

# Welcome to Ma 221! Lecture 26, Fall 24

## Conjugate Gradients (CG)

Krylov Subspace:  $\mathcal{K}_k = \text{span} \{ b, Ab, \dots, A^{k-1}b \}$

orthogonal  $Q_k = [q_1, \dots, q_k]$  orthogonal basis of  $\mathcal{K}_k$

Goal: find "best" solution to  $Ax = b$

for  $x \in \mathcal{K}_k(A, b)$

(3) Choose  $x_k$  so  $r_k = Ax_k - b \perp \mathcal{K}_k$   
 $r_k^\top Q_k = 0$

"orthogonal residual"

(4) A s.p.d : "best solution"

$$\begin{aligned} \text{minimizes } \|r_k\|_{A^*}^2 &= r_k^\top A^* r_k \\ &= \|x_k - x\|_A^2 \\ &= (x_k - x)^\top A (x_k - x) \end{aligned}$$

Lemma: CG "best" because it satisfies

(3) and (4)

$$(*) \quad x_k = Q_k (T_k)^{-1} Q_k^\top b = Q_k (T_k)^{-1} e, \|b\|_2$$

where  $T_k = \text{tridiagonal matrix}$

from Lanczos:

$$T_k = Q_k^\top A Q_k$$

$$\square = \boxed{\square} \boxed{\square}$$

## Intuition for (\*)

→ multiplying  $Q_k^T b = e_1 \|b\|_2$  projects  
 $b$  onto  $\mathcal{K}_k$

→ Multiplying by  $T_k^{-1}$  solves  
 projected problem exactly

→ Multiplying by  $Q_k$  maps  
 projection back to  $\mathbb{R}^n$

Proof of Lemma: drop subscript  $k$ :

$$\begin{aligned}
 Q &= Q_k, T = T_k, x = QT^{-1}e_1 \|b\|_2 \\
 r &= b - Ax, \quad T = Q^T A Q \\
 Q^T r &= Q^T (b - Ax) \\
 &= Q^T b - Q^T A x \\
 &= e_1 \|b\|_2 - Q^T A (QT^{-1}e_1 \|b\|_2) \\
 &= e_1 \|b\|_2 - \underbrace{(Q^T A Q)}_{T} T^{-1} e_1 \|b\|_2 \\
 &= 0 \Rightarrow \text{prop 3) holds}
 \end{aligned}$$

Show  $x$  minimizes  $\|r\|_{A^{-1}}^2$

$$x' = x + Qz \quad r' = b - Ax' = r - A(Qz)$$

$$\begin{aligned}
 \|r'\|_{A^{-1}}^2 &= r'^T A^{-1} r' \\
 &= (r - A(Qz))^T A^{-1} (r - A(Qz)) \\
 &= r^T A^{-1} r + \underbrace{= 0?}_{\text{middle}} + (AQz)^T A^{-1} (AQz) \\
 &= \|r\|_{A^{-1}}^2 + \underbrace{= 0?}_{\text{middle}} + \|AQz\|_{A^{-1}}^2
 \end{aligned}$$

$\geq \|r\|_{A^{-1}}^2$  which is minimum  
if  $\text{middle} = 0$

$$\begin{aligned}\text{middle} &= -(AQz)^T A^{-1} r - r^T A^{-1} (AQz) \\ &= -2 (AQz)^T \tilde{A}^{-1} r \\ &= -2 z^T \underbrace{Q^T A \tilde{A}^{-1}}_I r \\ &= -2 z^T \underbrace{Q^T r}_J \Rightarrow \text{prop 4) holds}\end{aligned}$$

$$r_k = b - Ax_k \in \mathcal{X}_{k+1} \\ \in \mathcal{X}_k$$

$r_k$  in span of  $\mathcal{X}_{k+1}$  but not  $\mathcal{X}_k$   
" " "  $Q_{k+1}$  but not  $Q_k$

