

Short Course

Robust Optimization and Machine Learning

Lecture 5: Robust Optimization

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“Nominal” optimization problem:

$$\min_x f_0(x) : f_i(x) \leq 0, \quad i = 1, \dots, m$$

f_0, f_i 's are **convex**.

- ▶ Includes many problems arising in decision making, statistics.
- ▶ Efficient (polynomial-time) algorithms.
- ▶ Convex relaxations for non-convex problems.

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Uncertainties are a pain!!

In practice, problem data is **uncertain**:

- ▶ *Estimation* errors affect problem parameters.
- ▶ *Implementation* errors affect the decision taken.

Uncertainties often lead to highly unstable solutions, or much degraded realized performance.

These problems are compounded in problems with multiple decision periods.

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Robust counterpart

“Nominal” optimization problem:

$$\min_x f_0(x) : f_i(x) \leq 0, \quad i = 1, \dots, m.$$

Robust counterpart:

$$\min_x \max_{u \in \mathcal{U}} f_0(x, u) : \forall u \in \mathcal{U}, f_i(x, u) \leq 0, \quad i = 1, \dots, m$$

- ▶ functions f_i now depend on a second variable u , the “uncertainty”, which is constrained to lie in given set \mathcal{U} .
- ▶ Inherits convexity from nominal. Very tractable in some practically relevant cases.
- ▶ Complexity is high in general, but there are systematic ways to get relaxations.

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Robust chance counterpart

(Assume for simplicity there are no constraints)

$$\min_x \max_{p \in \mathcal{P}} \mathbf{E}_p f_0(x, u).$$

- ▶ Uncertainty is now random, obeys distribution p .
- ▶ Distribution p is only known to belong to a class \mathcal{P} (e.g., unimodal, given first and second moments).
- ▶ Complexity is high in general, but there are systematic ways to get relaxations.
- ▶ Rich variety of related models, including Value-at-Risk constraints.

In this lecture: our main goal is to introduce some important concepts in robust optimization, e.g. robust counterparts, affine recourse, distributional robustness.

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Nominal problem:

$$\min_x c^T x : a_i^T x \leq b_i, \quad i = 1, \dots, m.$$

We assume that $a_i = \hat{a}_i + \rho u_i$, where

- ▶ \hat{a}_i 's are the nominal coefficients.
- ▶ u_i 's are the uncertain vectors, with $u_i \in \mathcal{U}_i$ but otherwise unknown.
- ▶ $\rho \geq 0$ is a measure of uncertainty.

Assumption that uncertainties affect each constraint independently is done without loss of generality.

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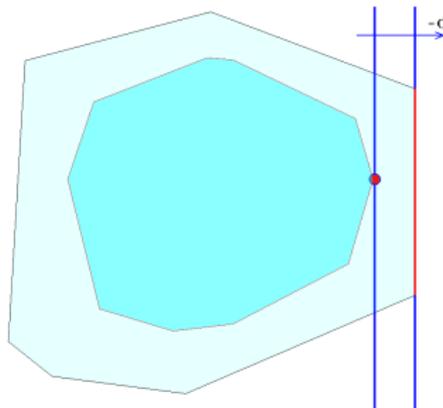
Robust counterpart

Robust counterpart:

$$\min_x c^T x : \forall u_i \in \mathcal{U}_i, (\hat{a}_i + \rho u_i)^T x \leq b_i, \quad i = 1, \dots, m.$$

Solution may be hard, but becomes easy when:

- ▶ \mathcal{U}_i are polytopic, given by their vertices (“scenarios”);
- ▶ \mathcal{U}_i 's are “simple” sets such as ellipsoids, boxes, LMI sets, etc.
- ▶ Complexity governed by the support functions of sets \mathcal{U}_i .



Robust LP with ellipsoidal uncertainty.

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Basic idea

Nominal LP:

$$\min_x c^T x : Ax \leq b.$$

We assume that A, b are affected by uncertainty in affine fashion. We assume uncertainty is available to “known by” some decision variables (*e.g.*, price revealed as time unfolds).

We seek an affinely adjusted robust solution (*i.e.*, a linear feedback).

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Nominal LP:

$$\min_x c^T x : Ax \leq b.$$

Assume that

- ▶ Right-hand side b is subject to uncertainty, $b(u) = \hat{b} + Bu$ with $u \in \mathcal{U}$.
- ▶ Decision variable can depend on (parts of) u : $x(u) = \hat{x} + Xu$.

Model information on u available to $x(\cdot)$ as $X \in \mathcal{X}$.

Affinely Adjustable Robust counterpart (AARC):

$$\min_{\hat{x}, X \in \mathcal{X}} \max_{u \in \mathcal{U}} c^T x(u) : \forall u \in \mathcal{U}, Ax(u) \leq b(u).$$

Above is tractable (provided \mathcal{U} is).

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Assume $\mathcal{U} = [-\rho, \rho]^m$, we obtain the AARC

$$\min_{\hat{x}, X \in \mathcal{X}} c^T \hat{x} - \rho \|c^T X\|_1 : A\hat{x} + \rho s \leq \hat{b}, \quad s_i \geq \|e_i^T (AX - B)\|_1, \quad i = 1, \dots, m.$$

We recover the “pure” robust counterpart with $X = 0$.

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Case with coefficient uncertainty

Approach can be extended to cases when A , c are also uncertain.

- ▶ AARC is usually not tractable.
- ▶ Efficient approximations via SDP.

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Chance constraints

Simple case

Consider an LP, and assume one of the constraints is $a^T x \leq b$, where $x \in \mathbf{R}^n$ is the decision variable.

