Math 55 - Spring 2004 - Lecture notes #20 - Apr 8 (Tuesday) random variables Goals for today: expection (average, mean) of random variables DEF: Let S be the sample space of a given experiment, with probability function P. A _random variable_ is a function f:S -> Reals. EX: Flip a biased coin once, $S1 = \{H,T\}$, P1(H) = p, $f1(x) = \{1 \text{ if } x=H, -1 \text{ if } x=T\}$ f1 = amount you win (if f1>0) (or lose, if f1<0) if you bet \$1 on H.EX: Flip a biased coin n times, S2 = {all sequences of H, T of length n} ASK&WAIT: What is P2(x), if x has i Heads? f2(x) = #H - #T = #H - (n-#H) = 2*#H - nf2 = amount you win (or lose) if you bet \$1 on H on each flip EX: Let S3 = result of rolling a die once, P3(any face) = 1/6 Let f3(outcome e) = value on top of die (an integer from 1 to 6) EX: Let S4 = result of rolling a pair of red and blue dice 24 times $= \{ ((1,1), (1,1), \dots, (1,1)), \dots, ((6,6), (6,6), \dots, (6,6)) \}$ <----> 24 times ----> <----> 24 times ----> What is P4(any outcome x in S4)? ASK&WAIT: Let f4(outcome x in S4) = { +1 if a pair of sixes appears in e } } { -1 otherwise We can interpret f4 as the amount of money we win (or lose) by betting on getting a pair of sixes EX: $S5 = \{US \text{ population}\}, P5(person x in S) = 1/|S5|,$ Let f5(person x in S) = { +1 if x has a particular disease } { 0 if x does not } EX: S6 = {all permutations of 1 to n}, P6(any permutation) = 1/n!f6(any permutation x) = time for your sorting algorithm to sort x EX: Suppose you flip a fair coin, and win \$1 if it comes up H, lose \$1 if it comes up T ASK&WAIT: What is the "average" amount you expect to win after N flips? DEF: Given S, P and random variable f, the _Expected Value_ (also called Mean or Average) of f is $E(f) = sum_{all outcomes x in S} P(x)*f(x)$

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This is the "average" value of f ones gets if one repeats the experiment a great number of times. EX: With S1, P1, f1 as before, E(f1) = (+1)*(p) + (-1)*(1-p) = 2*p-1= 0 if coin fair (p=1/2)Imagine betting \$1 on getting H. Then E(f1) is the amount you expect to win (if E(f1)>0) or lose (E(f1)<0) on the bet. If E(f1)=0, you break even EX: With S2, P2 and f2 as before, If we flip a coin N times, we expect E(f2) to be the amount we win betting \$1 on flip to get H; and intuititively this should be N*E(f1) = N*(2*p-1)Formally, we get $E(f2) = sum_{sequences x of n Hs and Ts} f2(x)*P2(x)$ = sum_{sequences x of n Hs and Ts} (#H in x)*P2(x) looks complicated, but later we will see that our intuition was right, and there is an easier way to do it that matches our intuitive approach EX: With S3, P3, f3 as before, $E(f3) = (1/6)*1 + (1/6)*2 + \dots + (1/6)*6 = 21/6 = 7/2$ EX: With S5, P5, f5 as before, $E(f5) = sum_{persons x} f5(x)*P5(x)$ = sum_{sick persons x} f5(x)*P5(x) + sum_{healthy persons x} f5(x)*P5(x)= sum_{sick persons x} 1*(1/|S5|) + sum_{healthy persons x} 0*(1/|S5|)= P(random person is sick) EX: S6, P6, f6 as before, then E(f6) = average time for your algorithm to sort EX: With S4, P4, f4, seem like you need to sum over all 6⁴⁸ sequences, We need a simpler way: DEF: $P(f=r) = sum_{all outcomes x in S such that f(x)=r} P(x)$ EX: With S1, P1 and f1 as before P1(f1=1) = P1(H) = p, P1(f1=-1) = P1(T) = 1-pEX: With S2, P2 and f2 as before ASK&WAIT: What is P2(f2=i)? EX: With S3, P3 and f3 as before P3(f3=k) = 1/6 for k=1,2,...,6 and P3(f3=k)=0 otherwise

