

A Tiger by the Tail: When Multiplicative Noise Stymies Control

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Abstract—This paper considers the stabilization of an unstable discrete-time linear system that is observed over a channel corrupted by continuous multiplicative noise. The main result is a converse bound that shows that if the system growth is large enough the system cannot be stabilized in a mean-squared sense. This is done by showing that the probability of the state magnitude remains bounded must go to zero with time.

It was known that a system with multiplicative observation noise can be stabilized using a simple linear strategy if the system growth is suitably bounded. However, it was not clear whether non-linear controllers could overcome arbitrarily large growth factors. One difficulty with using the standard approach for a data-rate theorem style converse is that the mutual information per round between the system state and the observation is potentially unbounded with a multiplicative noise observation channel. Our proof technique recursively bounds the conditional density of the system state (instead of focusing on the second moment) to bound the progress the controller can make.

I. INTRODUCTION

We consider the control of a system observed over a multiplicative noise channel, and provide a converse bound to define a region in which the system cannot be stabilized in the mean-square sense. Multiplicative observation noise can model the effects of fast-fading communication channel or sampling and quantization errors [1], [2]. A more detailed discussion of multiplicative noise models is available in [3].

Specifically, we analyze the following system (Fig. 1):

$$\begin{aligned} X_{a,n+1} &= a \cdot X_{a,n} - U_{a,n}, \\ Y_{a,n} &= Z_n \cdot X_{a,n}. \end{aligned} \quad (1)$$

In the preceding formulation, the system state is represented by $X_{a,n}$ at time n , and the control $U_{a,n}$ can be any function of the current and previous observations $Y_{a,0}$ to $Y_{a,n}$. The Z_n 's are i.i.d. random variables with a known continuous distribution. The realization of the noise Z_n is unknown to the controller, much like the fading coefficient of a channel might be unknown to the transmitter or receiver in non-coherent communication. The controller's objective is to stabilize the system in a mean-square sense, i.e. ensure that $\sup_n \mathbb{E}[|X_{a,n}|^2] < \infty$.

For a simple achievable strategy, suppose the Z_n 's are i.i.d. with mean 1 and variance σ^2 . It is easy to verify that a simple memoryless linear strategy can mean-square stabilize the system (1) if and only if $a^2 \leq (1 + \frac{1}{\sigma^2})$.

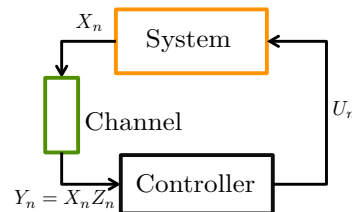


Fig. 1. The state X_n is observed through a multiplicative noise channel, modeled by $Y_n = X_n Z_n$. The gain a is suppressed in the figure notation.

A. Related work

This fundamental problem is inspired by many issues that have been long studied in control and we only mention some representative references here. It is connected to the vast body of work on data-rate theorems and control with communication constraints such as [4], [5]. Nair et al. provide a survey of related problems in [6]. For a system that is growing by a factor of a at each time, the data-rate theorems tell us that a noiseless observation data rate $R > \log |a|$ is necessary and sufficient to stabilize the system [4].

A related problem is that of estimating a linear system over multiplicative noise. While early work on this had been limited to exploring linear estimation strategies [7], some more recent work used a genie-based approach to show a general converse result for the estimation problem over multiplicative noise for both linear and non-linear strategies [2]. The control problem in this paper can also be interpreted as an “active” estimation problem for X_0 , and our converse (that also uses a side-information genie) applies to both linear and non-linear control strategies. However, techniques from the estimation converse or the data-rate theorems do not work for a converse here. Unlike the estimation problem, we cannot describe the distribution of X_n in our problem since the control U_n is arbitrary. For the same reason, we also cannot bound the range of X_n and we cannot bound the rate across the channel to use a data-rate theorem approach.

Another inspiration for our problem is the work on intermittent Kalman filtering [8], [9] and control over packet dropping networks [10] (i.e. estimation and control over Bernoulli multiplicative noise channels). The setup in this paper generalizes those setups to consider a general continuous multiplicative noise on the observation.

