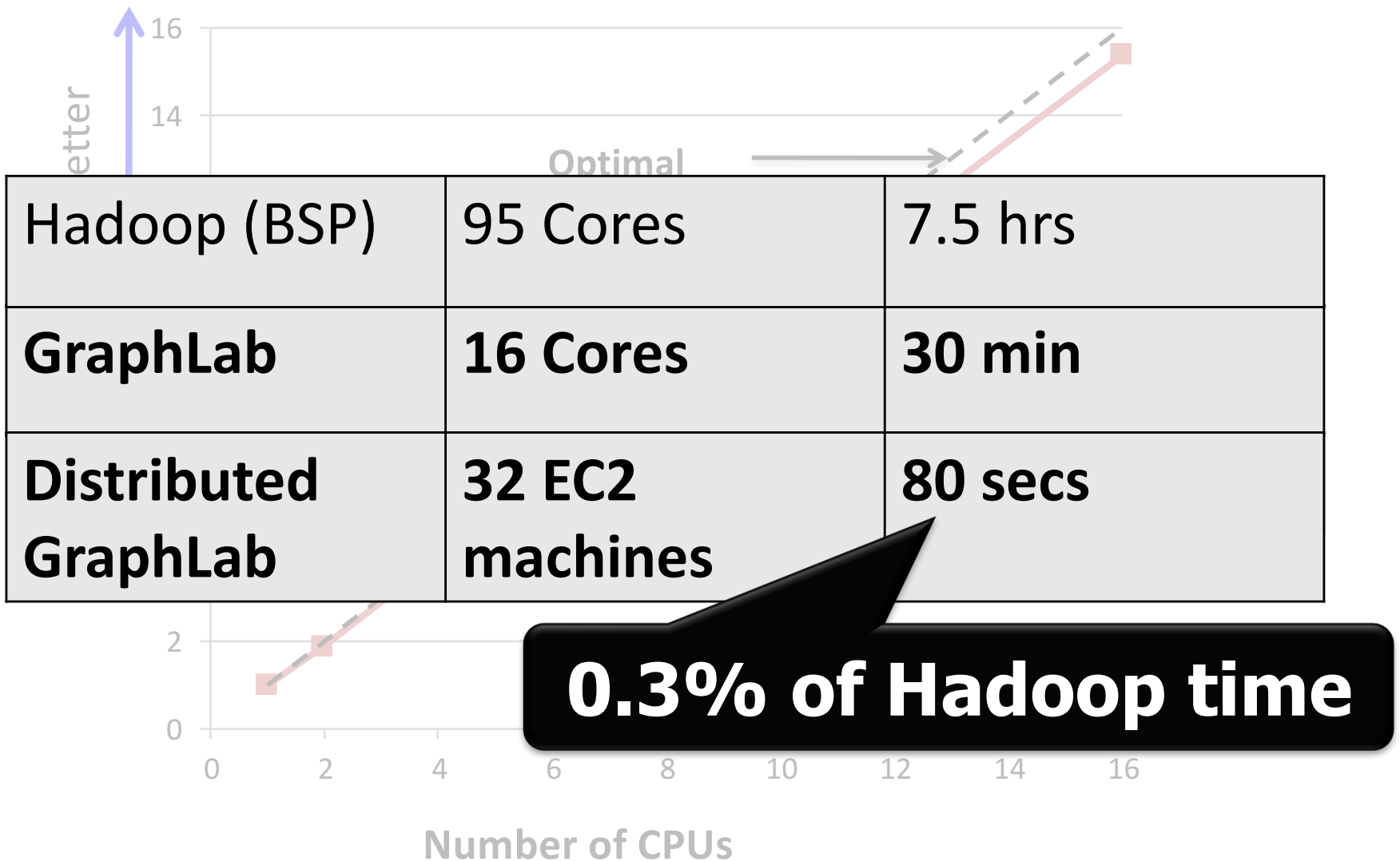
- Implemented as a C++ API
 - Widely downloaded open-source project
- **Multicore** and **distributed** versions:
 - **Hide Latency**: Pipelined locking
 - **Fault Tolerance**: Chandy-Lamport Snapshot
- Tested on a wide range of ML algorithms
 - ALS, BP, Gibbs, Lasso, CoEM, SVM, LDA, ...

GraphLab vs. Pregel (BSP)



PageRank (25M Vertices, 355M Edges)

Never Ending Learner Project (CoEM)



Summary: GraphLab

- **Generalizes** the GrAD Methodology
 - ALS, BP, Gibbs, Lasso, CoEM, SVM, PageRank, LDA, ...
- **Simplifies** the *design* and *implementation* of GrAD Algorithms
- Substantially outperforms existing systems
- Key Contributions:
 - Formalized the graph-parallel setting
 - Isolates computation from movement of data
 - Strong serializability guarantees
 - Evaluation on a wide range of algorithms

Thus far...

GraphLab provided exciting
scaling performance

But...

**We couldn't scale up to
Altavista Webgraph from 2002
1.4B vertices, 6.6B edges**

Parallel Algorithms for Probabilistic Inference

GraphLab & PowerGraph

Parallel Hardware

Joint Work With

Yucheng Low Aapo Kyrola Haijie Gu Danny Bickson
Carlos Guestrin Joe Hellerstein Guy Blelloch David O'Hallaron

