

Welcome back to Ma221! Lecture 40, Dec 1

Finish CG, some preconditioning

Derived CG from (\*)  $x_k = Q_k T_k^{-1} e_1 \|b\|_2$

Need recurrences for

$x_k =$  solution

$r_k =$  residual

$p_k =$  conjugate gradient

$$(Rr') \quad r_k = r_{k-1} - A \cdot p_k \cdot \nu_k$$

$$(Rx') \quad x_k = x_{k-1} + p_k \nu_k$$

$$(Rp') \quad p_k = r_{k-1} + \mu_k p_{k-1}$$

Need formulas for  $\nu_k, \mu_k$

(see book for details)

eg: take (Rr') multiply by  $r_{k-1}^T$

$$\Rightarrow r_{k-1}^T r_k = 0 = r_{k-1}^T r_{k-1} - r_{k-1}^T A p_k \nu_k$$

$$\Rightarrow \nu_k = \frac{r_{k-1}^T r_{k-1}}{r_{k-1}^T A p_k} \quad (\text{see book})$$

Final Formulas:

$$\nu_k = \frac{r_{k-1}^T r_{k-1}}{p_k^T A p_k}$$

$$\mu_k = \frac{r_{k-1}^T r_{k-1}}{r_{k-2}^T r_{k-2}}$$

Altogether: CG for  $Ax=b$

$$k=0, x_0=0, r_0=b, p_0=b$$

repeat

$$k=k+1$$

$$z = A \cdot p_k$$

$$\nu_k = (r_k^T \cdot r_{k-1}) / (p_k^T \cdot z)$$

$$x_k = x_{k-1} + \nu_k \cdot p_k$$

$$r_k = r_{k-1} - \nu_k \cdot z$$

$$\mu_k = \frac{r_k^T r_k}{r_{k-1}^T r_{k-1}}$$

$$p_{k+1} = r_k + \mu_k p_k$$

until  $\|r_k\|_2$  small enough

Convergence:

$$\lim \frac{\|r_k\|_{A^T}}{\|r_0\|_{A^T}} \leq \frac{2}{1 + \frac{2k}{\sqrt{\kappa(A)} - 1}}$$

when  $\kappa(A) \gg 1$  need

$O(\sqrt{\kappa(A)})$  steps to converge

proof uses Chebyshev polynomials.

For  $d$ -dimensional Poisson on  $n^d$  mesh

$\kappa(A) \sim O(n^2) \Rightarrow$  CG take  $O(n)$  steps  
to converge

## Preconditioning to accelerate convergence

How to change  $Ax=b$  cheaply to converge faster:

Simplest  $M^{-1}Ax = M^{-1}b$  where  $M \approx A$

(1)  $M^{-1}$  cheaper to multiply by than  $A$

(2)  $M^{-1}A$  better conditioned than  $A$

Idea straight forward for GMRES, just multiply by  $A$  and  $M^{-1}$ , CG trickier

$A$  s.p.d.  $\Rightarrow$  choose  $M$  s.p.d., but

$M^{-1}A$  not s.p.d.

$$M \text{ s.p.d.} \Rightarrow M = Q\Lambda Q^T \Rightarrow M^{1/2} = Q\Lambda^{1/2}Q^T$$

Imagine applying CG to

$$(*) \quad \underbrace{(M^{-1/2} A M^{1/2})}_{\text{s.p.d.}} (M^{1/2} x) = M^{-1/2} b \quad \left( \begin{array}{l} \text{Same as} \\ Ax=b \end{array} \right)$$

$$M^{-1}A \text{ and } M^{-1/2} A M^{-1/2} = M^{-1/2} (M^{-1}A) M^{-1/2} \text{ are similar}$$

$$\Rightarrow \text{same } \frac{\lambda_{\max}}{\lambda_{\min}}$$

Preconditioned CG (equivalent to CG on  $(*)$ )

$$k=0, x(0)=0, r(0)=b, p(0)=M^{-1}b, y(0)=M^{-1}r(0)$$

repeat

$$k = k + 1$$

$$z = A \cdot p(k)$$

$$\nu(k) = (y(k-1)^T \cdot r(k-1)) / (p(k)^T \cdot z)$$

$$X(k) = X(k-1) + \nu(k) \cdot p(k)$$

$$r(k) = r(k-1) - \nu(k) \cdot z$$

$$y(k) = M^{-1} \cdot r(k)$$

$$\nu(k+1) = (y(k)^T \cdot r(k)) / (y(k-1)^T \cdot r(k-1))$$

$$p(k+1) = y(k) + \nu(k+1) \cdot p(k)$$

until  $\|r(k)\|_2$  small enough

Thm 6.9: equivalent to (\*)

HW Q 6.14 (see errata on p 317)

How to pick a good  $M$ ?