Finally, the result in this paper can also be thought of as extending classic uncertainty threshold principle [11] to also understand multiplicative observation uncertainty. The uncertainty threshold principle provides a threshold for the stabilizability of a system with Gaussian uncertainty on the system gain and the control gain.

We would eventually like to establish the control capacity of the system (1), as done for systems with multiplicative noise on the actuation channel (control gain) in [12], [3]. If the controller for a system with multiplicative observation noise is restricted to using linear control strategies, then its performance limit is the same as that of system with multiplicative actuation noise (as in [12]). However, the approach from [12] does not seem to work for the system considered here.

B. Proof approach

This paper introduces a non-standard converse approach and we believe these techniques are a primary contribution of the work. Instead of focusing on the second-moment, our proof bounds the density of the state and thus shows the instability of any moment of the state.

A key element of the proof is that a “genie” observes the state of the system and provides a quantized version of the logarithm of the state to the controller at each time as extra side-information in addition to the multiplicative noise observation. This side information bounds the state in intervals of size 2^{-k} around the initial state X_0 (with k increasing as time increases). We know from results on non-coherent communication [13] and carry-free models [14] that only the order of magnitude of the message can be recovered from a transmission with multiplicative noise. As a result, this side-information does not effectively provide much extra information, but it allows us to quantify the rate at which the controller may make progress.

This paper focuses on a Gaussian distribution for the multiplicative noise, but the ideas can be extended to general distributions under mild assumptions.

II. PROBLEM STATEMENT AND SETUP

Consider the system S_a in (1). For simplicity, let the initial state $X_{a,0}$ be distributed as $X_{a,0} \sim \mathcal{N}(0, 1)$. Let Z_n be i.i.d. Gaussian random variables such that $Z_n \sim \mathcal{N}(1, \sigma^2)$, $f_Z(z) = \frac{1}{\sqrt{2\pi}} e^{-\frac{(z-1)^2}{2\sigma^2}}$. Without loss of generality we can assume that the Z_n have mean 1.

In this paper we show that there exists $a > 0$ such that $P(|X_{a,n}| < M) \rightarrow 0$ as $n \rightarrow \infty$ for all M . This implies that for the same a , the second moment $\mathbb{E}[X_{a,n}^2]$ is always unbounded, i.e. the system cannot be stabilized in the second-moment sense.

To obtain this threshold it suffices to consider the system, S (i.e. set $a = 1$.) The distributions on the Z_n variables and the initial state X_0 are the same as system S_a .

$$\begin{aligned} X_{n+1} &= X_n - U_n, \\ Y_n &= Z_n \cdot X_n. \end{aligned} \quad (2)$$

The system S_a behaves as S scaled by a^n when the control $U_{a,n} = a^n U_n$ is applied. A similar result is given in [3, Ch. 5, Lemma 5.4.4]. Hence, we can write:

$$\mathbb{P}(|X_{a,n}| < M) = \mathbb{P}(|X_n| < a^{-n} M).$$

A. Notation and definitions

Let $S_n = \sum_{i=0}^{n-1} U_i$ for $n \geq 1$. Hence, $X_n = X_0 - S_n$.

The goal of the controller is to have S_n be as close to X_0 as possible. We will track the progress of the controller through intervals I_n that contain X_0 and are decreasing in length.

Let $d(I_n, S) := \inf_{x \in I_n} |S - x|$ denote the distance of a point S from the interval I_n .

Definition 2.1: For all $n \geq 0$ and for any $k \in \mathbb{Z}$, there exists a unique integer $h(k)$ such that $X_0 \in [\frac{h(k)}{2^k}, \frac{h(k)+1}{2^k})$. Let $J(k) := [\frac{h(k)}{2^k}, \frac{h(k)+1}{2^k})$. We now inductively define the intervals I_n using the intervals $J(k)$. Define

$$K_0 := \min\{k \geq 0 \mid d(J(k), 0) \geq 2^{-k}\}, \text{ and}$$

$$K_n := \min\{k \mid k > K_{n-1}, d(J(k), S_n) \geq 2^{-k}\}.$$

Write $H_n := h(K_n)$ and $I_n := [\frac{H_n}{2^{K_n}}, \frac{H_n+1}{2^{K_n}})$.