If a is random, we can often deal with the chance constraint

$$\text{Prob} \left\{ a^T x \leq b \right\} \geq 1 - \epsilon$$

easily. For example, if a is Gaussian with mean \hat{a} and covariance matrix Γ , above is equivalent to

$$\hat{a}^T x + \kappa(\epsilon) \|\Gamma^{1/2} x\|_2 \leq b,$$

where $\kappa(\cdot)$ is a known function that is positive when $\epsilon < 0.5$.

More complicated chance constraints

Often, the random variable enters quadratically in the constraint. This happens for example when x includes affine recourse, and a depends linearly on some random variables.

We are led to consider

$$\text{Prob} \left\{ (u, 1)^T W(u, 1) > 0 \right\} \leq \epsilon$$

where W depends *affinely* on the decision variables. Above is hard, even in the Gaussian case.

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Distributional robustness

Consider instead

$$\sup_{p \in \mathcal{P}} \mathbf{Prob}_p \left\{ (u, 1)^T W(u, 1) > 0 \right\} \leq \epsilon$$

where the sup is taken with respect to all distributions p in a specific class \mathcal{P} , specifying e.g.:

- ▶ Moments.
- ▶ Symmetry, unimodality.

Fact: when \mathcal{P} is the set of distributions having zero mean and unit covariance, the condition $\sup_{p \in \mathcal{P}} P_{wc} \leq \epsilon$ is equivalent to the LMI in M, v :

$$\mathbf{Tr} M \leq \epsilon v, \quad M \succeq 0, \quad M \succeq vJ + W,$$

where J is all zero but a 1 in the bottom-right entry.

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Example

Transaction costs In many financial decision problems, the transaction costs can be modeled with

$$T(x, u) = \|A(x)u + b(x)\|_1,$$

for appropriate affine $A(\cdot), b(\cdot)$.

Example:

$$\sum_{t=1}^T |x_{t+1} - x_t|$$

with decision variable x_t an affine function of u .

This leads to consider quantities such as

$$\max_{u \sim (0, I)} \mathbf{E} T(x, u)$$

where $u \sim (0, I)$ refers to distributions with zero mean and unit covariance matrices.

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A useful result

For given $m \times d$ matrix A and d -vector b , define

$$\phi := \max_{u \sim (0, I)} \mathbf{E} \|Au + b\|_1$$

Let a_i denote the i -th row of A ($1 \leq i \leq m$). Then

$$\frac{2}{\pi} \psi \leq \phi \leq \psi,$$

where

$$\psi := \sum_{i=1}^m \left\| \begin{pmatrix} a_i \\ b_i \end{pmatrix} \right\|_2.$$

Note: ψ is convex in A, b , which allows to minimize it if A, b are affine in the decision variables.

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Dynamic programming

- ▶ Finite-state, discrete-time Markov decision process.
- ▶ Finite-horizon control problem: minimize expected cost.
- ▶ $a \in \mathcal{A}$ denote actions, $s \in \mathcal{S}$ states, and $c_t(s, a)$ the cost for action a in state s at time t .

Bellman recursion (value iteration):

$$v_t(s) = \min_{a \in \mathcal{A}} c_t(s, a) + p_t(a)^T v_{t+1}, \quad s \in \mathcal{S}$$

with $p_t(a)$ the transition probabilities at time t under action a .

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Uncertainty on transition matrix

We assume that *at each stage*, “nature” picks a transition probability vector $p_t(a)$ in a given set $\mathcal{P}_t(a)$.

Robust counterpart: the robust control problem, with “nature” the adversary.

Robust Bellman recursion:

$$v_t(s) = \min_{a \in \mathcal{A}} c_t(s, a) + \max_{p \in \mathcal{P}_t(a)} p^T v_{t+1}, \quad s \in \mathcal{S}.$$

For a wide variety of sets $\mathcal{P}_t(a)$, inner problem very easy to solve.

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Entropy uncertainty model

A natural way to model uncertainty in the transition matrices involves relative entropy bounds

$$\mathcal{P} = \left\{ p \geq 0 : \sum_j p_j \log \frac{p_j}{q_j} \leq \beta, \sum_j p_j = 1 \right\}.$$

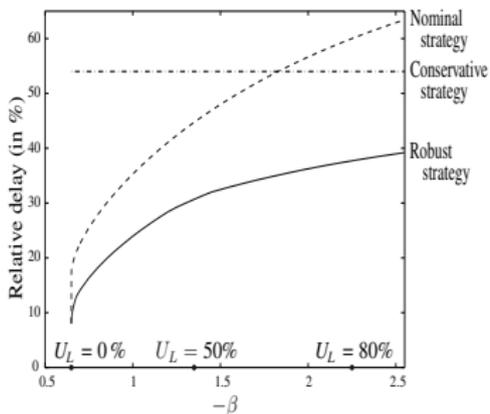
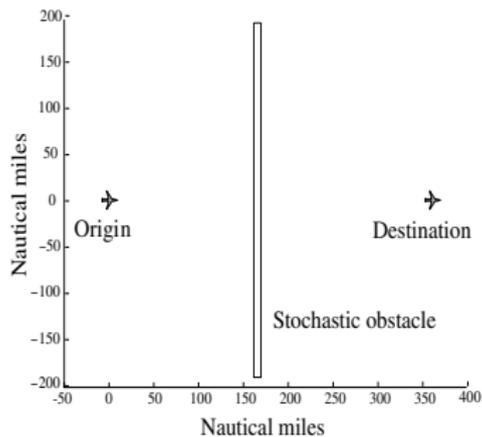
where $\beta > 0$ is a measure of uncertainty, and q is the nominal distribution.

The corresponding inner problem can be solved in $O(n)$ via bisection.

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Robust path planning



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