$\Rightarrow r_k$  must be a multiple of  $g_{k+1}$   
because  $r_k^T Q_k = 0$

$$\Rightarrow r_k = \pm \|r_k\|_2 \cdot g_{k+1}$$

Derive CG starting from (\*)  $x_k = Q_k T_k^{-1} e, \|b\|_2$

Need recurrences for

$x_k$  = solution

$r_k$  = residual

$p_k$  = conjugate gradient

only keep 3 most recent vectors in memory

(1)  $p_k$  called gradient because each step of CG moves  $x_k$  in direction  $p_k$

$$x_k = x_{k-1} + \nu \cdot p_k$$

until  $x_k$  minimizes  $\|r_k\|_{A^{-1}}$   
over all choices of  $\nu$

(2)  $p_k$  called conjugate (A conjugate):

orthogonal w.r.t. inner product

defined by A:  $p_k^T A p_j = 0$  if  $j \neq k$

Can also derive CG starting from (1) and (2), showing they satisfy (\*)

$T_k$  = s.p.d. and tridiagonal  $\Rightarrow$  use Cholesky

$$= L_k' L_k'^T, \quad L_k' \text{ lower bidiagonal}$$

$$\begin{bmatrix} & & & \\ & \diagdown & \diagup & \\ & O & & \\ & \diagup & \diagdown & \\ & & & \end{bmatrix}$$

$$T_k = L_k' L_k'^T = L_k D_k L_k^T$$

$$= \begin{bmatrix} 1 & & & \\ & \ddots & & \\ & & 1 & \\ & & & \ddots & 1 \end{bmatrix} \cdot \begin{bmatrix} & & & \\ & O & & \\ & & & \\ & O & & \\ & & & \ddots & 1 \end{bmatrix} \begin{bmatrix} 1 & & & \\ & \diagdown & \diagup & \\ & O & & \\ & \diagup & \diagdown & \\ & & & \ddots & 1 \end{bmatrix}$$

$L_k$  has unit diagonal,

$$D_k \text{ diagonal} \quad D_k(i,i) = (L_k'(i,i))^2$$

$$(*) \quad x_k = Q_k T_k^{-1} e_1 \|b\|_2$$

$$= Q_k (L_k D_k L_k^T)^{-1} e_1 \|b\|_2$$

$$= Q_k (L_k^{-T} D_k^{-1} L_k^{-1}) e_1 \|b\|_2$$

$$= [Q_k L_k^{-T}] \cdot [D_k^{-1} L_k^{-1} e, \|b\|_2]$$

$$= P'_k \cdot y_k$$

where  $P'_k = [P'_1, P'_2, \dots, P'_k]$ ,  $y_k = \begin{bmatrix} y_{k-1} \\ s_k \end{bmatrix}$

eventual conjugate gradients  $P_k$   
are scalar multiples of  $P'_k$

Proof of 2):  $P'_k$  are A-conjugate,

i.e.  $P'^T A P'_k$  diagonal:

$$= (Q_k L_k^{-T})^T A (Q_k L_k^{-T})$$

$$= L_k^{-T} (Q_k^T A Q_k) L_k^{-T}$$

$$= L_k^{-T} \cdot (T_k) \cdot L_k^{-T}$$

$$= \underbrace{L_k^{-T}}_{I} \cdot \underbrace{(L_k D_k L_k^T)}_{I} \underbrace{L_k^{-T}}_{I}$$

$$= D_k$$

Need recurrence for columns  $P'_k$  of  $P'_k$   
and components of  $y_k$

$$\text{need } P'_k = [P'_{k-1}, P'_k] \quad y_k = \begin{bmatrix} s_1 \\ \vdots \\ s_k \end{bmatrix} = \begin{bmatrix} y_{k-1} \\ s_k \end{bmatrix}$$

If true, get recurrence

$$(Rx) \quad x_k = P'_k y_k = [P'_{k-1}, P'_k] \cdot \begin{bmatrix} y_{k-1} \\ s_k \end{bmatrix}$$

$$= P_{k-1}' \cdot y_{k-1} + P_k' \cdot s_k$$

$$= x_{k-1} + P_k' \cdot s_k$$

also need recurrences for  $P_k'$  and  $s_k$

Since Lanczos constructs  $T_k$  row by row  
 $T_{k-1}$  is leading k-1 by k-1 submatrix of  $T_k$