Let Y_0^n indicate the observations Y_0 to Y_n , and let $\mathcal{F}_n := \{Y_0^n, K_0^n, H_0^n\}$, which is the total information available to the controller at time n . Let $f_{X_n}(x|\mathcal{F}_n)$ be the conditional density of X_n given \mathcal{F}_n .

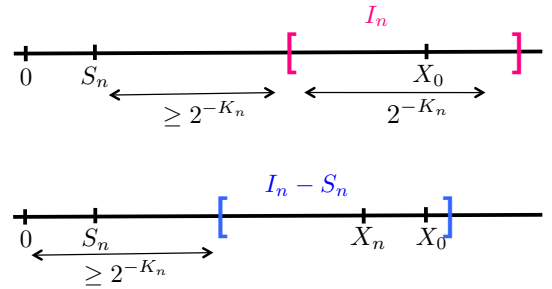


Fig. 2. A caricature illustrating the intervals I_n and $I_n - S_n$.

B. Side-information lemma

This lemma uses K_n to bound how fast S_n approaches X_0 . The easy proof is deferred to the Appendix.

Lemma 2.1: $2^{-K_n} \leq |X_0 - S_n|$, and if $K_n > K_{n-1} + 1$, then $|X_0 - S_n| \leq 2^{2-K_n}$.

III. CONVERSE RESULT

Theorem 3.1: There exists $a \in \mathbb{R}, 0 < a < \infty$ such that $\mathbb{P}(|X_{a,n}| < M) \rightarrow 0$ for all $M < \infty$.

Strategy: The proof provides an exponential upper bound on the density $f_{X_n}(x|\mathcal{F}_n)$, which implies that the controller cannot localize the state within an exponentially shrinking box. The bound on the density is obtained by estimating the change in the density from time n to $n+1$ due to the observation Y_n . The interval I_n (i.e. H_n and K_n) is provided to the controller as side information at time n , which helps generate the bound.

The first step of the proof uses Lemma 3.2 to recursively bound the ratio $\frac{f_{X_n}(x|\mathcal{F}_n)}{f_{X_n}(w|\mathcal{F}_n)}$. This leads to an exponential bound

on the density $f_{X_n}(x | \mathcal{F}_n)$ in terms of the side-information K_n in (8).

This bound helps us generate bounds on the probability the event of interest $\{|X_n| < a^{-n}M\}$, which we show then must go to 0 as $n \rightarrow \infty$.

Proof: Consider

$$\begin{aligned} & f_{X_n}(x | \mathcal{F}_n) \\ &= f_{X_n}(x | Y_n, K_n, H_n, \mathcal{F}_{n-1}) \\ &= \frac{f_{Y_n, K_n, H_n}(Y_n, K_n, H_n | X_n = x, \mathcal{F}_{n-1}) \cdot f_{X_n}(x | \mathcal{F}_{n-1})}{f_{Y_n, K_n, H_n}(Y_n, K_n, H_n | \mathcal{F}_{n-1})} \end{aligned}$$

Formally, f_{Y_n, K_n, H_n} is a density with respect to a product of the Lebesgue measure and two counting measures.

Since $X_0 \in I_n$, the controller knows that $X_n \in I_n - S_n$, where $I_n - S_n$ represents the interval I_n shifted by S_n . We can calculate the ratio of the densities at $x, w \in I_n - S_n$ as:

$$\frac{f_{X_n}(x | \mathcal{F}_n)}{f_{X_n}(w | \mathcal{F}_n)} = \frac{f_{X_n}(x | \mathcal{F}_{n-1}) f_{Y_n}(Y_n | X_n = x, \mathcal{F}_{n-1})}{f_{X_n}(w | \mathcal{F}_{n-1}) f_{Y_n}(Y_n | X_n = w, \mathcal{F}_{n-1})}. \quad (3)$$

Since K_n and H_n are defined by I_n , the conditional distributions of K_n and H_n given $X_n = x$ and $X_n = w$ are equal for $x, w \in I_n - S_n$. So these terms cancel when we consider a ratio, giving (3).

