

Chap 6+7 : Iterative Methods

Part 3: Splitting Methods

Given initial guess for $Ax=b$, x_0
cheaply compute $x_1, x_2, \dots \rightarrow x$

Def: A splitting of $A = M - K$, M nonsingular

Iterative Method: $Ax = Mx - Kx = b$

$$\Rightarrow Mx = Kx + b \Rightarrow \text{solve } Mx_{i+1} = Kx_i + b$$

$$\text{for } x_{i+1} \Rightarrow x_{i+1} = \underline{M^{-1}K}x_i + \underline{M^{-1}b} = \underline{R}x_i + \underline{c}$$

To work well need:

- ① it should converge (quickly)
- ② multiplying by R (solving $Mx_{i+1} = Kx_i + b$ for x_{i+1}) and computing $c = M^{-1}b$ should be much cheaper than solving with A

Recall operator norm: $\|A\| = \max_{x \neq 0} \frac{\|Ax\|}{\|x\|}$

Lemma: Let $\|\cdot\|$ be any operator norm:

Then if $\|R\| < 1$, then $x_{i+1} = Rx_i + c$ converges

for any x_0 .

Proof:
$$\begin{aligned} x_{i+1} &= Rx_i + c \\ - [x &= Rx + c] \\ \hline x_{i+1} - x &= R(x_i - x) \end{aligned} \quad e_i = x_i - x$$

$$e_{i+1} = R e_i = \dots = R^{i+1} e_0$$

$$\|e_{i+1}\| \leq \|R^{i+1}\| \cdot \|e_0\| \leq \|R\|^{i+1} \cdot \|e_0\|$$

$$\rightarrow 0 \text{ if } \|R\| < 1 \quad \text{QED}$$

Def The spectral radius of R is

$$\rho(R) = \max_i |\lambda_i(R)|$$

Lemma: For all operator norms, $\rho(R) \leq \|R\|$.

For all R , $\varepsilon > 0$, $\exists$ operator norm $\|\cdot\|_*$ such that $\|R\|_* \leq \rho(R) + \varepsilon$

Proof: To show $\rho(R) \leq \|R\|$, let $Rx = \lambda x$
 where $|\lambda| = \rho(R)$ $\|R\| \geq \frac{\|Rx\|}{\|x\|} = \frac{\|\lambda x\|}{\|x\|} = |\lambda|$
 $= \rho(R)$

To construct $\|\cdot\|_*$ use Jordan Form of R

Let $S^{-1}RS = J$ be Jordan form

$$J = \begin{bmatrix} \boxed{\lambda} & & \\ & \boxed{\lambda} & \\ & & \ddots \end{bmatrix} \quad \square = \begin{bmatrix} 1 & & \\ & \ddots & \\ & & \lambda \end{bmatrix}$$

Let $D = \text{diag}(1, \varepsilon, \varepsilon^2, \dots, \varepsilon^{n-1})$

$J_\varepsilon = D^{-1} J \cdot D$: has same diagonal as J
 each superdiagonal of J ($\neq 1$ or 0)
 is multiplied by ε

$$J_\varepsilon = \begin{bmatrix} \square & & \\ & \square & \\ & & \square \end{bmatrix} \quad \square = \begin{bmatrix} d & \varepsilon & & \\ & \ddots & \ddots & \\ & & \varepsilon & \\ & & & \lambda \end{bmatrix}$$

Construct $\|\cdot\|_*$ as follows:

$$\|x\|_* = \|(S \cdot D)^T x\|_\infty = \max \text{entry in } | \cdot |$$

(Recall Q 1.5 : shows $\|\cdot\|_*$ is a vector norm)

$$\begin{aligned} \|R\|_* &= \max_{x \neq 0} \frac{\|R x\|_*}{\|x\|_*} \\ &= \max_{x \neq 0} \frac{\|(S D)^T R x\|_\infty}{\|(S D)^T x\|_\infty} \quad y = (S D)^T x \\ &= \max_{y \neq 0} \frac{\|(S D)^T R (S D) y\|_\infty}{\|y\|_\infty} \\ &= \max_{y \neq 0} \frac{\|(D^T S^T R S D) y\|_\infty}{\|y\|_\infty} \\ &= \max_{y \neq 0} \frac{\|J_\varepsilon y\|_\infty}{\|y\|_\infty} \\ &= \|J_\varepsilon\|_\infty \leq \max_i |\lambda| + \varepsilon \\ &= \rho(R) + \varepsilon \quad \forall \varepsilon > 0 \end{aligned}$$

Thm: $x_{i+1} = R x_i + c$ converges to
 solution of $Ax = b$ for all x_0
 if and only if $\rho(R) < 1$

proof: if $\rho(R) \geq 1 \Rightarrow$ choose x_0 so
 $x_0 - x$ is an eigenvector for R
 where $|\lambda| = \rho(R) \Rightarrow x_1 - x = R(x_0 - x)$
 $= \lambda(x_0 - x)$

does not converge since $|\lambda| \geq 1$

if $\rho(R) < 1$ use Lemma to build
 and operator norm $\|R\|_\infty \leq \rho(R) + \epsilon < 1$
 $\Rightarrow \|x_1 - x\|_\infty \leq \|R\|_\infty^i \|x_0 - x\|_\infty \rightarrow 0 \quad QED$