Since Cholesky works "top-down",  $L_{k-1}$  and  $D_{k-1}$   
are also leading k-1 by k-1 submatrices  
of  $L_k$  and  $D_k$

$$\begin{aligned} T_k &= L_k D_k L_k^T \\ &= \left[ \begin{array}{c|c} L_{k-1} & 0 \\ \hline 0 & D_{k-1} \end{array} \right] \cdot \left[ \begin{array}{c|c} D_{k-1} & 0 \\ \hline 0 & d_k \end{array} \right] \left[ \begin{array}{c|c} L_{k-1} & 0 \\ \hline 0 & D_{k-1} \end{array} \right]^T \\ \Rightarrow L_k^{-1} &= \left[ \begin{array}{c|c} L_{k-1}^{-1} & 0 \\ \hline \text{stuff} & 1 \end{array} \right] \end{aligned}$$

$$\begin{aligned} y_k &= D_k^{-1} L_k^{-1} e_1 \|b\|_2 \\ &= \left[ \begin{array}{c|c} D_{k-1}^{-1} & 0 \\ \hline 0 & d_k^{-1} \end{array} \right] \left[ \begin{array}{c|c} L_{k-1}^{-1} & 0 \\ \hline \text{stuff} & 1 \end{array} \right] e_1 \|b\|_2 \\ &= \left[ \begin{array}{c} D_{k-1}^{-1} L_{k-1}^{-1} e_1 \|b\|_2 \\ s_k \end{array} \right] = \left[ \begin{array}{c} y_{k-1} \\ s_k \end{array} \right] \end{aligned}$$

$$\begin{aligned} P_k' &= Q_k \cdot L_k^{-T} = [Q_{k-1}, q_k] \cdot \left[ \begin{array}{c|c} L_{k-1}^{-T} & \text{stuff} \\ \hline 0 & 1 \end{array} \right] \\ &= [Q_{k-1} L_{k-1}^{-T}, P_k'] = [P_{k-1}', P_k'] \end{aligned}$$

To get recurrence for  $p_k'$ , write

$$Q_k = P_k' \cdot L_k^+ \quad L_k^+ = \begin{bmatrix} I & 0 \\ 0 & I \end{bmatrix}$$

equal last column:

$$g_k = p_k' + p_{k-1}' L_k(k, k)$$

or

$$(R_p) \quad p_k' = g_k - p_{k-1}' l_{k-1}$$

Also need recurrence for  $r_k$ : use  $(R_r)$

$$\begin{aligned} (R_r) \quad r_k &= b - A s_k \\ &= \underbrace{b - A(x_{k-1} + p_k' s_k)}_{r_{k-1} - A p_k' s_k} \end{aligned}$$

All recurrences:

$$(R_r) \quad r_k = r_{k-1} - A p_k' s_k$$

$$(R_x) \quad x_k = x_{k-1} + p_k' \cdot s_k$$

$$(R_p) \quad p_k' = g_k - l_{k-1} p_{k-1}'$$

Substitute  $g_k = r_{k-1} / \|r_{k-1}\|_2$

$$p_k' = p_{k-1}' \cdot \|r_{k-1}\|_2$$

$$\begin{aligned} (R_r') \quad r_k &= r_{k-1} - A p_k (s_k / \|r_{k-1}\|_2) \\ &= r_{k-1} - A p_k \cdot v_k \end{aligned}$$

$$(R_x') \quad x_k = x_{k-1} + p_k \cdot v_k$$

$$\begin{aligned} (R_p') \quad p_k &= r_{k-1} - \left( \|r_{k-1}\|_2 \cdot l_{k-1} / \|r_{k-1}\|_2 \right) p_{k-1} \\ &= r_{k-1} + N_k \cdot p_{k-1} \end{aligned}$$

just 3 vectors:  $P$ ,  $x$ ,  $r$

Need formulas for  $\nu_k$  and  $\mu_k$ :