Published Results

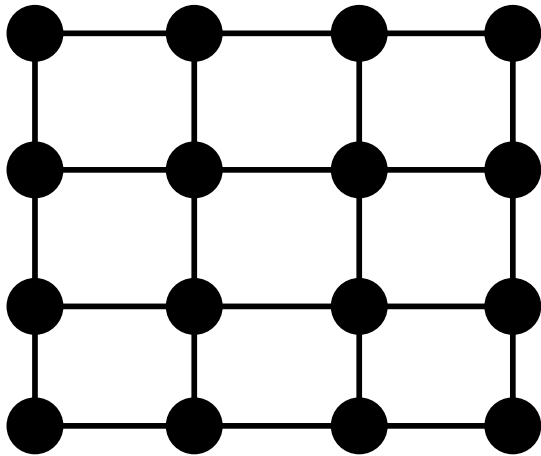
OSDI'12

Natural Graphs

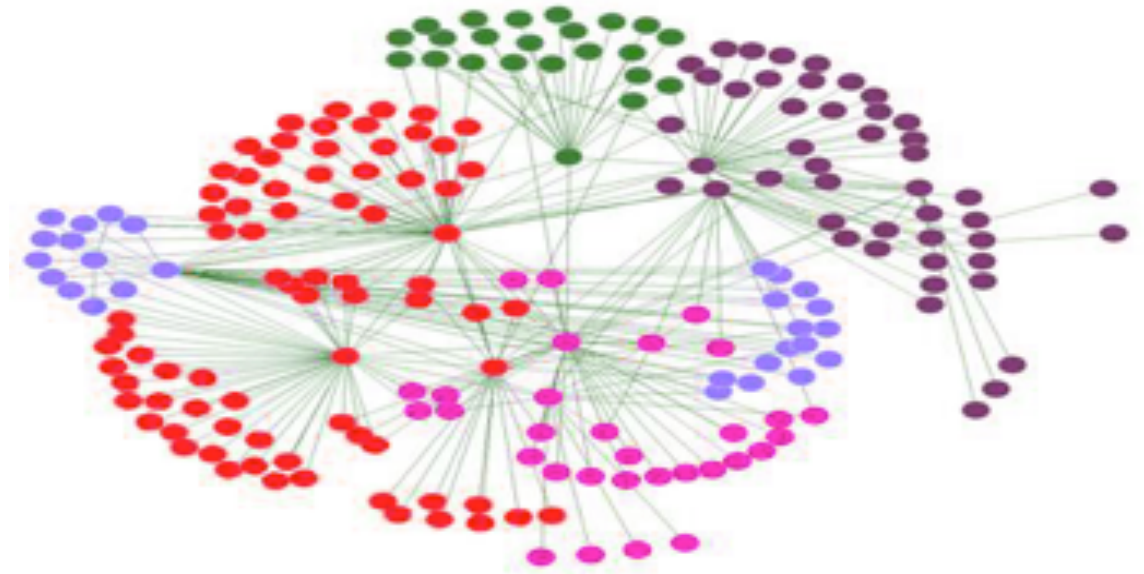
Graphs derived from natural phenomena



Properties of Natural Graphs



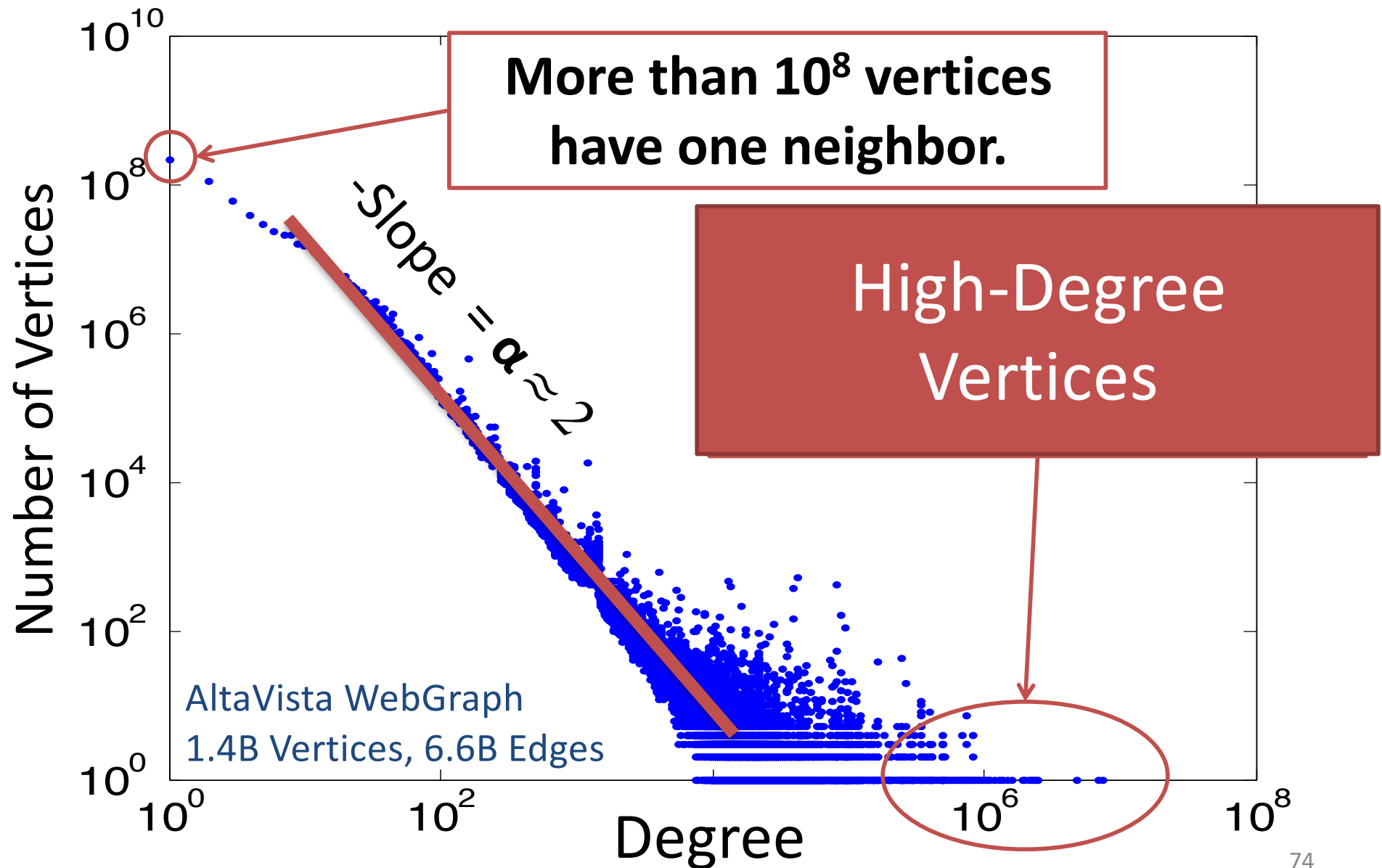
Regular Mesh



Natural Graph

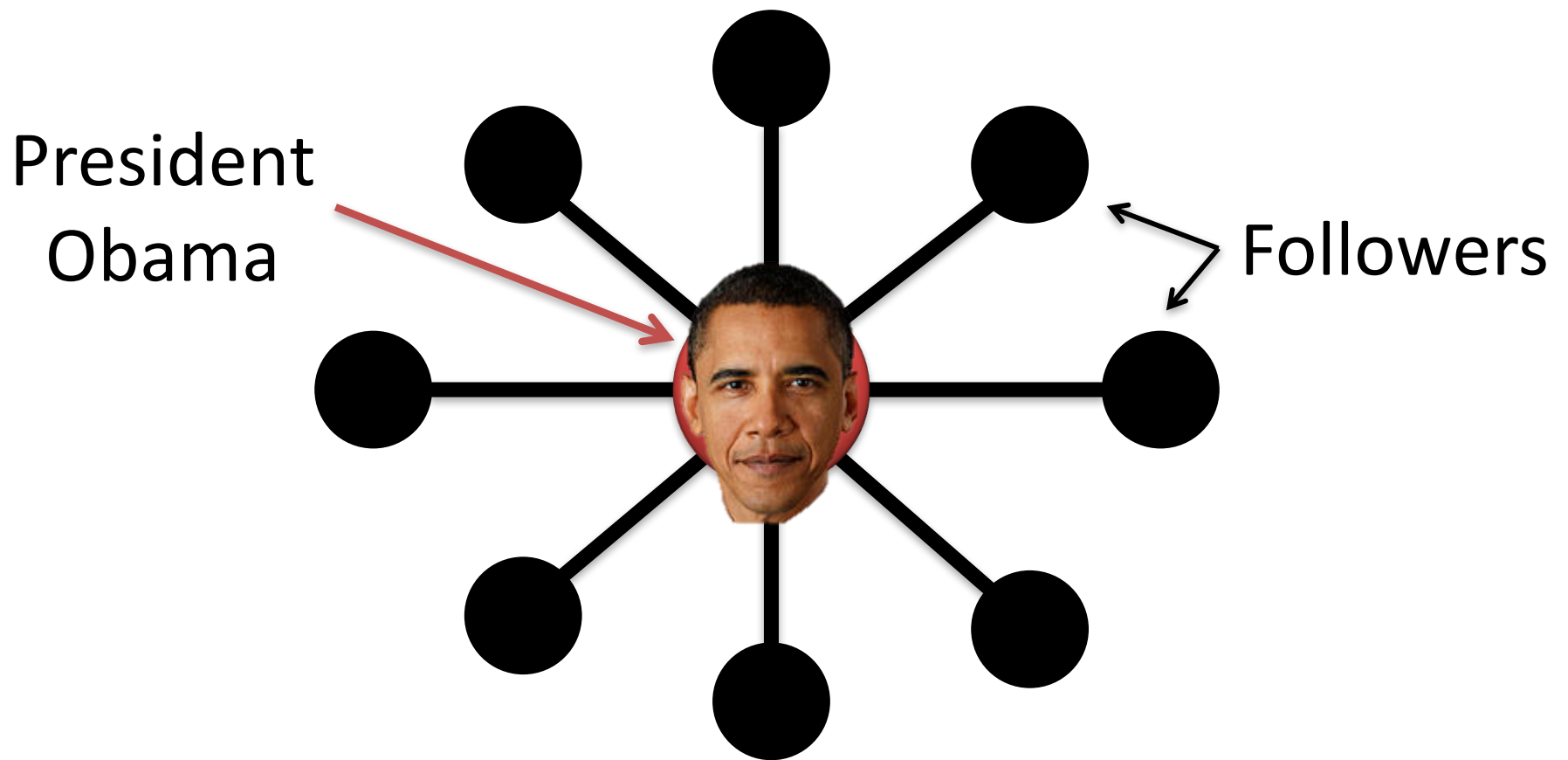
Power-Law Degree Distribution

Power-Law Degree Distribution

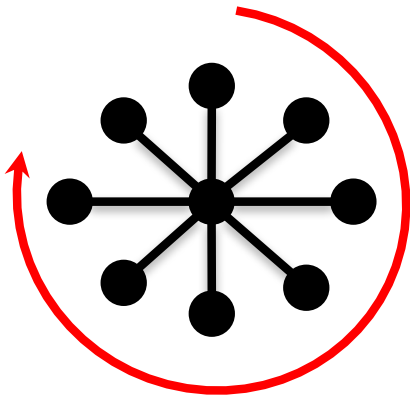


Power-Law Degree Distribution

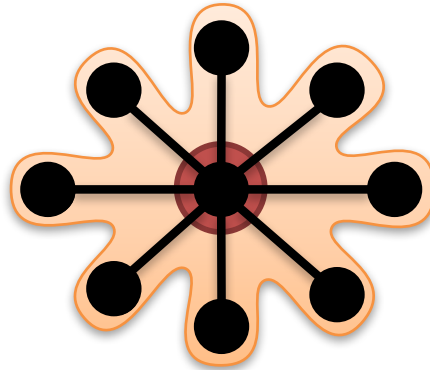
“Star Like” Motif



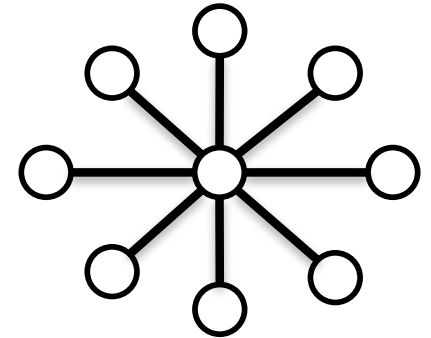
Challenges of High-Degree Vertices



Sequentially process
edges



Touches a large
fraction of graph



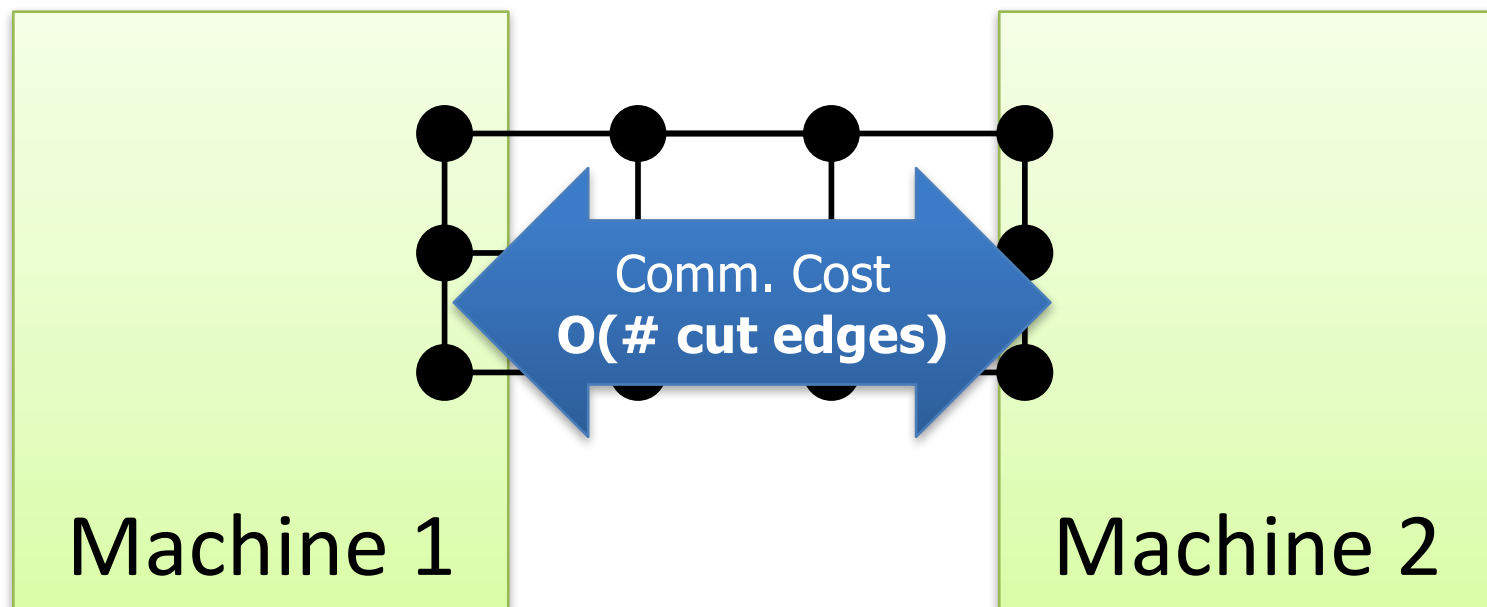
Edge meta-data
too large for single
machine



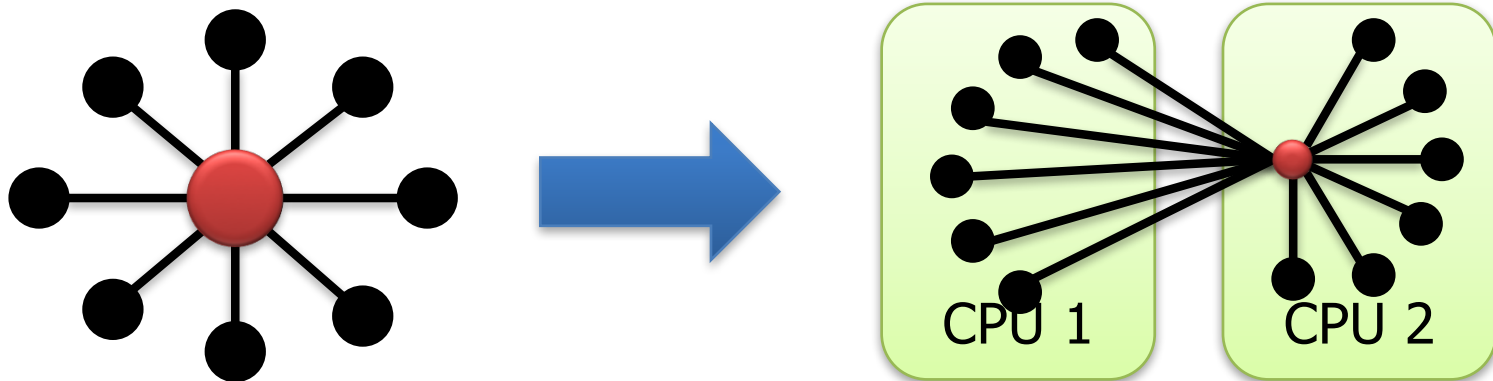