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EX: With S4, P4 and f4 as before,
      P4(f4=1) = sum_{all outcomes x in which a pair of sixes appears} P4(x)
               = P4(a pair of sixes appears)
ASK&&WAIT: What is P4(f4=-1)?
EX: With S5, P5 and f5 as above,
ASK&WAIT: what is P5(f5=1)? P5(f5=0)?
 Thm: E(f) = sum_{numbers r in range of f} r*P(f=r)
   Proof: Write down proof for S finite, but same for S countable
           Let {r1,r2,...,rk} be numbers in range of f, and write
           S = S1 U S2 U \dots U Sk where
           Si = {outcomes x in S such that f(x)=ri}
           and so P(Si) = P(f=ri)
           Note that all Si are pairwise disjoint, so we can write
           E(f) = sum_{x in S} f(x)*P(x)
                = sum_{x in S1} f(x)*P(x) + sum_{x in S2} f(x)*P(x)
                  + ... + sum_{x in Sk} f(x) * P(x)
                = sum_{x in S1} r1*P(x) + sum_{x in S2} r2*P(x)
                  + ... + sum_{x in Sk} rk*P(x)
           Look at one term:
           sum_{x in Si} ri*P(x) = ri * sum_{x in Si} P(x)
                                 = ri * P(Si)
                                 = ri * P(f=ri)
           so E(f) = r1*P(f=r1) + r2*P(f=r2) + ... + rk*P(f=r3)
                   = sum_{number r in range of f} r*P(f=r)
           as desired.
EX: With S3, P3 and f3 as above,
     E(f3) = sum_{k=1} to 6 k*P(f=k) = sum_{k=1} to 6 k*(1/6) = 7/2 as before
EX: With S4, P4, f4 as above,
     E(f4) is the average amount one wins (if E(f4)>0) or loses (if E(f4)<0)
     every time one plays.
     E(f4) = sum_{numbers r in range of f} r*P(f4=r)
           = +1*P4(getting pair of sixes) + (-1)*P4(not getting pair of sixes)
           = P4(getting pair of sixes) - P4(not getting pair of sixes)
ASK&WAIT:
              What is P4(not getting pair of sixes)?
     P4(getting pair of sixes) = 1 - P4(not getting pait of sixes)
                               \sim 1-.5086 = .4914
     and E(f4) = .4914 - .5086 = -.0172, so you lose in the long run
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Note: In 1654 the gambler Gombaud asked Fermat and Pascal whether
           this was a good bet, inadvertently starting the field of
           probability theory
     Note: If we do 25 rolls instead of 24,
           P4(not getting a pair of sixes) drops to (35/36)<sup>25</sup> ~ .4945
           P4(getting pair of sixes) grows to .5055, so it is a good bet.
 EX: Let S5, P5, f5 be as above. Then
     E(f5) = (+1)*P(f5=1) + O*P(f5=0)
           = P(f5=1) = P(person sick)
     This is a special case of the following lemma:
 Lemma: Let S be a sample space, E subset S any event, and
        f(x) = \{1 \text{ if } x \text{ in } E\}
                                  }
                \{0 \text{ if } x \text{ not in } E\}
        Then E(f) = P(E)
ASK&WAIT: proof?
EX: S2, P2, f2 as above:
    E(f2) = expected win betting $1 on a coin N times
          = sum_{i=0 to N} i*P2(getting i heads)
          = sum_{i=0 to N} i*C(N,i)*p^i^(1-p)^(N-i)
    still isn't simple, so need a new idea:
 Thm: Let S and P be a sample space and probability function, and
      let f and g be two random variables. Then
            E(f+g) = E(f) + E(g)
      Proof: Let h=f+g be a new random variable.
            Then E(h) = sum_{outcomes x in S} h(x)*P(x)
                       = sum_{outcomes x in S} (f(x)+g(x))*P(x)
                       = sum_{x} f(x)*P(x) + sum_{x} g(x)*P(x)
                       = E(f)
                                            + E(g)
  Corollary: Let S and P be as above, and h = f1 + f2 + \ldots + fn
      Then E(h) = E(f1) + E(f2) + ... + E(fn)
  EX: Let S2, P2, f2 be as before. Then we can write
      f2 = g1 + g2 + ... + gN where
      gi(x) = \{ +1 \text{ if } i-th flip = H \}
               \{-1 \text{ if } i-th flip = T \}
      and E(f2) = E(g1) + E(g2) + ... + E(gN)
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For any i E(gi) = (+1)*P(H) + (-1)*P(T) = p - (1-p) = 2*p-1so E(f2) = N*(2*p-1)which matches our original intuition about making N independent bets in a row (whew!)

EX: Let S4, P4, and f4 be as before. Suppose you also make the side bet that you win 2 if at least 8 fives come up, and lose 2.5 if fewer than 8 fives come up. Is this joint bet worth making? Answer: Let $g(x) = \{ +2 \text{ if at least 8 fives come up in x } \}$ $\{ -2.5 \text{ if at most 7 fives come up in x } \}$ P(g=+2) = P(at least 8 fives) $= \text{sum}_{\{i=8 \text{ to } 48\}} C(48,i) * (1/6)^i * (5/6)^(48-i)$ $\tilde{.}55992$ P(g=-2.5) = P(at most 7 fives)= 1 - P(at least 8 fives)= 1 - P(at least 8 fives)= 1 - .55992 = .44008E(g) ~ +2*.55992 - 2.5*.44008 ~ .0196Then the value of the joint bet f4+g is E(f4+g) = E(f4)+E(g) ~ -.0172+.0196 = .0024and being positive, is worth making.

