

Welcome to Ma 221! Lecture 27, Oct 25

Power Method  $x_{i+1} = Ax_i / \|x_i\|_2$

Inverse Iteration  $x_{i+1} = (A - \sigma I)^{-1} x_i / \|x_i\|_2$

Orthogonal Iteration

$Z_i^{n \times p}$   $Y_{i+1} = AZ_i$ ,  $Z_{i+1} R_{i+1} = QR(Y_{i+1})$

Thm: Orthog Iter. with  $Z_0 = I^{n \times n}$

$|\lambda_1| > |\lambda_2| > \dots \Rightarrow$

$[A_i = Z_i^T A Z_i \rightarrow \text{Schur form}]$

Next: QR iteration (let's us introduce inverse iter.)

$A_0 = A$

$i = 0$

repeat

$[A_i = Q_i R_i \dots \text{QR decomp}]$

$[A_{i+1} = R_i Q_i]$

$i = i + 1$

until convergence

Thm:  $A_i$  from QR iteration same as

$Z_i^T A Z_i$  from orthog. iter

proof (induction) assume  $A_i = Z_i^T A Z_i$

One step of Orthog iter

$A Z_i = Z_{i+1} R_{i+1} \dots \text{QR decomp}$

$A_i = Z_i^T A Z_i \dots \text{induct. hypoth.}$

$$\begin{aligned}
 &= \underbrace{Z_c^T Z_{i+1}}_{\text{orthog}} \underbrace{R_{i+1}}_{\nabla} \dots \text{ from orthog. iter.} \\
 &= \text{QR decomp of } A_{i+1}, \text{ by uniqueness} \\
 Z_{i+1}^T A Z_{i+1} &= (Z_{i+1}^T A Z_{i+1}) (Z_c^T Z_{i+1}) \\
 &= (R_{i+1}) (Z_c^T Z_{i+1}) \\
 &= R Q \\
 &= A_{i+1}
 \end{aligned}$$

Add inverse iteration to QR iteration

$$A_0 = A$$

$$i = 0$$

repeat

$$\text{factor } A_i - \sigma I = Q_i R_i$$

$$A_{i+1} = R_i Q_i + \sigma I$$

$$i = i + 1$$

until convergence

Lemma:  $A_i$  and  $A_{i+1}$  are orthog. similar

$$\text{proof: } A_{i+1} = R_i Q_i + \sigma I$$

$$= Q_i^T (Q_i R_i) Q_i + \sigma I$$


$$= Q_i^T (A_i - \sigma I) Q_i + \sigma I$$

$$= Q_i^T A_i Q_i - \underbrace{\sigma Q_i^T Q_i + \sigma I}_{\text{cancel}}$$

$$= Q_i^T A_i Q_i$$

Note: if  $R_i$  nonsingular, can also write

$$\begin{aligned}
 A_{ii} &= R_i Q_i + \sigma_i I \\
 &= R_i \underbrace{Q_i R_i^{-1}} + \sigma_i I \\
 &= R_i (A_i - \sigma_i I) R_i^{-1} + \sigma_i I \\
 &= R_i A_i R_i^{-1} - \underbrace{\sigma_i R_i R_i^{-1}}_{\text{cancel}} + \sigma_i I
 \end{aligned}$$

$$= R_i A_i R_i^{-1}$$


Upper Hessenberg

If  $\sigma_i = \text{exact eval of } A$

QR converges in one step

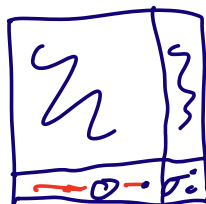
$A_i - \sigma_i I$  singular  $\Rightarrow R_i$  singular

$\Rightarrow R_i(k, k) = 0$ , suppose  $R_i(n, n) = 0$

$\Rightarrow$  last row of  $R_i = 0$

$\Rightarrow$  last row of  $R_i Q_i = 0$

$\Rightarrow$  last row of  $A_{ii} = R_i Q_i + \sigma_i I$   
is zero, except for  $A_{ii}(n, n) = \sigma_i$