Serializability Requires
Heavy Locking

Graph Partitioning

- Graph parallel abstractions rely on partitioning:
 - Minimize communication
 - Balance computation and storage



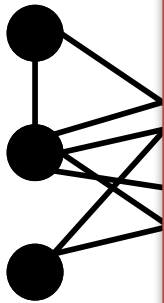
Power-Law Graphs are Difficult to Partition



- Power-Law graphs do not have **low-cost** balanced cuts [*Leskovec et al. 08, Lang 04*]
- Traditional graph-partitioning algorithms perform poorly on Power-Law Graphs. [*Abou-Rjeili et al. 06*]

Random Partitioning

- GraphLab resorts to **random** (hashed) partitioning on **natural graphs**

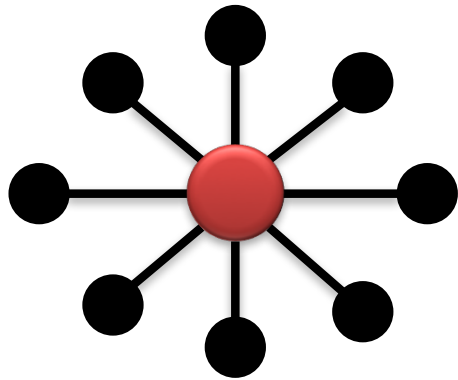

$$\mathbb{E} \left[\frac{|Edges \text{ Cut}|}{|E|} \right] = 1 - \frac{1}{p}$$

10 Machines → 90% of edges cut

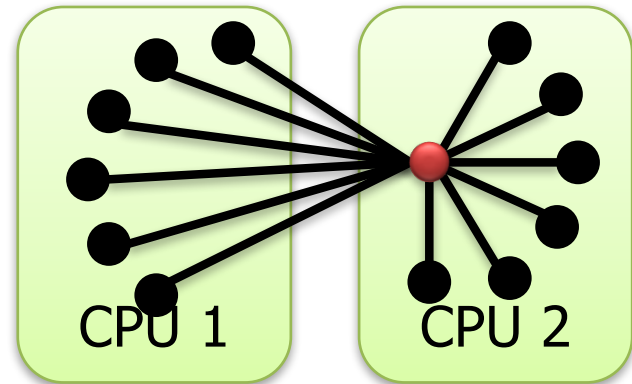
100 Machines → 99% of edges cut!