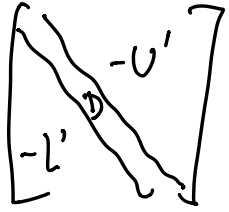
Goals: make $\rho(R)$ as small as possible
 multiplication by R easy

One choice: $M = I$ and $K = I - A$
 $\Rightarrow R = M^{-1}K = I - A$, but $\rho(R)$ may
 not be small

Good news: another choice of diagonal M
 that works in many cases = Jacobi's
 method

Describe 3 basic Splitting Methods:
 Jacobi, Gauss-Seidel (GS)
 Successive Overrelaxation (SOR)

Notation: $A = D - L' - U'$
 $= D(I - L - U)$



Jacobi:

In words: for $j = 1$ to n , pick $x_{i+1,j}$
 to solve equation j

As a loop:

for $j = 1$ to n , $\underline{x_{i+1,j}} = \frac{b_j - \sum_{k \neq j} A_{jk} x_{i,k}}{A_{jj}}$

As a splitting: $A = D - (L' + U')$
 $= M - K$

so $R_j = M^{-1}K = D^{-1}(L' + U') = L + U$

For 2D Poisson $T_N V + V T_N = h^2 F$
 $\square$

to get $\pm$ from V_i to V_{i+1} :

for $j = 1$ to n , for $k = 1$ to n

$$V_{i+1}(j,k) = (V_i(j-1,k) + V_i(j+1,k) + V_i(j,k-1) + V_i(j,k+1) + h^2 \cdot F(j,k)) / 4$$

= "average" of 4 nearest neighbors
 and $h^2 F$

Gauss-Seidel:

In words: improve on Jacobi: by using most recently updated solution values

As a loop:

$$\text{for } j=1 \text{ to } n \quad x_{i+1,j} = (b_j - \sum_{k < j} A_{jk} x_{i+1,k} - \sum_{k > j} A_{jk} x_{i,k}) / A_{jj}$$

^{updated}
 $x_{i+1,k}$

older x_i

As a splitting:

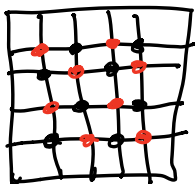
$$A = \underbrace{(D - L')}_{M} - \underbrace{U'}_{K} \text{ so}$$

$$R_{GS} = (D - L')^{-1} U' = (D(I - D^{-1}L'))^{-1} U' \\ = (I - L)^{-1} U$$

Difference from Jacobi: answer depends on order of equations

For 2D Poisson: 2 orders of equations to consider:

"Natural order": rowwise or columnwise update $V(i,j)$



Red-Black Ordering:

Color the entries in 2D mesh like a checkerboard

Property: Red nodes only have Black neighbors, and vice versa

⇒ When updating Red nodes, I can update them in any order, Given all data at Black nodes (and vice versa)

for all (j, k) nodes that are Red ($j+k$ even)

all be done
in
parallel →

$$V_{i+1}(j, k) = (V_i(j-1, k) + V_i(j+1, k) + V_i(j, k-1) + V_i(j, k+1) + h^2 F(j, k)) / 4$$

— use old data (step i)

for all (j, k) nodes that are Black ($j+k$ odd)

$$V_{i+1}(j, k) = (V_{i+1}(j-1, k) + V_{i+1}(j+1, k) + V_{i+1}(j, k-1) + V_{i+1}(j, k+1) + h^2 F(j, k)) / 4$$

— use new data (step $i+1$)

SOR:

In words: depends on parameter w :

Let ~~the~~ result x_{i+1} be a weighted combination of x_i and what GS would have computed:

$$x^{\text{SOR}}(w)_{i+1} = (1-w)x_i + w x^{\text{GS}}_{i+1}$$

If $w=1$, same as GS

If $w \leq 1$, would call this "underrelaxation"

If $w > 1$, overrelaxation.

Intuition: if moving x_i toward x^{GS}_{i+1} was a good idea, just go farther in same direction

- Question for later: how much farther, i.e. how big w ?

As a loop: for $j=1$ to n

$$x_{i+1,j} = (1-w)x_{j,j} + w \cdot (b_j - \sum_{k < j} A_{j,k} x_{i+1,k} - \sum_{k > j} A_{j,k} x_{i,k}) / A_{j,j}$$

As a Splitting: Multiply by A_{jj}

$$(D - wL')x_{i+1} = ((1-w)D + wU)x_i + wb$$

Divide by w to get

$$A = (\frac{1}{w}D - L') - (\frac{1}{w}D - D + U') = M - K$$

$$R_{SOR}(w) = M^{-1}K = (\frac{1}{w}D - L')^{-1}(\frac{1}{w}D - D + U') \\ = (I - wL')^{-1}((1-w) \cdot I + wU)$$

For 2D Poisson with Red Black Ordering:

for all Red (j, k) nodes

$$V_{i+1}(j, k) = (1-w)V_i(j, k) + \\ w \cdot (V_i(j-1, k) + V_i(j+1, k) \\ + V_i(j, k-1) + V_i(j, k+1) \\ + h^2 F(j, k)) / 4 \\ \dots \text{old (Black) data}$$

for all Black (j, k) nodes

$$V_{i+1}(j, k) = (1-w)V_i(j, k) \\ + w(V_{i+1}(j-1, k) + V_{i+1}(j+1, k) \\ + V_{i+1}(j, k-1) + V_{i+1}(j, k+1) \\ + h^2 F(j, k)) / 4 \\ \dots \text{new (Red) data}$$