Q4.1

If  $\sigma_i$  not exact eval, but "close enough", i.e.  $\|A_{i+1}(n, 1:n-1)\|_2 = O(\epsilon)$

set  $A_{i+1}(n, 1:n-1) = 0$   
backward stable

Previous analysis: expect  $A_{i+1}(n, 1:n-1)$  to shrink by factor

$$\frac{|\lambda_k - \sigma_i|}{\min_{j \neq k} |\lambda_j - \sigma_i|} \quad \lambda_k \text{ is closest eval to } \sigma_i$$

$$A_i - \sigma_i I = Q_i R_i$$

$$Q_i^T (A_i - \sigma_i I) = R_i$$

$$\text{if } \sigma_i \text{ were } = \lambda_k \Rightarrow R_i(n, n) = 0$$

$$\Rightarrow \text{last row of } Q_i^T (A_i - \sigma_i I) = 0$$

$\Rightarrow$  last col of  $Q_i$  is left evec of  $A_i$  for  $\sigma_i$

Now suppose  $\sigma_i$  just close to eval

$$A_i - \sigma_i I = Q_i R_i$$

$$(A_i - \sigma_i I)^T = R_i^T Q_i^T$$

$$(A_i - \sigma_i I)^{-T} = Q_i R_i^{-T}$$

$$(A_i - \sigma_i I)^{-T} R_i^T = Q_i$$

△

$$(A_i - \sigma_i I)^T e_n R_i(n, n) = Q_i \cdot e_n$$

⇒ one step of inverse iteration starting with  $e_n$ , expect last col of  $Q_i$  to be approx evec

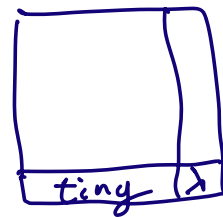
⇒ last col of  $Q_i$  closer to evec of  $A_i^T$

⇒ last col of  $A_i^T Q_i$  closer to  $\lambda \cdot$  last col of  $Q_i$

⇒ <sup>last col of</sup>  $Q_i^T A_i^T Q_i$  closer to  $\lambda \cdot e_n$

⇒ last row of  $Q_i^T A_i^T Q_i$  closer to  $\lambda e_n^T$

⇒



Where do we get  $\sigma_i$ ?  
use  $\sigma_i = A_i(n, n)$

⇒ yields quadratic convergence

$$\|A_i(n, 1:n-1)\|_2 = \epsilon \ll 1$$

$$|A_i(n, n) - \lambda|_2 = O(\epsilon) \leftarrow$$

expect  $\|A_i(n, 1:n-1)\|_2$  gets  
multiplied by  $|A_i(n, n) - \lambda| = O(\epsilon)$

$\Rightarrow$  becomes  $O(\epsilon^2)$

Matlab demo:  
see typed notes  
for code

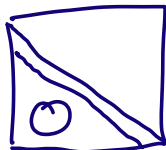
---

Making QR iteration practical

- 1) Each iteration costs 1 QR  
+ 1 matmul  $= O(n^3)$ , if we did  
only  $O(1)$  iterations per eval  
 $\Rightarrow O(n^4)$  cost total, want  $O(n^3)$
- 2) How to shift to converge to  
real Schur form?
- 3) How do decide on convergence?
- 4) How to minimize communication?

Answers:

- (1) preprocess  $A = Q H Q^T$  where  
Q orthog, H Upper Hessenberg



QR iteration preserves  $O_s$ ,  
 lowers cost to  $O(n^2)$   
 $\Rightarrow \text{cost} = O(n^3)$

$$A = A^T \Rightarrow H = H^T \quad H^T = (Q^T A Q)^T$$

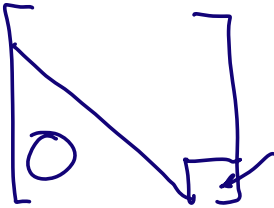
$\Rightarrow H$  tridiagonal 

$\Rightarrow \text{cost} = O(n)$  (chap 5)

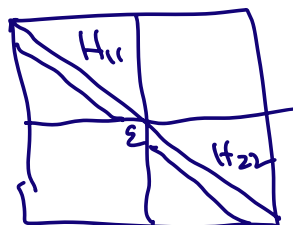
(2) Converge to Real Schur Form:

since evals appear in  $\lambda, \bar{\lambda}$  pairs,  
 turns out that 1 step of QR  
 iteration with shift =  $\lambda$ , followed  
 by  $\bar{\lambda}$  as shift  $\Rightarrow A_{i+2}$  real

$\Rightarrow$  don't compute any imaginary parts

  $\lambda, \bar{\lambda} = \text{eigen values}$   
 of  $H(n-1:n, n-1:n)$

(3) Detect convergence?



if  $\epsilon \leq \text{macheps} \cdot \|A\|$   
 set  $\epsilon$  to 0

work independently on  $H_{11}$ ,  $H_{22}$