In Summary

GraphLab is not well suited for
natural graphs



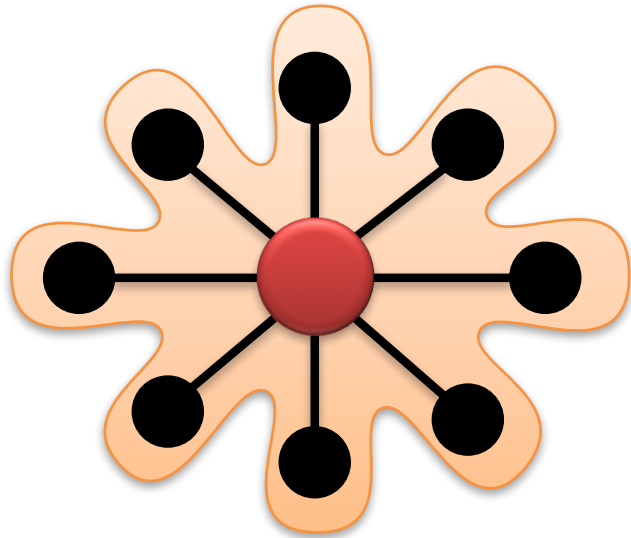
Challenges of **high-degree vertices**



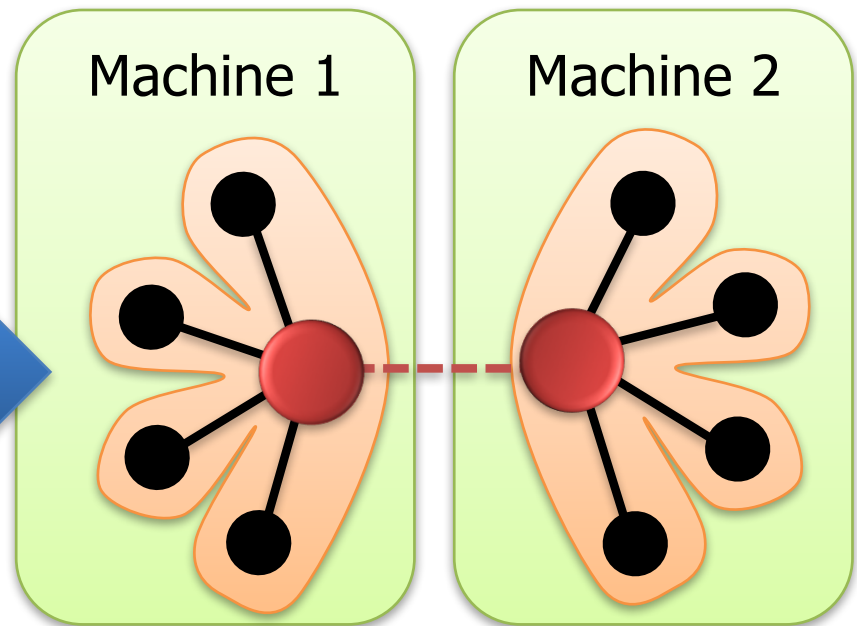
Low quality
partitioning

PowerGraph

Program
For This



Run on This



- Split **High-Degree** vertices
- **New Abstraction** → Equivalence on Split Vertices

A Common Pattern for Vertex-Programs

GraphLab_PageRank(i)

```
// Compute sum over neighbors  
total = 0  
foreach( j in neighbors(i)):  
    total = total + R[j] * wji
```

**Gather Information
About Neighborhood**

```
// Update the PageRank  
R[i] = total
```

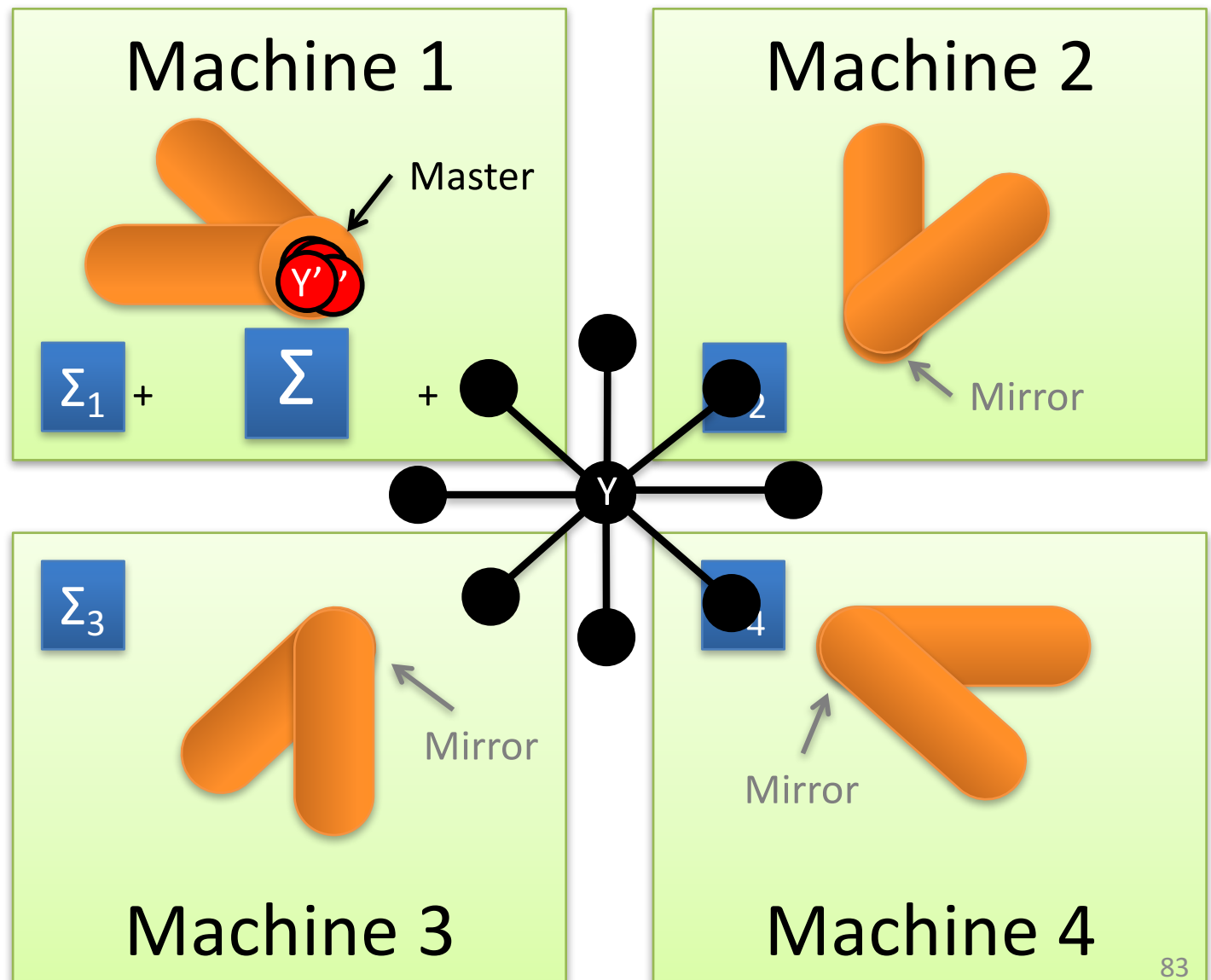
Update Vertex

```
// Trigger neighbors to run again  
priority = |R[i] - oldR[i]|  
if R[i] not converged then  
    signal neighbors(i) with priority
```

**Signal Neighbors &
Modify Edge Data**

GAS Decomposition

Gather
Apply
Scatter



Minimizing Communication in PowerGraph

New Theorem:

*For **any edge-cut** we can directly construct a vertex-cut which requires **strictly less** communication and storage.*

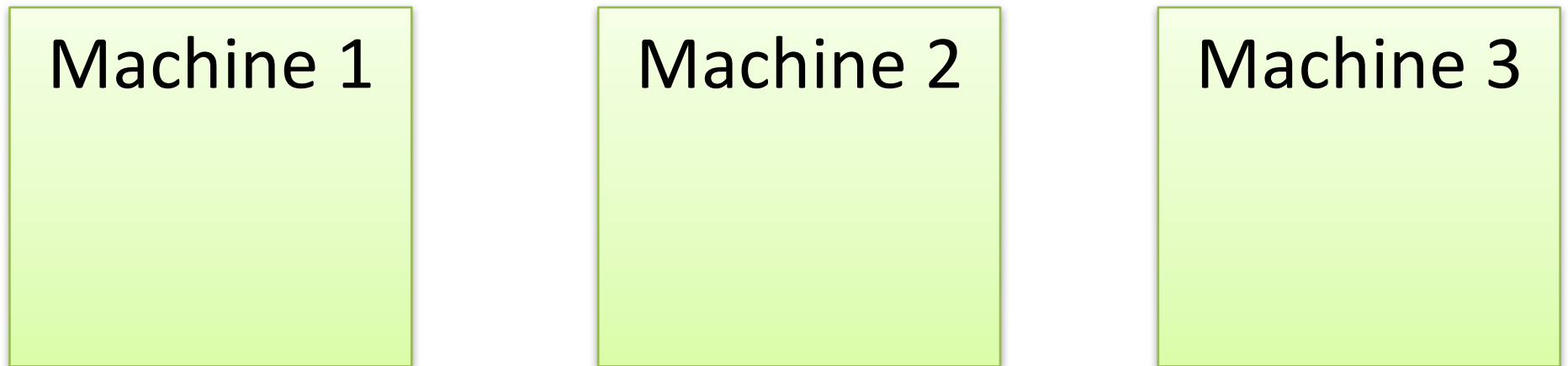
*Percolation theory suggests that power law graphs have **good vertex cuts**. [Albert et al. 2000]*

Constructing Vertex-Cuts

- **Evenly** assign **edges** to machines
 - Minimize machines spanned by each vertex
- Assign each edge **as it is loaded**
 - Touch each edge only once
- Propose two **distributed** approaches:
 - *Random Vertex Cut*
 - *Greedy Vertex Cut*