Taking logarithms and using the triangle inequality gives the following recursive lemma, which we prove in the Appendix.

Lemma 3.2:

$$\left| \log \frac{f_{X_n}(x | \mathcal{F}_n)}{f_{X_n}(w | \mathcal{F}_n)} \right| \leq \frac{1}{\sigma^2} (\sigma^2 + 2Z_n(2Z_n - 1)) 2^{K_n} |x - w| + \left| \log \frac{f_{X_n}(x | \mathcal{F}_{n-1})}{f_{X_n}(w | \mathcal{F}_{n-1})} \right|. \quad (4)$$

Now, based on the control law we know that

$$f_{X_n}(x | \mathcal{F}_{n-1}) = f_{X_{n-1}}(x + U_{n-1} | \mathcal{F}_{n-1}),$$

since U_{n-1} is \mathcal{F}_{n-1} measurable. Substituting this into (4) and unfolding recursively gives:

$$\left| \log \frac{f_{X_n}(x | \mathcal{F}_n)}{f_{X_n}(w | \mathcal{F}_n)} \right| \leq \sum_{i=1}^n \frac{1}{\sigma^2} (\sigma^2 + 2Z_i(2Z_i - 1)) 2^{K_i} |x - w| + \left| \log \frac{f_{X_0}(x + S_n)}{f_{X_0}(w + S_n)} \right|. \quad (5)$$

The inequality (5) separates the effect of the uncertainty due to X_0 and the subsequent uncertainty due to the observations and control.

To bound the effect of the initial state let $\eta_n = \max_{x, w \in I_n - S_n} \left| \log \frac{f_{X_0}(x + S_n)}{f_{X_0}(w + S_n)} \right|$. Since I_n is an interval of size at most 2^{-n} that contains X_0 we get that:

$$|\eta_n| \leq \frac{1}{2} |(X_0 + 2^{-n})^2 - (X_0 - 2^{-n})^2| \leq 2^{1-n} |X_0|. \quad (6)$$

Now, we define

$$\Psi_n = 2^{-K_n} \frac{1}{\sigma^2} \sum_{i=1}^n (\sigma^2 + 2Z_i(2Z_i - 1)) 2^{K_i}, \quad (7)$$

and rewrite (5) as:

$$\left| \log \frac{f_{X_n}(x | \mathcal{F}_n)}{f_{X_n}(w | \mathcal{F}_n)} \right| \leq \Psi_n \cdot 2^{K_n} \cdot |x - w| + \eta_n.$$

Finally, we are in a position to bound $f_{X_n}(x | \mathcal{F}_n)$:

$$f_{X_n}(x | \mathcal{F}_n) \leq (\exp\{\Psi_n 2^{K_n} |x - w| + \eta_n\}) f_{X_n}(w | \mathcal{F}_n). \quad (8)$$

We integrate (8) over an interval of length $\gamma = 2^{(-K_n - T)}$ with x at one end point. So $|x - w| \leq 2^{(-K_n - T)}$. Here, $T \in \mathbb{R}$ is a constant that we will choose later. Such an interval can be fit into I_n to the left or right of any x depending on where x is in the interval. Assuming without loss of generality that x is the left endpoint of the integration interval we compute that

$$\int_x^{x+\gamma} f_{X_n}(x | \mathcal{F}_n) dw \leq \int_x^{x+\gamma} (e^{\Psi_n 2^{K_n} |x-w| + \eta_n}) f_{X_n}(w | \mathcal{F}_n) dw$$

