

Welcome back to Ma221! Lecture 24, Oct 18

## Eigenproblems

Lemma: if  $Ax_i = \lambda_i x_i$  for  $i=1:n$

and  $S = [x_1, \dots, x_n]$  nonsingular

$$\Rightarrow A = S \begin{bmatrix} \lambda_1 & & \\ & \ddots & \\ & & \lambda_n \end{bmatrix} S^{-1}$$

Conversely if  $A = S \Lambda S^{-1}$ ,  $\Lambda = \text{diag}(\lambda_1, \dots, \lambda_n)$

then columns of  $S$  are evecs

Proof:  $A = S \Lambda S^{-1}$  iff  $AS = S\Lambda$

$$\text{iff } AS[:, i] = S[:, i] \cdot \lambda_i \quad \text{i.e.}$$

$S[:, i]$  evec for eval  $\lambda_i$

But we can't always diagonalize  $A$ :

may be mathematically impossible  
(recall Jordan form)

may be unstable even if it exists  
(when evals close together)

Recall Jordan Form: For any  $A$

exists similarity  $S$ :  $SAS^{-1} = J = \text{diag}(J_1, \dots, J_k)$

$$J_i = \begin{bmatrix} \lambda_i & & \\ & \ddots & \\ & & \lambda_i \end{bmatrix}$$

upto permuting order of  $J_i$ , unique.

$\lambda_i$  can be same or different eg  $A = I$

only one right or left evec per  $\lambda_i$ :

$$\begin{bmatrix} \lambda & 0 \\ 0 & \lambda \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \lambda \begin{bmatrix} 1 \\ 0 \end{bmatrix} \Rightarrow \text{can't diagonalize}$$

Why not compute Jordan Form?

Consider slightly perturbed  $2 \times 2$  I

①  $\begin{bmatrix} 1 & 0 \\ 0 & 1+\epsilon \end{bmatrix} : (1, \begin{bmatrix} 1 \\ 0 \end{bmatrix}), (1+\epsilon, \begin{bmatrix} 0 \\ 1 \end{bmatrix})$

②  $\begin{bmatrix} 1 & \epsilon \\ \epsilon & 1 \end{bmatrix} : (1+\epsilon, \begin{bmatrix} 1 \\ 1 \end{bmatrix}), (1-\epsilon, \begin{bmatrix} 1 \\ -1 \end{bmatrix})$   
evecs rotate  $45^\circ$

③  $\begin{bmatrix} 1 & \epsilon \\ 0 & 1+\epsilon^2 \end{bmatrix} : (1, \begin{bmatrix} 1 \\ 0 \end{bmatrix}), (1+\epsilon^2, \begin{bmatrix} 1 \\ \epsilon \end{bmatrix})$   
evecs nearly parallel

④  $\begin{bmatrix} 1 & \epsilon \\ 0 & 1 \end{bmatrix} : (1, \begin{bmatrix} 1 \\ 0 \end{bmatrix})$  just one evec

⑤  $\begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} : (1: \text{anything}), (1, \text{anything})$

Jordan form very ill conditioned

Our goal: exact evals, evecs of

$$A + E, \quad \frac{\|E\|}{\|A\|} = O(\varepsilon)$$

Backward Stable approach to eigen problems  
 Only use orthogonal similarities

Recall Chap 3: Multiplying by multiple  
 orthog. matrices is backward stable

$$f_l(Q_k(\dots(Q_2(Q_1 A)^\dagger)\dots)) = Q(A + E)$$

exact orthog       $\frac{\|E\|}{\|A\|} = O(\epsilon)$

Apply twice to orthog. similarity

$$f_l(Q_k(\dots Q_2(Q_1 A Q_1^\dagger) Q_2^\dagger \dots) Q_k^\dagger) = Q(A + F) Q^\dagger$$

$Q Q^\dagger = I$ ,  $\frac{\|F\|}{\|A\|} = O(\epsilon)$

If we restrict to orthog. similarities,  
 how close to Jordan form can we get?