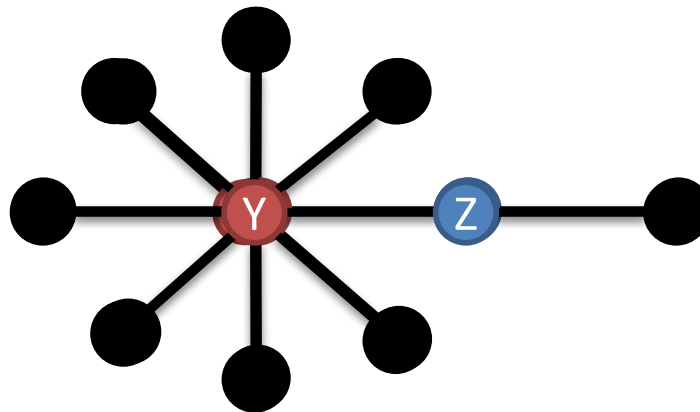
Random Vertex-Cut

- Randomly assign edges to machines



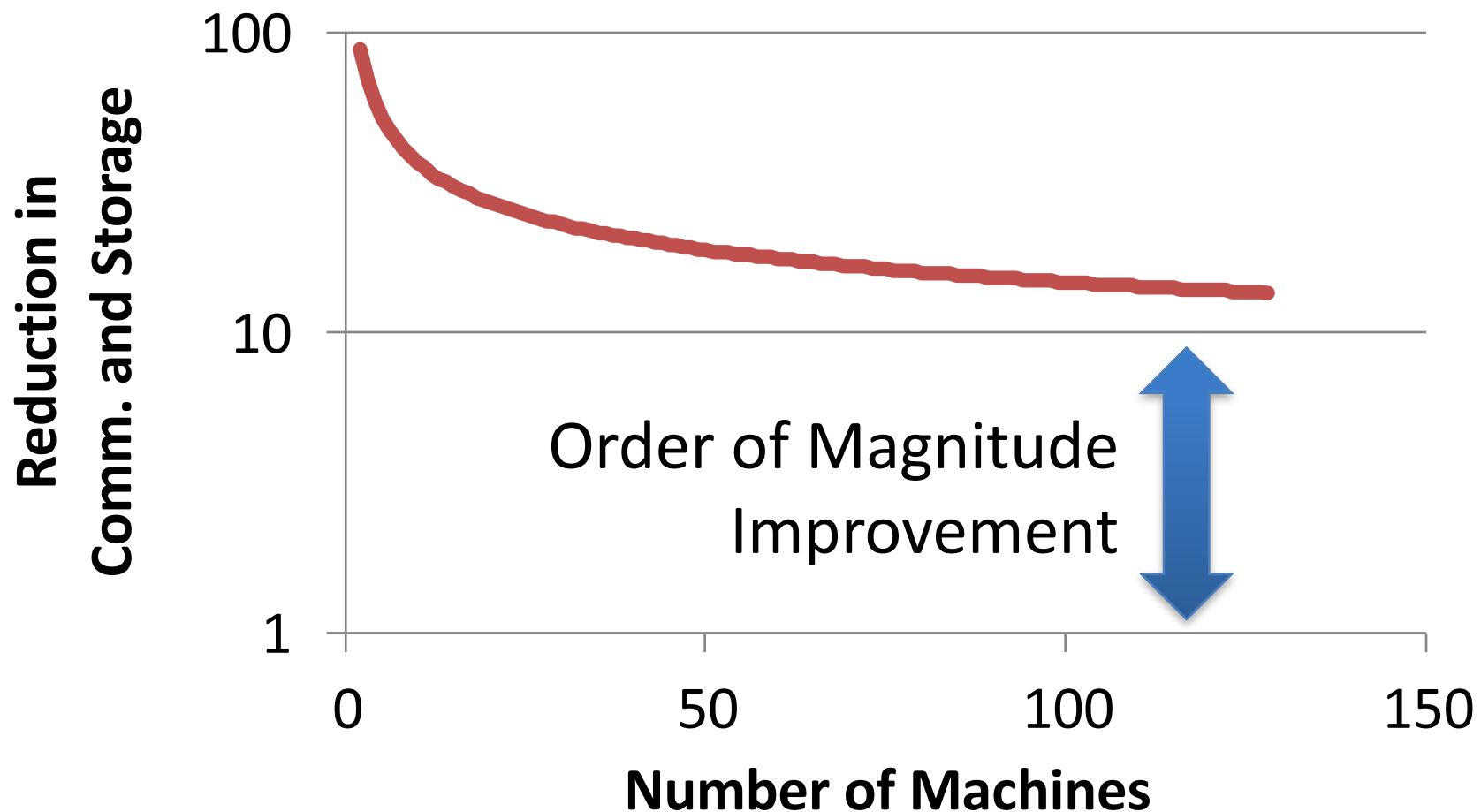
Balanced Vertex-Cut

- Y** Spans 3 Machines
- Z** Spans 2 Machines
- Not cut!



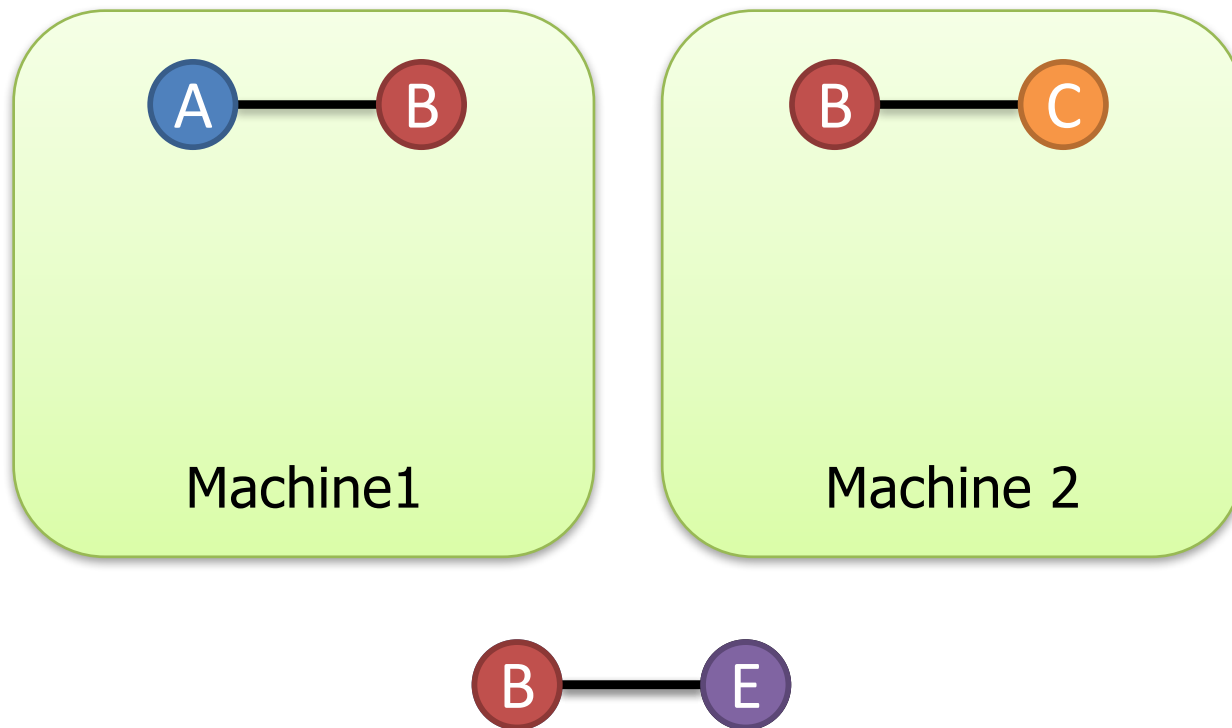
Random Vertex-Cuts vs. Edge-Cuts

- Expected improvement from vertex-cuts:

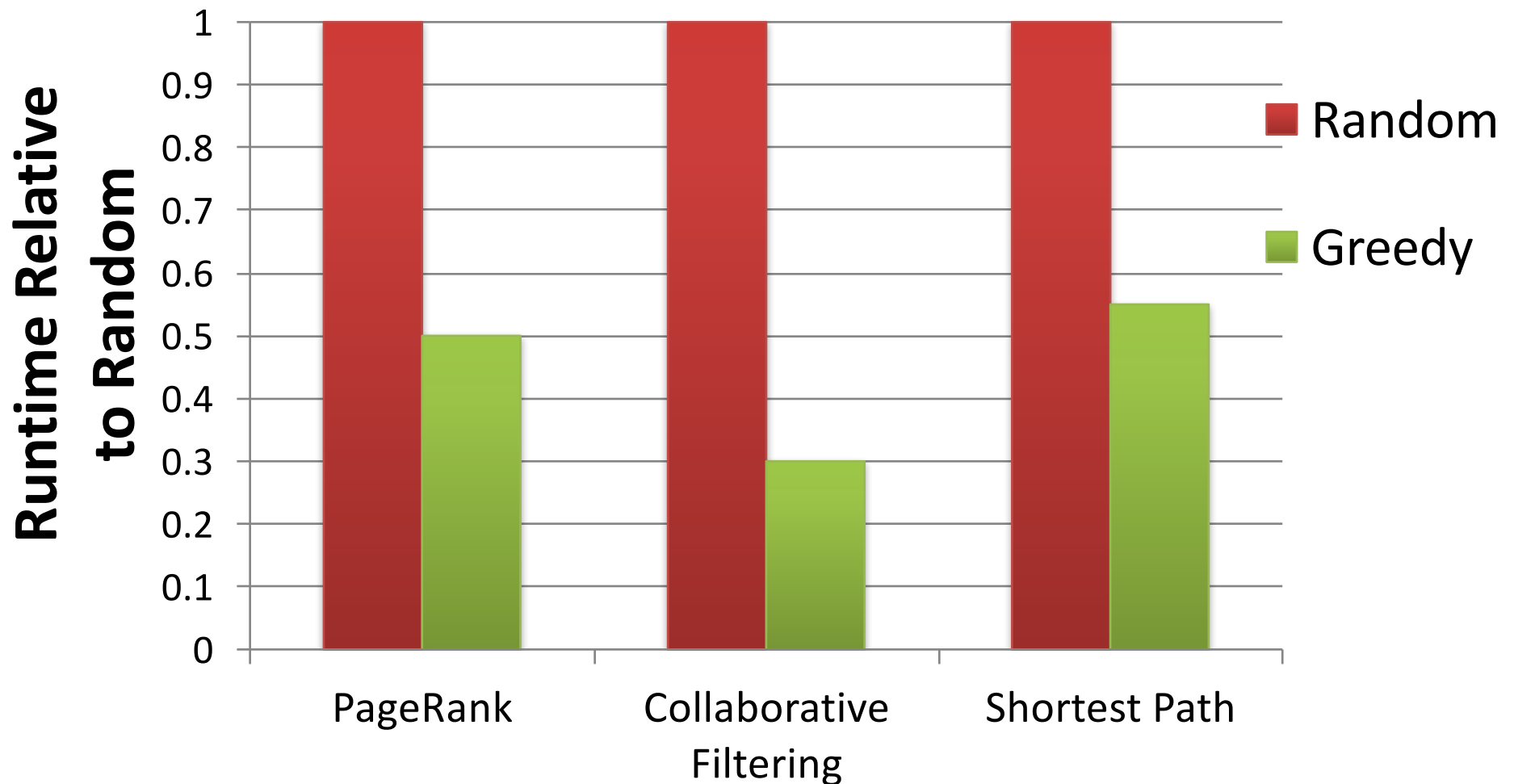


Greedy Vertex-Cuts

- Place edges on machines which already have the vertices in that edge.

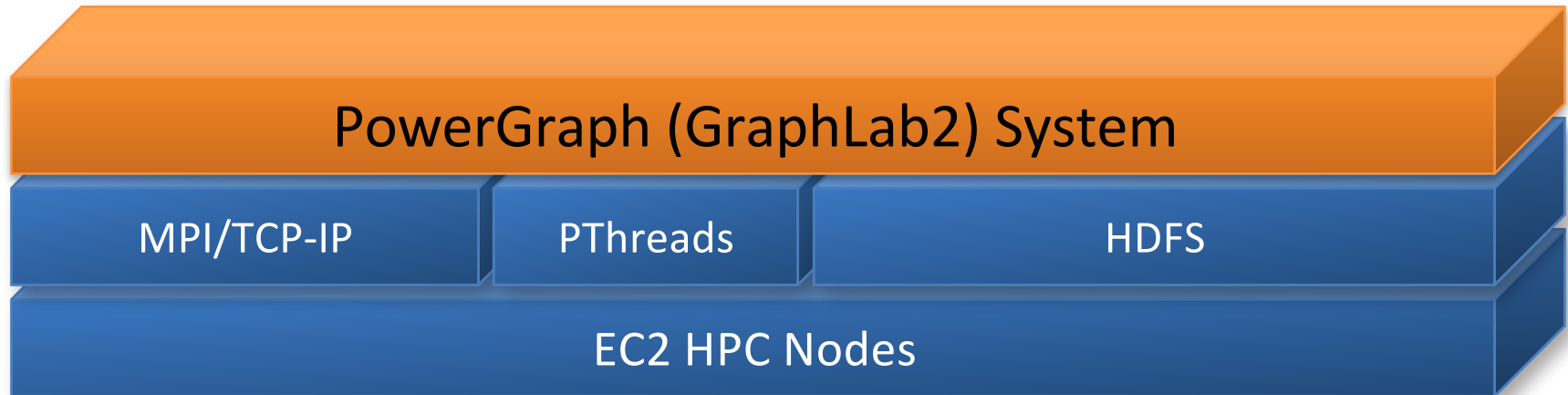


Greedy Vertex-Cuts Improve Performance



Greedy partitioning improves computation performance.

System Design

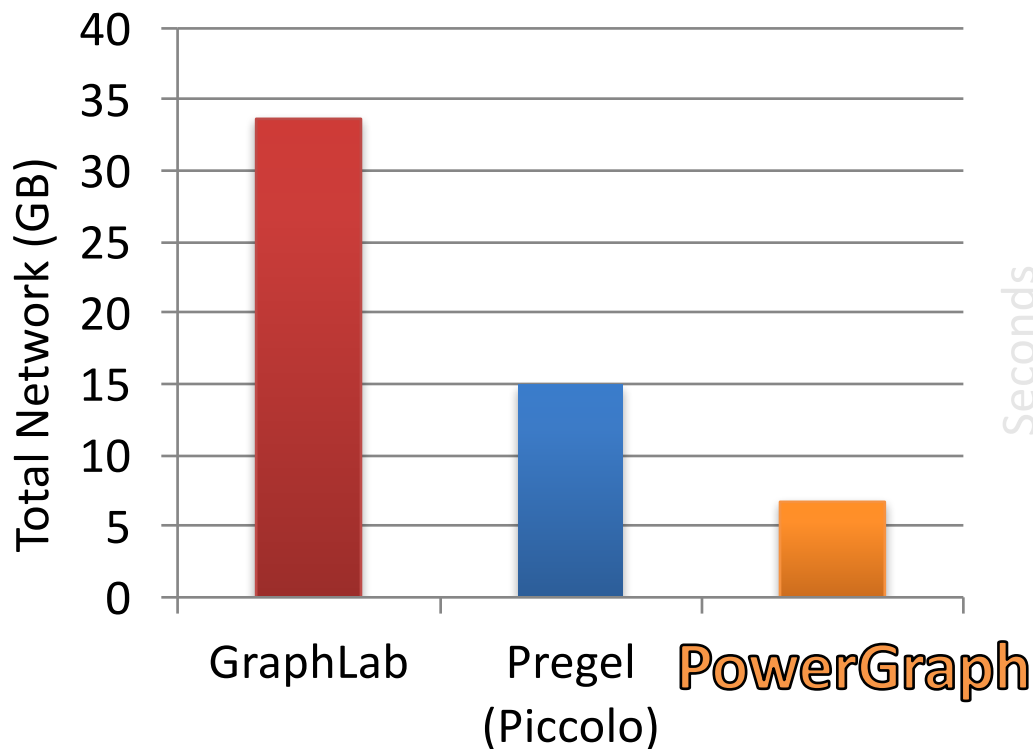


- Implemented as C++ API
- Uses HDFS for Graph Input and Output
- Fault-tolerance is achieved by check-pointing
 - Snapshot time < 5 seconds for twitter network