Bound $|x - w|$ on the RHS by $2^{(-K_n - T)}$. Integrating gives

$$\begin{aligned} \gamma \cdot f_{X_n}(x | \mathcal{F}_n) &\leq \int_x^{x+\gamma} (e^{\Psi_n 2^{K_n} 2^{(-K_n - T)} + \eta_n}) f_{X_n}(w | \mathcal{F}_n) dw \\ &\leq (e^{\Psi_n 2^{-T} + \eta_n}) \cdot 1, \end{aligned}$$

since the density integrates out to 1. Hence,

$$f_{X_n}(x | \mathcal{F}_n) \leq e^{\Psi_n 2^{-T} + \eta_n} 2^{K_n + T}. \quad (9)$$

This gives us a bound on the density of X_n in terms of K_n . Since the K_n 's are tracking the magnitude of X_n what remains to be done is to bound the growth of the K_n 's. To do this we first state the following lemma about the crucial quantity Ψ_n .

Lemma 3.3: For $T > C_\sigma$, where C_σ is a constant that depends only on σ , the expectation $\mathbb{E}[e^{\Psi_n 2^{-T}}]$ is uniformly bounded for all n .

The next lemma uses Lemma 3.3 and (9) to show that the K_n grow at most linearly.

Lemma 3.4: There exists a constant $C = C_{\sigma, T}$ depending only on (σ, T) such that

$$\mathbb{P}(K_n - K_0 > C_{\sigma, T} \cdot n) \rightarrow 0.$$

The proofs of Lemmas 3.3, 3.4 are deferred to the Appendix.

Let G_n denote the event that $K_n - K_0 > Cn$, and G_n^c its complement. Then the event $\{|X_n| < a^{-n}M\}$ is covered by:

$$\begin{aligned} & \mathbb{P}(|X_n| < a^{-n}M) \\ & \leq \mathbb{P}(G_n) + \mathbb{P}(K_0 > n) + \mathbb{P}(|X_n| \leq a^{-n}M, G_n^c, K_0 \leq n). \end{aligned} \quad (10)$$

We evaluate the three terms one by one. For the first term in (10), we have $\mathbb{P}(G_n) = \mathbb{P}(K_n - K_0 > Cn) \rightarrow 0$ as $n \rightarrow \infty$ from Lemma 3.4.

The second term, $\mathbb{P}(K_0 > n)$, captures the case where the initial state X_0 might be very close to zero. However, eventually this advantage dies out for large enough n , since $\mathbb{P}(X_0 < 2^{-n}) \rightarrow 0$ as $n \rightarrow \infty$.

The last term in (10) remains. By the law of iterated expectation:

$$\begin{aligned} & \mathbb{P}(|X_n| < a^{-n}M, G_n^c, K_0 \leq n) \\ &= \mathbb{E}[\mathbb{P}(|X_n| < a^{-n}M, G_n^c, K_0 \leq n \mid \mathcal{F}_n)]. \end{aligned}$$

We focus on the term conditioned on \mathcal{F}_n :

$$\begin{aligned} & \mathbb{P}(|X_n| < a^{-n}M, G_n^c, K_0 \leq n \mid \mathcal{F}_n) \\ &= \mathbb{E}[\mathbb{1}_{\{|X_n| < a^{-n}M\}} \mathbb{1}_{\{G_n^c\}} \mathbb{1}_{\{K_0 \leq n\}} \mid \mathcal{F}_n] \\ &= \mathbb{P}(|X_n| < a^{-n}M \mid \mathcal{F}_n) \cdot \mathbb{1}_{\{G_n^c\}} \mathbb{1}_{\{K_0 \leq n\}}. \end{aligned} \quad (11)$$

Now, we can apply (9) to get

$$\begin{aligned} \mathbb{P}(|X_n| < a^{-n}M \mid \mathcal{F}_n) &= \int_{-a^{-n}M}^{a^{-n}M} f_{X_n}(x \mid \mathcal{F}_n) dx \\ &\leq \int_{-a^{-n}M}^{a^{-n}M} e^{\eta + \Psi_n 2^{-T}} 2^{K_n + T} dx \\ &= 2Ma^{-n} \cdot e^{\eta + \Psi_n 2^{-T}} \cdot 2^{K_n + T}. \end{aligned}$$