EX: Suppose you shoot at a target, and miss it with probability p each time you try. What is the expected number of times you have to try before getting a hit? $S = \{ H, MH, MMH, MMH, \ldots \}$ $P(MM...MH) = p^{H}M * (1-p)$ f(MM...MH) = #shots = #M + 1 We want $E(f) = sum_{m=0}^{infinity (m+1)*p^m*(1-p)}$ $sum_{m=0}^{infinity p^m = 1/(1-p)}$ Recal so d/dp (sum_{m=0}^infinity p^m) = d/dp (1/(1-p)) sum_{m=0}^infinity m*p*{m-1} = 1/(1-p)^2 or sum_{m=0}^infinity m*p^m*(1-p) = p/(1-p) or so sum_{m=0}^infinity $(m+1)*p^m*(1-p) =$ p/(1-p) + (1-p)/(1-p) = 1/(1-p)so E(f) = 1/P(hit)So if P(M)=.99, you need to take 1/(1-.99) = 100 shots on average to hit

EX: Suppose homework from 350 students is collected, graded, randomly shuffled, and handed back. What is the expected number of students who get their own homework back? S = {permutations of 1 to 350}, P(any permutation) = 1/350! ~ 8e-741 f(permutation x) = #homeworks returned to right students,

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We want E(f)
     Let fi(x) = { 1 if student i gets right homework back }
                  { 0 otherwise }
     Then f(x) = f1(x) + f2(x) + \ldots + f350(x)
     and E(f) = E(f1) + ... + E(f350)
     Now E(fi) = P(student i gets right homework)
                = (# permutations where student i gets right homework)/350!
               = (# permutations of other 349 homeworks)/350!
               = 349! / 350! = 1/350
     so E(f) = 350*(1/350) = 1
     Result would be true for any number of students!
 EX: Recall definition of independent sets:
     P(A \text{ inter } B) = P(A) * P(B); intuitively, this means that knowing whether
     or not you are a member of A tells you nothing about whether you are
     a member of B
 EX: S = \{ 2 \text{ coin flips} \} = \{ HH, HT, TH, TT \} \text{ with } P(H)=p, P(T)=q=1-p
     A = \{ HH, HT \}, B = \{ HH, TH \}
     Then P(A \text{ inter } B) = P(HH) = p^2 = P(A)*P(B)
     Let f1(e) = \{ 1 \text{ if first coin H} \\ f2(e) = \{ 1 \text{ if second coin H} \}
                  { 0 otherwise
                                                    { 0 otherwise
     Then P(A \text{ inter } B) = P(f1=1 \text{ and } f2=1) = P(A)*P(B) = P(f1=1)*P(f2=1)
     Can also check that both A and complement(A) are independent with
     B and complement(B), i.e.
         P(f1=r1 and f2=r2) = P(f1=r1)*P(f2=r2) for r1, r2 = 0, 1
     (this was homework due yesterday!)
DEF Let f, g be random variables. Then we call
        f and g _independent_ if for all values of r and s
        P({x: f(x)=r and g(x)=s}) = P({x: f(x)=r}) * P({x: g(x)=s})
    This generalizes situation where f,g= 1 or 0 only. It still means
    intuitively that knowing about the value of f says nothing about
    the value of g. In this case, the result P(A \text{ inter } B)=P(A)*P(B)
    generalizes as follows:
 Thm 1: If f and g are independent, then E(f*g)=E(f)*E(g)
  proof: E(f*g) = sum_{x in S} f(x)*g(x)*P(x)
                 = sum_{pairs (r,s)} sum_{x such that f(x)=r and g(x)=s}
                                 f(x)*g(x)*P(x)
                      because we do the same sum over x, just grouping
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as desired.

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EX: Flip a biased coin 10 times
    S = {all sequences of 10 Hs and Ts}, P(x) = 1/2^{10}
    Let f = #H - #T in first 5 flips
    Let g = "turn last 5 flips into binary number b, via 1=H and 0=T"
    Then f and g are independent, because they depend on independent
      events (first 5 flips vs last 5 flips) so E(f*g)=E(f)*E(g)
    What are E(f)? E(g)? E(f*g)?
        E(f) = E(\#H) - E(\#T) = 5*p - 5*(1-p) = 10*p-5
        To compute E(g), let b4 = { 1 if coin 6 = H, 0 otherwise }
                               b3 = \{ 1 \text{ if } coin 7 = H, 0 \text{ otherwise } \}
                               . . .
                               b0 = \{ 1 \text{ if coin } 10 = H, 0 \text{ otherwise } \}
        Then b = b4*2^4 + b3*2^3 + b2*2^2 + b1*2 + b0
        so E(g) = E(b) = E(b4)*2^4 + \ldots + E(b0)
                         = p*2^4 + ...
                                           + p
                         = p*31
        Finally, E(f*g)=E(f)*E(g)=(10*p-5)*p*31
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