Goals: (1) cheap to multiply by  $M^{-1}$   
(2)  $\text{cond}(M^{-1}A) \ll \text{cond}(A)$

(1) If  $A$  s.p.d.  $M = \text{diag}(A_{11}, A_{22}, \dots, A_{nn})$

"Jacobi preconditioning"

$$M^{-1}A = I - R_s$$

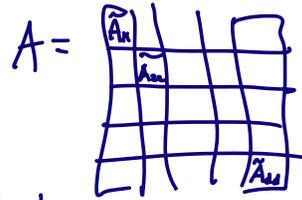
$\rho(R_j) \ll 1 \Rightarrow M^{-1}A$  well conditioned

Thm (van der Sluis, 1969) if  $M$  diagonal

then  $M^{-1}A = \text{diag}(A_{11}, \dots, A_{nn})$  within factor

$n$  of optimal for minimizing  $\text{cond}(M^{-1/2}AM^{1/2})$

$$(2) M = \text{diag}(\tilde{A}_{11}, \dots, \tilde{A}_{dd})$$



more expensive than plain diagonal, but can be faster

"block Jacobi preconditioning"  
each  $\tilde{A}_{ii}$  could correspond to a different subset of physical simulation

Thm (D, 2023) above choice of  $M$  within factor of  $d$  of optimal over any block diagonal  $M$ , for minimizing  $\text{cond}(M^{-1/2} A M^{-1/2})$

(3) "Incomplete Cholesky"

complete partial Cholesky  $LL^T \approx A$

by limiting fill-in in  $L$

Easiest: no fill-in  $A_{ij} = 0 \Rightarrow L_{ij} = 0$

ILU(0)  $\Rightarrow$  need triangular solve to

(can also do ILU(k)) Multiply by  $M^{-1} = (L^T)^{-1} = L^{-T} L^{-1}$

cost =  $O(nnz(L)) = O(nnz(A))$

same cost as  $Ax$

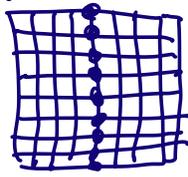
(4) A few steps of multigrid (perhaps just applied to  $\tilde{A}_{ii}$ , or to coarsest mesh)

(5) Domain Decomposition

Ex: discretize FE problem on 2D mesh

Left is steel, right is concrete

Interface between them  
steel interface



concrete

$n \times n$  mesh

$$A = \begin{bmatrix} A_{ss} & \text{O} & A_{si} \\ \text{O} & A_{cc} & A_{ci} \\ A_{is} & A_{ic} & A_{ii} \end{bmatrix}$$

$\begin{matrix} \frac{n(n-1)}{2} & \frac{n(n-1)}{2} & n \end{matrix}$

$$A = \begin{bmatrix} I & \text{O} & \text{O} \\ \text{O} & I & \text{O} \\ A_{is} A_{ss}^{-1} & A_{ic} A_{cc}^{-1} & I \end{bmatrix} \begin{bmatrix} I & \text{O} & \text{O} \\ \text{O} & I & \text{O} \\ \text{O} & \text{O} & S \end{bmatrix} \begin{bmatrix} A_{ss} & \text{O} & A_{si} \\ \text{O} & A_{cc} & A_{ci} \\ \text{O} & \text{O} & I \end{bmatrix}$$

Schur complement =  $S = A_{ii} - A_{is} A_{ss}^{-1} A_{si} - A_{ic} A_{cc}^{-1} A_{ci}$

$$A^{-1} = \begin{bmatrix} A_{ss}^{-1} & \text{O} & -A_{ss}^{-1} A_{si} \\ \text{O} & A_{cc}^{-1} & -A_{cc}^{-1} A_{ci} \\ \text{O} & \text{O} & I \end{bmatrix} \cdot \begin{bmatrix} I & & \\ & I & \\ & & S^{-1} \end{bmatrix} \cdot \begin{bmatrix} I & \text{O} & \text{O} \\ \text{O} & I & \text{O} \\ A_{is} A_{ss}^{-1} & -A_{ic} A_{cc}^{-1} & I \end{bmatrix}$$

need to multiply by  $A^{-1}$  approximately  
useful if good preconditioners for  
 $A_{ss}^{-1}, A_{cc}^{-1}$  available

Ex: multiply  $(A_{ss}^{-1} A_{si})x = A_{ss}^{-1} (A_{si}x)$

↑  
↑ spmv  
use preconditioned solver for  $A_{ss}$

What about  $S^{-1}$ ?

Can multiply by  $S$ :

$$Sx = (A_{ii}x) - A_{is}(A_{ss}^{-1}(A_{si}x)) - A_{ic}(A_{cc}^{-1}(A_{ci}x))$$

Given ability to multiply  $S$ , use

CG to solve  $Sx=b$ , ie  $S^{-1}b$

Often  $S$  much better conditioned, and smaller than  $A$

Then  $\text{cond}(S) \leq \text{cond}(A)$  if  $A$  s.p.d.

Ex for Poisson  $\kappa(S) = \sqrt{\kappa(A)}$

called "non overlapping domain decomposition"  
can also do "overlapping domain decomposition"