Then we can bound (11) as

$$\begin{aligned} & \mathbb{P}(|X_n| < a^{-n}M \mid \mathcal{F}_n) \cdot \mathbb{1}_{\{G_n^c\}} \mathbb{1}_{\{K_0 \leq n\}} \\ &\leq 2Ma^{-n} \cdot e^{\eta + \Psi_n 2^{-T}} \cdot 2^{(C+1)n+T}, \end{aligned} \quad (12)$$

since $K_n \leq Cn + K_0$ and $K_0 \leq n$ implies $K_n \leq (C+1)n$. Taking expectations on both sides we get:

$$\begin{aligned} & \mathbb{P}(|X_n| < a^{-n}M, G_n^c, K_0 \leq n) \\ &\leq 2Ma^{-n} \cdot 2^{(C+1)n+T} \cdot \mathbb{E}[e^{\eta + \Psi_n 2^{-T}}]. \end{aligned} \quad (13)$$

By Lemma 3.3 and (6), the above expression (13) tends to 0 for $a > 2^{C+1}$. Thus, all three probabilities in (10) converge to 0 as $n \rightarrow \infty$. Hence, if $a > 2^{C+1}$ then $\mathbb{P}(|X_n| < a^{-n}M) \rightarrow 0$ for all M . ■

IV. CONCLUSION AND FUTURE WORK

This paper provides a first proof-of-concept converse for a control system observed over continuous multiplicative noise. However, there is an exponential gap between the scaling behavior of the achievable strategy and the converse. It still remains to be seen if the achievable region can grow with non-linear controllers to significantly beat the $a^2 < (1 + \frac{1}{\sigma^2})$ bound, and computing the control capacity of the system remains open.

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APPENDIX

A. Proof of Lemma 2.1

From the definition of I_n , we know that $d(I_n, S_n) \geq 2^{-K_n}$. Hence, I_n cannot contain S_n . This gives $2^{-K_n} \leq |X_0 - S_n|$, since $X_0 \in I_n$.

To show the second half of the inequality, suppose that $|X_0 - S_n| > 2^{2-K_n}$. Then,

$$2^{1-K_n} < |X_0 - S_n| - 2^{1-K_n}.$$

Hence, there exists a smaller interval $J(K_n - 1)$ that contains X_0 such that

$$2^{1-K_n} < d(J(K_n - 1), S_n),$$

where $J(K_n - 1)$ is an interval of length $2^{1-K_n} > 2^{-K_n}$. Since we also assumed that $K_n > K_{n-1} + 1$, this contradicts the assumption that K_n was the minimal $k > K_{n-1}$ such that $d(J(k), S_n) \geq 2^{-k}$. ■

B. Proof of Lemma 3.2

We take logarithms on both sides of (3) and apply the triangle inequality to get

$$\begin{aligned} & \left| \log \frac{f_{X_n}(x \mid \mathcal{F}_n)}{f_{X_n}(w \mid \mathcal{F}_n)} \right| \\ &\leq \left| \log \frac{f_{Y_n}(Y_n \mid X_n = x, \mathcal{F}_{n-1})}{f_{Y_n}(Y_n \mid X_n = w, \mathcal{F}_{n-1})} \right| + \left| \log \frac{f_{X_n}(x \mid \mathcal{F}_{n-1})}{f_{X_n}(w \mid \mathcal{F}_{n-1})} \right|. \end{aligned} \quad (14)$$

Using $f_{Y_n}(Y_n|X_n = x, \mathcal{F}_{n-1}) = \frac{1}{\sqrt{2\pi\sigma|x|}} e^{-(\frac{Y_n}{x}-1)^2 \frac{1}{2\sigma^2}}$, and the applying the triangle inequality again we get,

$$\begin{aligned} & \left| \log \frac{f_{Y_n}(Y_n | X_n = x, \mathcal{F}_{n-1})}{f_{Y_n}(Y_n | X_n = w, \mathcal{F}_{n-1})} \right| \\ & \leq \left| \log |x| - \log |w| \right| + \frac{1}{2\sigma^2} \left| \left(\frac{Y_n}{x} - 1 \right)^2 - \left(\frac{Y_n}{w} - 1 \right)^2 \right|. \end{aligned} \quad (15)$$