Take  $(R_P')$ , multiply by  $r_{k-1}^T$ , get

$$r_{k-1}^T r_k = r_{k-1}^T r_{k-1} - r_{k-1}^T A P_k \cdot \nu_k$$

$$0 = "$$

$$\Rightarrow \nu_k = \frac{r_{k-1}^T r_{k-1}}{r_{k-1}^T A P_k}$$

More stable version:

$$(R_P')^T \cdot A P_k \Rightarrow$$

$$(P_k = r_{k-1} + N_k P_{k-1})^T A P_k$$

$$\Rightarrow P_k^T A P_k = r_{k-1}^T A P_k + N_k \cdot P_{k-1}^T A P_k$$

$$\Rightarrow P_k^T A P_k = r_{k-1}^T A P_k$$

$$\Rightarrow \nu_k = \frac{r_{k-1}^T r_{k-1}}{P_k^T A P_k}$$

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Formula for  $\mu_k$

$$(R_P')^T A P_{k-1} \Rightarrow$$

$$(P_k = r_{k-1} + N_k P_{k-1})^T A P_{k-1}$$

$$\Rightarrow P_k^T A P_{k-1} = r_{k-1}^T A P_{k-1} + N_k P_{k-1}^T A P_{k-1}$$

$$\Rightarrow \mu_k = \frac{-r_{k-1}^T A P_{k-1}}{P_{k-1}^T A P_{k-1}}$$

avoid extra dot product in numerator

$$(Rr')^+ r_k \Rightarrow$$

$$(r_k = r_{k-1} - A p_k \cdot v_k)^T r_k$$

$$\Rightarrow r_k^T r_k = r_{k-1}^T r_k - p_k^T A r_k \cdot v_k \\ = 0$$

$$\Rightarrow v_k = \frac{-r_k^T r_k}{p_k^T A r_k}$$

$$\Rightarrow \frac{-r_k^T r_k}{p_k^T A r_k} = \frac{r_{k-1}^T r_{k-1}}{p_{k-1}^T A p_{k-1}}$$

$$\Rightarrow \frac{r_k^T r_k}{r_{k-1}^T r_{k-1}} = - \frac{p_k^T A r_k}{p_{k-1}^T A p_{k-1}} = \nu_{k+1}$$

$$\Rightarrow \underline{\nu_k = \frac{r_{k-1}^T r_{k-1}}{r_{k-2}^T r_{k-2}}}$$

Put it altogether to get CG for  
 $Ax=b$

$$k=0, x_0=0, r_0=b, p_0=b$$

repeat

$$k=k+1$$

$$z = A \cdot p_k$$

$$v_k = (r_{k-1}^T \cdot r_{k-1}) / (p_k^T \cdot z)$$

$$(Rx') \quad x_k = x_{k-1} + v_k \cdot p_k$$

$$(Rr') \quad r_k = r_{k-1} - v_k \cdot z$$

$$\nu_{k+1} = \frac{r_k^T \cdot r_k}{r_{k-1}^T \cdot r_{k-1}}$$

$(R\rho')$   $P_{k+1} = r_k + N_{k+1} \cdot P_k$   
 until  $\|r_k\|_2$  small enough

cost: 2 dot products  $r_k^T r_k$  and  $P_k^T z$   
 3 saxpy's  
 1  $A \cdot P_k$

Convergence:

$$\text{Thm } \frac{\|r_k\|_{A^{-1}}}{\|r_0\|_{A^{-1}}} \leq \frac{2}{1 + \sqrt{\frac{2}{\text{cond}(A) - 1}}}$$

if  $\text{cond}(A)$  large, need  $O(\sqrt{\text{cond}(A)})$   
 steps to converge

For d-dimensional Poisson on  
 mesh with n vertices in each  
 direction,  $\text{cond}(A) \sim O(n^2)$

$\Rightarrow$  CG takes  $O(n)$  steps to converge