

Math 55 - Spring 04 - Lecture notes # 1 - Jan 20 (Tuesday)

Name, class, URL ([www.cs.berkeley.edu/~demmel/ma55](http://www.cs.berkeley.edu/~demmel/ma55)) on board  
Head TA Mike West speaks on bureaucracy  
Advertise CS 70 (T Th 2-3:30) as an "honors" alternative to Ma55  
All material on web page and at Copy Central Northside  
Read Course Outline on web page for course rules and grading policies  
You are responsible for reading this and knowing the rules!

Class outline:

Basics (Chap 1) - common language for rest of course

logic (basis of proofs,  
logical operations in programs (CS61)  
hardware design (CS150))

sets and functions (same proof of 3 results:

- (1) can show that can't write a program to  
implement every possible function,  
because there are more functions than programs;
- (2) can show it is impossible to to  
write a debugger that finds "infinite loops"  
in other programs (CS 172))
- (3) there are more real numbers than rational numbers  
(even though there are infinitely many of both

Used in most of math curriculum

Integer algorithms (Chap 2)

Big-0 notation (measure approximately how fast a function  $f(n)$   
increases as  $n$  increases)

Using Big-0 to analyze speed of algorithms (used in CS61, CS170)

Prime numbers, GCD (Greatest common divisors),

"modular arithmetic" ( $a \bmod b$ ), Chinese remainder theorem

Applications: generating random numbers

integer with "bignums"

cryptography (eg RSA used in Netscape)

(Ma 113, CS 170)

More proof techniques (Chap 3)

induction (knowing when a theorem, or program, is correct)

(CS 17x, most of math curriculum)

Probability theory and counting (Chaps 4, 5, Lenstra's notes)

analyzing, designing algorithms (CS 170, CS 162, CS 188...)

(sometimes an algorithm that uses "random numbers" is faster than one that does not)  
 how many ways to pair up n students into teams of 2 or 3 or ...  
 how to gamble in Reno  
 algorithms that use random numbers  
 EX suppose you build a web search engine  
 (or web page to sell books or ...)  
 you buy 100 PCs and put them in a room  
 when a request comes in, you pick a PC at random, and send it to that PC  
 what is the average waiting time to service a request?

Read Chapter 1 of book, homework due Wednesday Jan 28  
 at the beginning of section

Do problems

sec 1.1: 6, 16(a,c,e,g), 40, 42, 48, 60  
 1.2: 4a, 8(a,b), 10 (for 8a and 8b), 20, 38 (assume results of 36, 37)  
 1.3: 8, 14, 26(a,c,e), 56(a,b)  
 1.4: 4, 8(a,c,e), 28(a,c,e,g,i)  
 1.5: 6, 8(a,c,e), 22, 52

Begin course:

DEF: Proposition is a statement that must be either true (T) or false (F)

EG:  $2+2 = 4$  (Y)

$2+2 = 3$  (Y)

Is  $2+2=4$ ? (N)

Let  $x = 3$ . (N)

$x+1 = 2$  (N, unless we know the value of x)

Every even number  $> 2$  is the sum of two primes. (Y)

Notation: propositions denoted p,q,r ...

DEF: not p , p or q (disjunction), p and q (conjunction),  
 p xor q (exclusive or)

Truth tables

DEF: a compound proposition is formed by simple propositions connected by logical operators and parentheses;

EG: (p and q) or (r and s)

programming example below

DEF:  $p \rightarrow q$ , "if p, then q", "p implies q", "p sufficient for q",  
 "q necessary for p";  $p \rightarrow q$  same as  $\text{not}(p) \text{ or } q$ ;

EG: if it is sunny, then we go to the beach

p = it is sunny, q = we go to the beach  
(T unless it is sunny (p) and we don't go (not q))

Truth table

ASK&WAIT EG: if (today is Friday), then (2+3 = 6);  
on which days is this true?

p = (today is Friday), q = (2+3=6)  
T every day except Friday

ASK&WAIT EG: if (1+1=3) then (2+2=4)

Note that p and q need not have anything to do with one another  
in order to form the expression  $p \rightarrow q$

DEF: The converse of  $p \rightarrow q$  is  $q \rightarrow p$ ; they need not be true  
at the same time

ASK&WAIT EG: If I am exactly 18 years old then I am allowed to vote;  
converse is

if I am allowed to vote, then I am exactly 18 years old  
note that converse not necessarily true if original is

DEF: The contrapositive of  $p \rightarrow q$  is  $\text{not } q \rightarrow \text{not } p$ ;  
these are the same

ASK&WAIT EG: If I am exactly 18 years old then I am allowed to vote;  
converse is

if I am not allowed to vote, then  
I am not exactly 18 years old  
contrapositive is true

Truth table proof of  $\text{prop} \Leftrightarrow \text{contrapositive}$

DEF:  $p \leftrightarrow q$  "p if and only q", " $p \rightarrow q$  and  $q \rightarrow p$ ", biconditional

Application to programming (C or C++ syntax)

if (compound proposition, or logical expression) then do something...  
if ((n>0) && (m<k)) then do something ...  
(n>0) and (m<k) are propositions whose value (T or F)  
can be evaluated when you get to this line in the program,  
and && means "and", || means "or", etc.

in computer 1 represents true and 0 false (representation details  
in CS61C, CS150)

Application to Web search

In [www.google.com](http://www.google.com) you can type into the Advanced Search form:

With all the words: guano

and With at least one of the words: geometry tickle

Same as "find all web pages W for which the proposition is true"

(W contains guano) and  
 ((W contains geometry) or (W contains tickle))

ASK&WAIT: what web pages do you get?