We can use the derivatives of the functions to bound the two function differences above. Using $\left| \frac{d}{dx} \log |x| \right| = \frac{1}{|x|}$ and $\frac{d}{dx} \left(\frac{Y_n}{x} - 1 \right)^2 = 2 \left(\frac{Y_n}{x} - 1 \right) \left(\frac{-Y_n}{x^2} \right)$, we can bound (15) as:

$$\begin{aligned} & \left| \log \frac{f_{Y_n}(Y_n | X_n = x, \mathcal{F}_{n-1})}{f_{Y_n}(Y_n | X_n = w, \mathcal{F}_{n-1})} \right| \leq \left(\max_{t \in I_n - S_n} \frac{1}{|t|} \right) |x - w| + \\ & \left(\max_{t \in I_n - S_n} \frac{1}{2\sigma^2} \left| 2 \left(\frac{X_n Z_n}{t} - 1 \right) \left(\frac{X_n Z_n}{t^2} \right) \right| \right) |x - w|. \end{aligned} \quad (16)$$

Since $X_n \in I_n - S_n$, the maximizations are over $t \in I_n - S_n$ in (16). First, note that $\max_{t \in I_n - S_n} \frac{1}{|t|} \leq 2^{K_n}$. Second, for all $t \in I_n - S_n$, we have $\frac{1}{2} \leq \frac{X_n}{t} \leq 2$. Hence, the second term on the RHS of (16) can upper bounded by:

$$\leq \left(\max_{t \in I_n - S_n} \frac{1}{2\sigma^2} \left| 2(2Z_n - 1) \left(\frac{2Z_n}{t} \right) \right| \right) |x - w|.$$

Combining these two ideas we upper bound (16) as:

$$\begin{aligned} & \leq 2^{K_n} |x - w| + \frac{1}{2\sigma^2} (4Z_n(2Z_n - 1) 2^{K_n}) |x - w| \\ & = \left(1 + \frac{2Z_n(2Z_n - 1)}{\sigma^2} \right) 2^{K_n} |x - w|. \end{aligned}$$

This implies the desired bound for (14). \blacksquare

C. Proof of Lemma 3.3

$$\begin{aligned} & \mathbb{E}[e^{\Psi_n 2^{-T}}] \\ & = \mathbb{E}[\exp\{2^{-T-K_n} \frac{1}{\sigma^2} \sum_{i=1}^n (\sigma^2 + 2Z_i(2Z_i - 1)) 2^{K_i}\}]. \end{aligned}$$

Since the K_i 's increase by at least 1 each time, we can replace $2^{-K_n} 2^{K_i}$ by 2^{i-n} . We take absolute values and use the triangle inequality to get the upper bound

$$\mathbb{E}[e^{\Psi_n 2^{-T}}] \leq \mathbb{E}[\exp\{2^{-T} \sum_{i=1}^n \frac{1}{\sigma^2} |\sigma^2 + 2Z_i(2Z_i - 1)| 2^{i-n}\}].$$

Substituting Z_i as $Z_i = \sigma \tilde{Z}_i + 1$, noting that the Z_i are independent and some simple algebra gives the bound below. α, β are constants that depend on σ .

$$\mathbb{E}[e^{\Psi_n 2^{-T}}] \leq \prod_{j=1}^n \mathbb{E} \left[\exp \left\{ 2^{-T} \left(\alpha \tilde{Z}_j^2 + \beta \right) 2^{-j} \right\} \right]. \quad (17)$$

For each term in the product of (17):

$$\begin{aligned} & \mathbb{E} \left[\exp \left\{ 2^{-T} \left(\alpha \tilde{Z}_j^2 + \beta \right) 2^{-j} \right\} \right] \\ & = \exp \left\{ \beta 2^{-T-j} \right\} \int \frac{1}{\sqrt{2\pi}} e^{(\alpha 2^{-T-j} - \frac{1}{2}) z^2} dz. \end{aligned}$$