PageRank on the Twitter Follower Graph

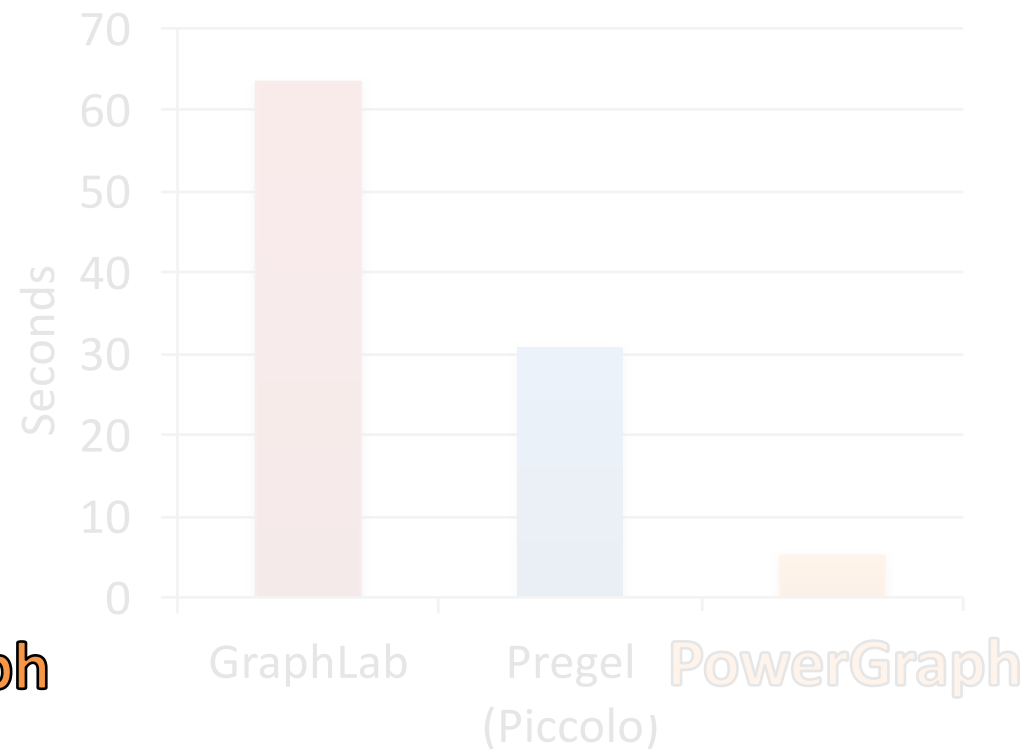
Natural Graph with 40M Users, 1.4 Billion Links

Communication



Reduces Communication

Runtime



Runs Faster

32 Nodes x 8 Cores (EC2 HPC cc1.4x)

PowerGraph is Scalable

Yahoo Altavista Web Graph (2002):

One of the largest publicly available web graphs

1.4 Billion Webpages, 6.6 Billion Links

7 Seconds per Iter.

1B links processed per second

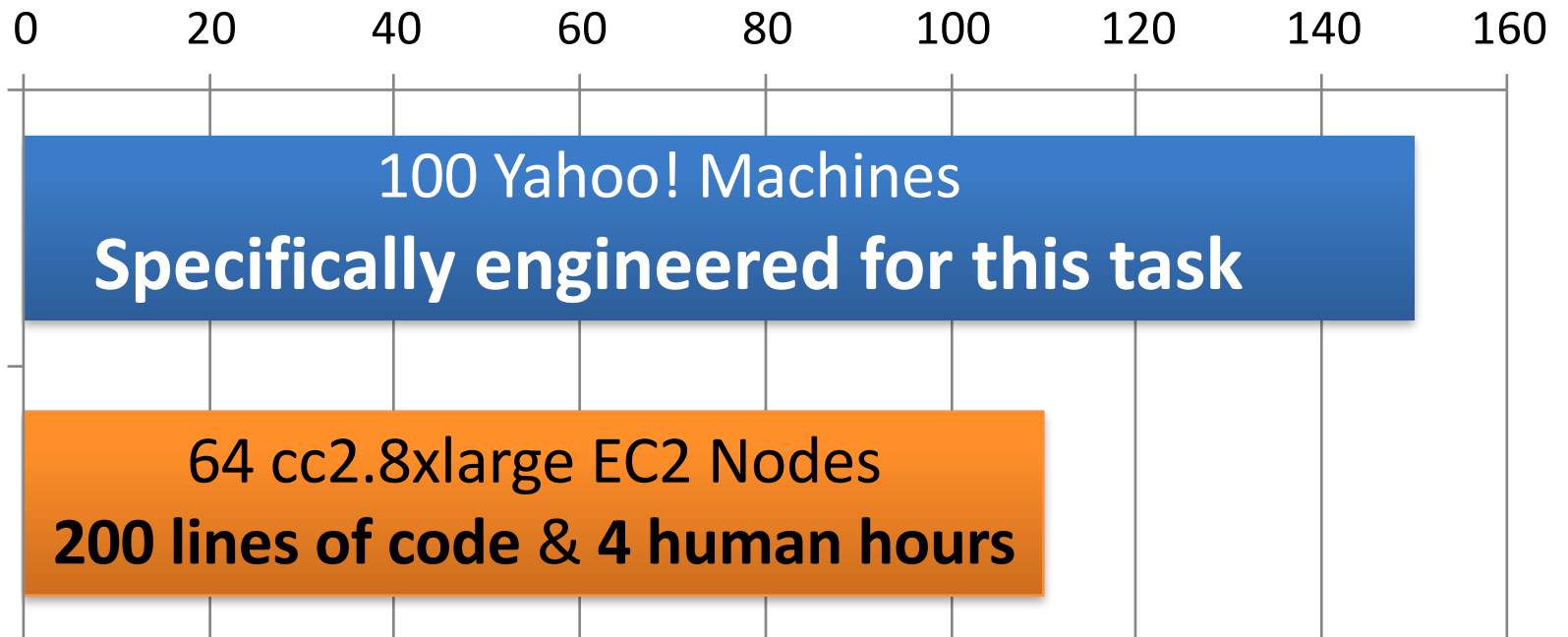
30 lines of user code

Topic Modeling



- English language Wikipedia
 - 2.6M Documents, 8.3M Words, 500M Tokens
 - Computationally intensive algorithm

Million Tokens Per Second

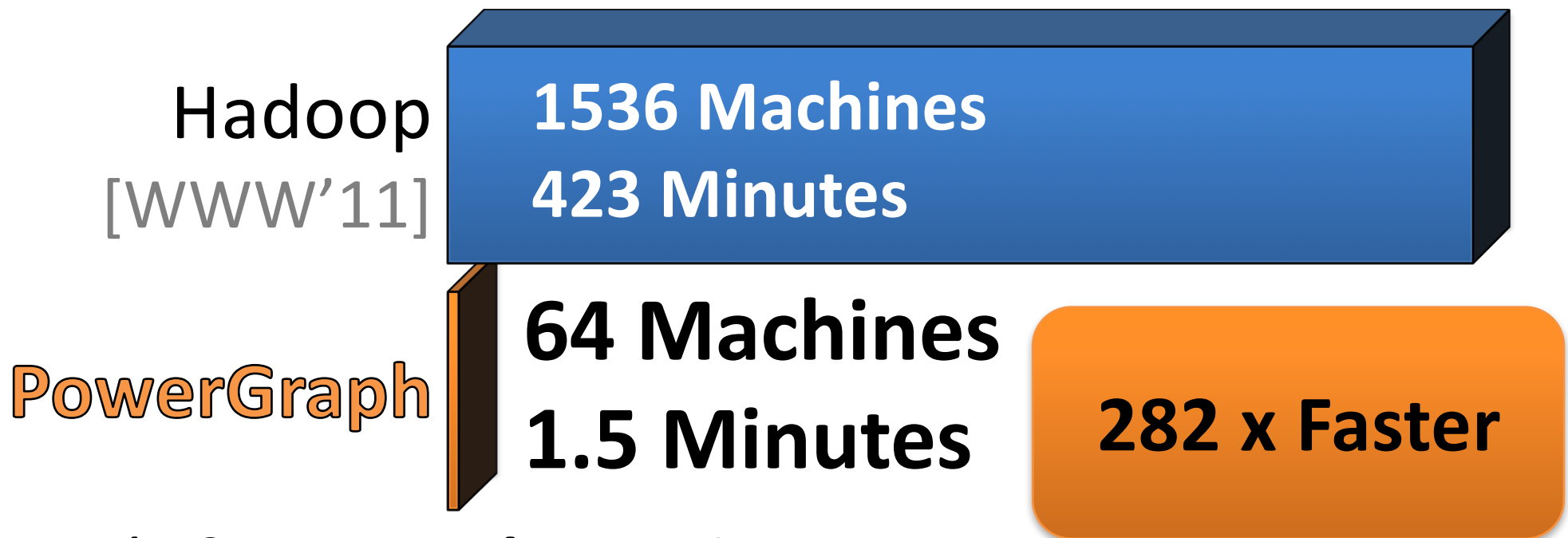


PowerGraph

Triangle Counting on The Twitter Graph

Identify individuals with **strong communities**.