EX: Geometry Learning center page on disease histoplasmosis,  
 caused by fungus growing in bat guano,  
 contains guano and geometry, not tickle

EX: Magazine article about "tickle down economics"  
 of Christmas toys and a town in Florida  
 that smells like alligator guano  
 contains guano and tickle, not geometry

EX: poem in on-line poetry magazine Lynx  
 contains all 3

Now let's try google with  
 With all the words: guano  
 and With at least one of the words: geometry tickle  
 and Without the words the

ASK&WAIT: Same as "find all web pages W for which the proposition is true"  
 (W contains guano) and  
 ((W contains geometry) or (W contains tickle)) and  
 (not(W contains the))

ASK&WAIT: what web pages do you get?

EX: "G" page of an Esperanto dictionary  
 contains guano, geometry, not tickle or the

.....

Next Goal: simplifying compound propositions or replacing  
 an expression depending on p, q, r, ... combined with  
 "and", "or", etc. with another "equivalent" and "simpler"  
 one that has the same truth/false value for all values of  
 p, q, r, ... (just like when you learned algebra in high school;  
 algebra with variables that can only have values True and False  
 is called "Boolean Algebra")

What is "equivalence"?

DEF:  $p \leftrightarrow q$  means  $(p \rightarrow q)$  and  $(q \rightarrow p)$ ; true if p and q are  
 both true or both false (show truth table)

DEF: propositions p and q are called logically equivalent if  
 $p \leftrightarrow q$  is always true (a tautology); write  $p \Leftrightarrow q$ ;

EG:  $p \Leftrightarrow \text{not}(\text{not}(p))$

EG:  $p \text{ or } \text{not}(p) \Leftrightarrow \text{True}$  ( $p \text{ or } \text{not}(p)$  a "tautology")

EG:  $p \text{ and } \text{not}(p) \Leftrightarrow \text{False}$  ( $p \text{ and } \text{not}(p)$  a "contradiction")

When is q "simpler" than p? Roughly, if q is a "smaller" compound expression, with fewer operations like and, or etc

Motivation:

- 1) Proofs: if the compound prop. p equivalent to True (p a tautology), then you have proved a theorem
- 2) In programming, simple expressions (in "if" statements) are easier to understand and debug, faster to evaluate, since each "and" or "or" costs an operation
- 3) In computer design: A computer stores and computes with "bits" which are either 1 (true) or 0 (false), represented by electrical signals. The computations are done with little pieces of hardware that, say, take two input bits and compute their "and" to get an output bit. Each little piece of hardware corresponds to one of the operations "and", "or", etc. that we have talked about, and so computer hardware is just evaluating compound propositions. A "simpler" compound proposition needs fewer pieces of hardware to evaluate it, which is smaller, cheaper, faster... See also CS150, Chap 10, other "CAD" courses in EECS, and whole industries.

Rules we can use to simplify compound propositions:

ASK&WAIT EG: True or q  $\Leftrightarrow$  True (domination)  
ASK&WAIT EG: (I am a student at Stanford) or (2+2=4)  $\Leftrightarrow$  True  
EG: False and q  $\Leftrightarrow$  False (domination)  
EG: p or q  $\Leftrightarrow$  q or p (commutativity)  
EG: ( p or q ) or r  $\Leftrightarrow$  p or ( q or r ) (associativity)  
EG: DeMorgan's laws  
not(p or q)  $\Leftrightarrow$  (not p) and (not q)  
not(p and q)  $\Leftrightarrow$  (not p) or (not q)  
proof: Truth table  
ASK&WAIT EG: Prove that not(p  $\rightarrow$  q)  $\rightarrow$  p a tautology, i.e.  $\Leftrightarrow$  True  
not(p  $\rightarrow$  q)  $\rightarrow$  p  
not(not(p) or q)  $\rightarrow$  p Def  $\rightarrow$   
not(not(not(p) or q)) or p Def  $\rightarrow$   
(not(p) or q) or p Double negative  
not(p) or (q or p) associativity  
not(p) or (p or q) commutativity  
(not(p) or p) or q associativity  
(True ) or q not(p) or p  $\Leftrightarrow$  True  
True true or q  $\Leftrightarrow$  true

Like musical instrument: practice

Also: Truth Table (Ref: EE290H - logic verification, 150 vars)  
EG:  $p \text{ or } (q \text{ and } r) \Leftrightarrow (p \text{ or } q) \text{ and } (p \text{ or } r)$  distributivity  
EG:  $p \text{ and } (q \text{ or } r) \Leftrightarrow (p \text{ and } q) \text{ or } (p \text{ and } r)$  distributivity  
(notice that names of rules -- distributivity etc -- are  
same as rules for regular algebra with numbers, since the rules  
are analogous)

How does one interpret an expression like not p or q?  
as (not p) or q, or as not(p or q), which are quite different  
The precedence for operators in absence of parentheses is  
Highest to Lowest: not, ( and, or ), ( -> , <-> )  
If there is ambiguity, I will use parentheses