Choose $T > \log_2(2\alpha)$, so that $\xi := \alpha 2^{-T} < \frac{1}{2}$. Note $0 < \xi < \infty$. The integral is equal to $\frac{1}{\sqrt{1-\xi 2^{1-j}}}$. Then,

$$\prod_{j=1}^n \frac{1}{\sqrt{1-\xi 2^{1-j}}} \leq \prod_{j=1}^n e^{\xi 2^{1-j}} = e^{2\xi \sum_{j=1}^n 2^{-j}} \leq e^{4\xi}.$$

A similar bound can be obtained for the terms with β . The details are omitted and will be included in the full version. \blacksquare

D. Proof of Lemma 3.4

By construction, $K_{n+1} \geq K_n + 1$. In the case where $K_{n+1} > K_n + 1$, we can apply Lemma 2.1 and get that for $\ell \geq 2$

$$\begin{aligned} \mathbb{P}(K_{n+1} - K_n \geq \ell | \mathcal{F}_n) & \leq \mathbb{P}(|X_{n+1}| \leq 2^{2-K_n-\ell} | \mathcal{F}_n) \\ & = \mathbb{P}(|X_n - U_n| \leq 2^{2-K_n-\ell} | \mathcal{F}_n). \end{aligned}$$

This is because the control U_n must have been very close to X_n for K_{n+1} to be much larger than K_n . Then we calculate this probability by integrating out the density.

$$\begin{aligned} \mathbb{P}(|X_n - U_n| \leq 2^{2-K_n-\ell} | \mathcal{F}_n) & = \int_{U_n - 2^{2-K_n-\ell}}^{U_n + 2^{2-K_n-\ell}} f_{X_n}(t | \mathcal{F}_n) dt \\ & \leq 2 \cdot 2^{2-K_n-\ell} \left(\max_t f_{X_n}(t | \mathcal{F}_n) \right). \end{aligned}$$

Combined with (9), this gives us that

$$\begin{aligned} \mathbb{P}(K_{n+1} - K_n \geq \ell | \mathcal{F}_n) & \leq 2 \cdot 2^{2-K_n-\ell} e^{\eta_n + \Psi_n 2^{-T}} 2^{K_n+T} \\ & = 2^{3-\ell+T} \cdot e^{\eta_n + \Psi_n 2^{-T}}. \end{aligned} \quad (18)$$

Let $D_n = K_{n+1} - K_n$ and

$$\tilde{K}_n = \sum_{i=0}^{n-1} (D_i - \mathbb{E}[D_i | \mathcal{F}_i]).$$

It is clear that (\tilde{K}_n) is a martingale with respect to \mathcal{F}_n . In addition, (18) yields that the conditional distribution of D_n given \mathcal{F}_n is stochastically dominated by the distribution of

$$G + \left(3 + T + \frac{|\eta_n + \Psi_n 2^{-T}|}{\ln 2} \right), \quad (19)$$

where G is an independent geometric variable with mean 2. Combined with (6) and (7), it yields that

$$\mathbb{E}(D_i - \mathbb{E}[D_i | \mathcal{F}_i])^2 \leq \mathbb{E}D_i^2 \leq C_1,$$

where C_1 is a constant depending on (T, σ) . This implies

$$\mathbb{P}(\tilde{K}_n \geq n) \rightarrow 0 \text{ as } n \rightarrow \infty. \quad (20)$$

Now, using (19) again, we get that

$$\mathbb{E}[D_i | \mathcal{F}_i] \leq 2(5 + T + |\eta_i + \Psi_i 2^{-T}|). \quad (21)$$

A careful analysis of the quantity $\Delta_n = |\eta_n + \Psi_n 2^{-T}|$ can be used to show that

$$\mathbb{P}\left(\sum_{i=1}^n \mathbb{E}[D_i | \mathcal{F}_i] \geq 10^6(1 + T + \sigma^{-2})n\right) \rightarrow 0,$$

which combined with (20) proves the lemma. Details will be included in the full version. \blacksquare