Thm: Schur Canonical Form:

Given any  $n \times n A \exists$  unitary  $Q$

$Q Q^H = I$  s.t.  $Q^H A Q = T = \text{upper triangular}$   
 evals of  $T$  are  $T(i,i)$

Compute evals of  $T$  just triangular solve.

$$\begin{matrix} i-1 & & & & n-i \\ \vdots & \left[ \begin{array}{ccc} T_{11} & T_{12} & \cdots & T_{1n-i} \\ 0 & T(i,i) & & T_{i,n-i} \\ & & \ddots & \\ 0 & 0 & & T_{n-i,n-i} \end{array} \right] & \left[ \begin{array}{c} x_1 \\ x_2 \\ \vdots \\ x_n \end{array} \right] & = & \left[ \begin{array}{c} k_1 \\ k_2 \\ \vdots \\ k_n \end{array} \right] \end{matrix}$$

$$\begin{aligned} T_{11}x_1 + T_{12}x_2 + T_{13}x_3 &= T(i,i)x_1 \\ T(i,i)x_2 + T_{23}x_3 &= T(i,i)x_2 \\ T_{33}x_3 &= T(i,i)x_3 \end{aligned}$$

Suppose  $T(i,i)$  unique

$$\underbrace{(T_{33} - T(i,i)I)}_{\text{non zero diagonal}} x_3 = 0 \Rightarrow x_3 = 0$$

$$T(i,i)x_2 = T(i,i)x_2 \quad x_2 = 1$$

$$(T_{11} - T(i,i)I)x_1 = -T_{12}x_2$$

triangular solve

What does Matlab do with

$$[V, D] = \text{eig}([0 \ 1 \ 0 \ 0]) \ ? \ \text{try it!}$$

Proof that Schur form exists:

choose any  $\lambda$ , with evcc  $x$ :  $Ax = \lambda x$   
 $\|x\|_2 = 1$

Let  $Q = [x, Q']$  be unitary

$$Q^H A Q = \begin{bmatrix} x^* \\ Q'^* \end{bmatrix} A \begin{bmatrix} x \\ Q' \end{bmatrix}$$

$$= \begin{array}{c|c} x^* A x & x^* A Q' \\ \hline Q'^* A x & Q'^* A Q' \end{array}$$

$$= \left[ \begin{array}{c|c} x^* \lambda x & " \\ \hline Q'^H \lambda x & " \end{array} \right]$$

$$= \left[ \begin{array}{c|c} \lambda & " \\ \hline 0 & Q'^H A Q' \end{array} \right]$$

induction on  $Q'^* A Q' = U^H T U$

$$= \left[ \begin{array}{c|c} \lambda & x^* A Q' \\ \hline 0 & U^H T U \end{array} \right]$$

$$= \left[ \begin{array}{c|c} 1 & 0 \\ \hline 0 & U^* \end{array} \right] \left[ \begin{array}{c|c} \lambda & x^* A Q' U^* \\ \hline 0 & T \end{array} \right] \left[ \begin{array}{c|c} 1 & 0 \\ \hline 0 & U \end{array} \right]$$

$$\Rightarrow \underbrace{\left[ \begin{array}{c|c} 1 & 0 \\ \hline 0 & U \end{array} \right]}_{\text{unitary}} Q^H A Q \underbrace{\left[ \begin{array}{c|c} 1 & 0 \\ \hline 0 & U^* \end{array} \right]}_{\text{inverse}} = \left[ \begin{array}{c|c} \lambda & x^* A Q' U^* \\ \hline 0 & T \end{array} \right]$$

What about real matrices with complex evals?