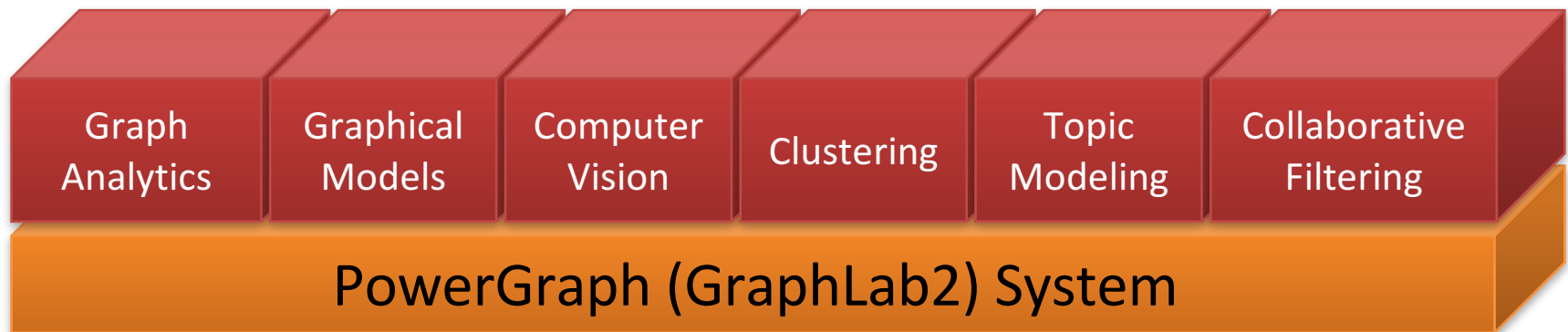
Counted: 34.8 Billion Triangles



Why? **Wrong Abstraction** →

Broadcast $O(\text{degree}^2)$ messages per Vertex

Machine Learning and Data-Mining Toolkits

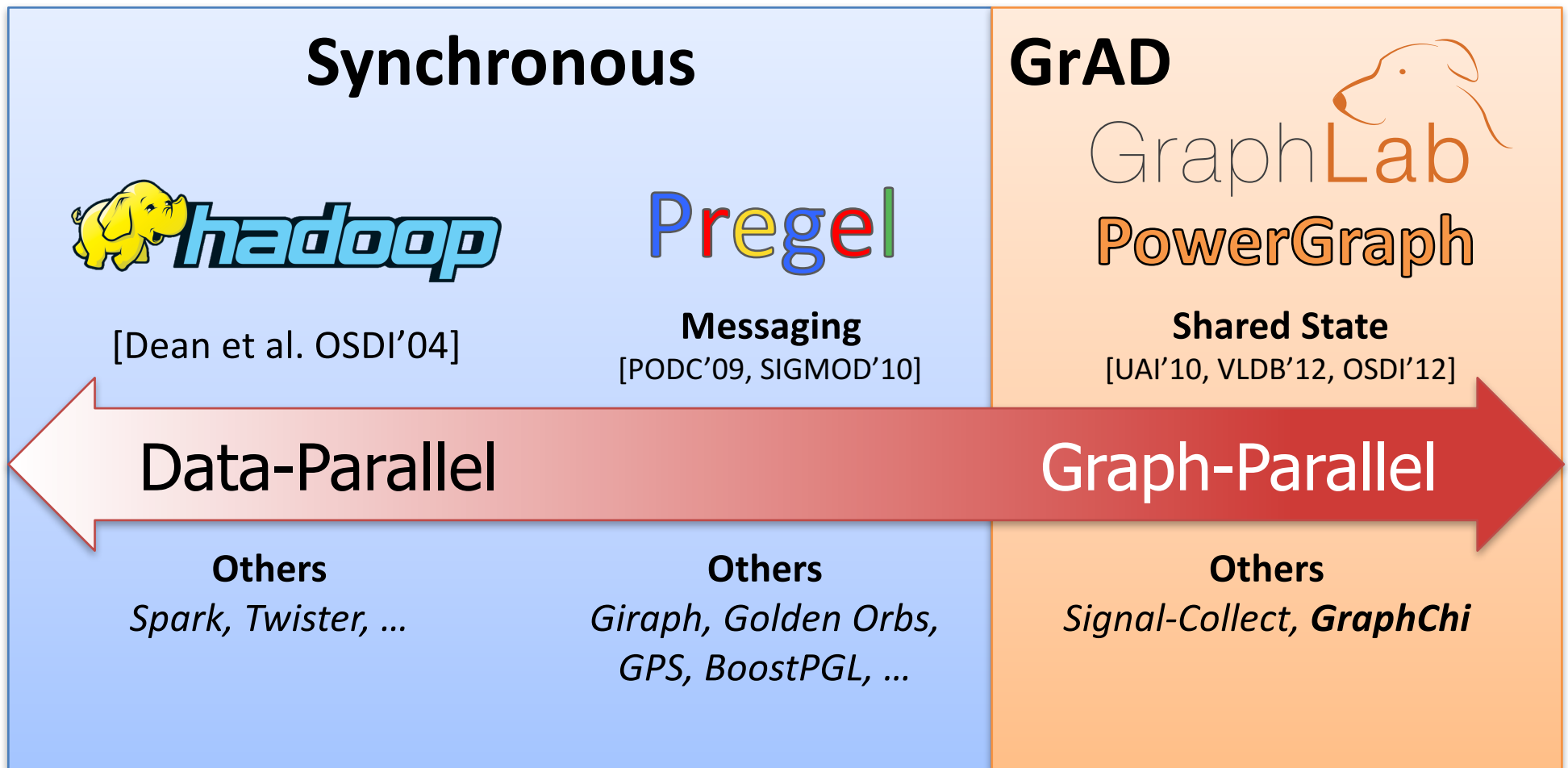


Demonstrates the Applicability
of the GrAD Methodology

Summary: **PowerGraph**

- Identify the **challenges** of Natural Graphs
 - High-degree vertices, Low-quality edge-cuts
- Solution **PowerGraph** System
 - **GAS Decomposition**: split vertex programs
 - **Vertex-partitioning**: distribute natural graphs
- PowerGraph **theoretically** and **experimentally** outperforms existing graph-parallel systems.

Related High-Level Abstractions





facebook

flickr

You Tube

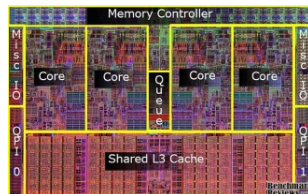
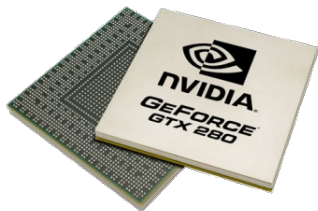
Massive Structured Problems

Probabilistic Graphical Models

Parallel and **Distributed** Algorithms
for Probabilistic **Inference**

GraphLab & PowerGraph

Advances Parallel Hardware



Thesis Statment

*Efficient **parallel** and **distributed** systems for probabilistic reasoning follow the **GrAD Methodology***

1. Graphically decomposition:

- Expose parallelism and distribute state

2. Asynchronous scheduling

- Improved convergence and correctness

3. Dynamic prioritization

- Eliminated wasted work

Observations

- *Graphical models* encode **statistical**, **computational**, and **parallel** structure
- *Tradeoff: Convergence and Parallelism*
 - Many things can be computed in parallel
 - Not all parallel computation is productive
- **Approximation** \rightarrow *Increased Parallelism*
 - τ_ε -approximation, approximate sampling
- Power of high-level abstractions
 - Enables the exploration of GrAD methodology

Future: Declarative Models

- Models are *recursive relationships*
 - BP, Gibbs Sampling, PageRank, ...

My Interests Sum of my friends interests

$$A[x_i] = a \left(\sum_{j \in \mathcal{N}[i]} g(A[x_i], A[x_i, x_j], A[x_j]) \right)$$

“Closeness” number of overlapping posts

$$A[x_i, x_j] = s(A[x_i], A[x_i, x_j], A[x_j])$$

- System determines the optimal schedule

Future: Online Probabilistic Reasoning

- The world is rapidly evolving:
 - Make friends and rate movies in real-time
- How do we define and maintain models?
 - **Declarative specification:** *time invariant*
 - τ_ϵ -approximation: **small** change \rightarrow **local** effect
- Exploit ***Power-Law*** structure in change
 - Popular items are rated more frequently
 - Exploit burstiness for better caching

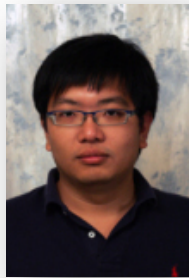
Contributions & Broader Impact

- Theoretically and experimentally characterized
 - Importance of **dynamic asynchronous** scheduling
 - Effect of model **structure** and **parameters** on parallelism
 - Effect of **approximation accuracy** on parallelism
 - Tradeoff between **parallelism** and **convergence**
- Developed two **graph-parallel** abstractions
 - **GraphLab**: vertex-centric view of computation
 - **PowerGraph**: *Distributed* vertex-centric view of computation
- Fostered a community around GraphLab/PowerGraph
 - Substantial industry and academic interest
- Built a foundation for the future design of scalable systems for probabilistic reasoning

Thank You!



Sue Ann
Hong



Yucheng
Low



Aapo
Kyrola



Haijie
Gu



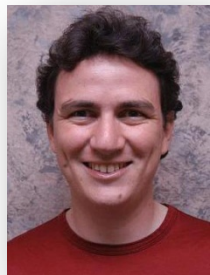
Danny
Bickson



Arthur
Gretton



Andreas
Krause



Carlos
Guestrin



Alex
Smola



Jeff
Bilmes



David
O'Hallaron



Guy
Blelloch



Joe
Hellerstein

The Select Lab & My Family