# Z3str4: A Multi-Armed String Solver 

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## String Solvers and Queries by Example

How Can We Automatically Reason About This Code?

```
def check_password(x: str) -> bool:
    if re.search(r"\d", x) is None:
        return False # no number
    if "12345" in x:
        return False # contains "123456"
    if len(x) < 5:
        return False # too short
    if x[:-len(x)//2] * 2 == x:
        return False # same string twice
    return True
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```
`Model
x\in .*[\d]+..*
"12345" \in 
    |x|<5
    y\cdoty=x
```


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```
x\in .*[\d]+.*
```

" 12345 " $\in x$
$|x|<5$
$y \cdot y=x$

## Solver

## Goal of Z3str4: Solve as many queries as possible, as fast as possible.



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Membership Predicate


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```
F ::=Atom | F F^F | F\veeF | \negF
```



```
Are }::=\mp@subsup{t}{str}{}\inR
A int }\quad::=\mp@subsup{t}{int}{}=\mp@subsup{t}{int}{}\quad|\quad\mp@subsup{t}{int}{}<\mp@subsup{t}{int}{
```




```
    str.to_int (t tstr ) where m\inConint & v\inVarint
t str }::=s | v | t tstr 牴t⿱strr | str.from_int (tint ) | replace (t tstr, , tstr , tstr ) |
```




## Results Overview



10,000 20,000 30,000 40,000 50,000 60,000 70,000 80,000 90,000100,00010,00020,000
Solved instances

## Results on PyEx



## Results on Automark25



## Contributions Overview



## Contribution I: Length Abstraction Solver



## Contribution 2: Extended Arrangement Solver



## Contribution 3: Selection Architecture



## Length Abstraction Solver

## Length Abstraction Solver: Refinement Loop



## Length Abstraction Solver: Solve Length Constraints



## Length Abstraction Solver: Solve Length Constraints



## Length Abstraction Solver: Generate Character Query



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## Length Abstraction Solver: Learn Length Constraint



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## Length Abstraction Solver: Learn Length Constraint



## Length Abstraction Solver: Learn Length Constraint



## Length Abstraction Solver: Refinement Loop



## Length Abstraction Solver: Refinement Loop



## Length Abstraction Solver Summary

- Returns UNSAT when there are no length solutions left to explore
- Returns SAT when character solver returns SAT
- For a given length solution
- Learns length constraints that block character UNSAT Cores
- Works well for character queries that are only conjunctions
- Faster solving
- Better learning


Arrangement Solver Extension

## Existing Arrangement Solver

Basic Idea:

- split equations into simpler ones (arrangements)
- until their satisfiability is "easily decided" (solved form)
- backtrack on conflicts


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$$
X Y \text { Y M N }
$$

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## Arrangement Solver Extension

## Basic Idea:

- check some arrangements with length-abstraction solver (LAS)
- helps block infinite loops
- helps model generation



## Selection Architecture

## Selection Architecture



## Selection Architecture: Conjunctive Fragment



## Conjunctive Fragment Static Analysis

$$
\begin{aligned}
& F \quad::=\text { Atom }|\quad F \wedge F| \neg G\left|A_{\text {int }} \vee A_{\text {int }}\right| \neg A_{\text {int }} \\
& G \quad::=G \vee G \quad \mid \quad \neg F \\
& \text { Atom }::=t_{\text {str }}=t_{\text {str }} \quad\left|\quad A_{\text {int }} \quad\right| \quad A_{\text {ext }} \\
& A_{\text {int }} \quad::=t_{i n t}=t_{\text {int }} \quad \mid \quad t_{\text {int }}<t_{\text {int }} \\
& A_{\text {ext }}::=\operatorname{prefix}\left(t_{s t r}, t_{s t r}\right) \quad \mid \quad \text { suffix }\left(t_{s t r}, t_{s t r}\right) \\
& t_{i n t} \quad::=m|v| \operatorname{len}\left(t_{\text {str }}\right)\left|t_{\text {int }}+t_{\text {int }}\right| m \cdot t_{i n t} \\
& \text { where } m \in \text { Con }_{\text {int }} \& v \in \text { Var }_{\text {int }}
\end{aligned}
$$

$$
\begin{aligned}
& \text { where } s \in \text { Con }_{\text {str }} \& v \in \text { Var }_{\text {str }}
\end{aligned}
$$

## Conjunctive Fragment Static Analysis

Theorem 2 (Conjunctive Fragment). Let $L$ be the language generated by the grammar in Figure 3. If $\varphi \in L$ is an input query, then LAS will always call ReduceToBV such that it produces a conjunction of bit-vector equations.

Evaluation

## Experimental Setup

All experiments were performed on a server running

- Ubuntu I8.04.4 LTS
- with two AMD EPYC 7742 processors and
- 2TB RAM
- using the ZaligVinder benchmarking framework.
- The timeout for solving an instance was set at 20 seconds.


## Overall Evaluation



10,000 20,000 30,000 40,000 50,000 60,000 70,000 80,000 90,000100,00010,00020,000
Solved instances

## Length Abstraction Solver Analysis

- LAS solves more queries per second compared to the arrangement solver in the conjunctive fragment (I $128.5 \%$ )
- than it does outside the conjunctive fragment (904.6\%)
- High percentages due to dynamic difficulty estimation
- Really good at cutting off LAS


## Arrangement Solver Analysis

- Without the extension, the arrangement solver
- solves 724I7 instances
- in 423949.163 seconds
- With the extension, the arrangement solver
- solves 10740 I instances (148.3\% of the queries without)
- in 262047.893 seconds ( $61.8 \%$ of the time without)


## Thank You!

z3str4.github.io