[if  $A = A^T$ , all evals real, see Chap 5]

Prefer real arithmetic!

reduce #flops

less memory

make sure that all evals, evecs appear in complex conjugate pairs

$$\underset{\text{real}}{A} \underset{\text{complex}}{x} = \lambda x \Leftrightarrow A \bar{x} = \bar{\lambda} \bar{x}$$

Instead of  $T = \text{upper triangular}$   
use  $T = \text{block triangular}$

$$T = \begin{bmatrix} T_{11} & T_{12} & \cdots & T_{1k} \\ & T_{22} & & \\ & & \ddots & \\ & & & T_{kk} \end{bmatrix} \quad T_{ii} \text{ square}$$

$$\text{evals}(T) = \bigcup_{i=1}^k \text{evals}(T_{ii}) \quad (\text{HW 4.1})$$

Goal: have all  $T_{ii}$  be  $1 \times 1$  for real evals  
or  $2 \times 2$  for complex conjugate pairs

$\lambda, \bar{\lambda}$   
Thm: (Real Schur Canonical form)

Given any real  $A$ ,  $\exists$  real orthog  $Q$   
st  $Q^T A Q$  is block upper triangular  
with  $1 \times 1$  and  $2 \times 2$  diag. blocks

Generalize to "invariant subspaces"

Def:  $V = \text{span}\{x_1, \dots, x_m\} = \text{span}(X)$

$X = [x_1, \dots, x_m]$  subspace of  $\mathbb{R}^n$

$V$  invariant if  $A \cdot V = \text{span}(AX) \subseteq V$

Ex:  $V = \text{span}\{x\} = \{\alpha x \text{ for all } \alpha \in \mathbb{R}\}$

where  $Ax = \lambda x$

$$AV = \{A(\alpha x) \mid \alpha \in \mathbb{R}\}$$

$$= \{\alpha \lambda x, \forall \alpha\}$$

$$\subseteq V, = V \text{ unless } \lambda = 0$$

Ex:  $V = \text{span}\{x_1, \dots, x_k\} = \left\{ \sum_{i=1}^k \alpha_i x_i \mid \alpha_i \in \mathbb{R} \right\}$   
 where  $Ax_i = \lambda_i x_i$

$$AV = \left\{ A \sum \alpha_i x_i \right\} = \left\{ \sum \alpha_i Ax_i \right\} \\ = \left\{ \sum \alpha_i \lambda_i x_i \right\} \subseteq V$$

Lemma:  $V = \{x_1, \dots, x_n\}$  invariant

$$\text{then } \exists B: AX = XB^{n \times n}$$

$$\boxed{\quad} \boxed{\quad} = \boxed{\quad} \boxed{\quad}$$

the evals of  $B$  are evals of  $A$

proof: existence of  $B$  follows from  
 def of invariance

$$Ax_i \in V \Rightarrow \exists \text{ scalars } B(1,i), B(2,i) \dots B(n,i)$$

$$Ax_i = \sum_{j=1}^n x_j B(j,i) \text{ i.e. } AX = XB$$

$$By = \lambda y \Rightarrow A(Xy) = (Ax)y =$$

$$(XB)y = X(By) = \underbrace{(Xy)}_{\text{evcc of } A}$$

Lemma: Let  $V = \text{span}(X)$  be  $n$ -dimensional

Invariant subspace of  $A$ ,

$$AX = X \mathcal{B}, \quad X = QR$$

Let  $[Q, Q']$  be square orthogonal matrix

$$[Q, Q']^T A [Q, Q'] = \begin{bmatrix} A_{11} & A_{12} \\ 0 & A_{22} \end{bmatrix}$$

$A_{11} = RBR^{-1}$  has same evals as  $B$

$$\begin{aligned} \text{Proof: } [Q, Q']^T A [Q, Q'] &= \begin{bmatrix} Q^T A Q & Q^T A Q' \\ Q'^T A Q & Q'^T A Q' \end{bmatrix} \\ &= \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} = \begin{bmatrix} RBR^{-1} & A_{12} \\ 0 & A_{22} \end{bmatrix} \end{aligned}$$

$$AQ = AXR^{-1} = XBR^{-1} = QRBR^{-1}$$

$$A_{11} = Q^T A Q = \underbrace{Q^T}_{Q'} \underbrace{QRBR^{-1}}_{RBR^{-1}} = RBR^{-1}$$

$$A_{21} = \underbrace{Q'^T}_{0} RBR^{-1} = 0$$