

Welcome back to Ma221! Lecture 20, Oct 9

Dealing with (nearly) low rank matrices

Motivation: Real data often low rank
(or nearly redundant)

- ① take precautions to avoid inaccurate LS solution
- ② use it to compress data, go faster, both deterministic and randomized

use LS to illustrate compression
useful elsewhere.

Solving a LS problem when A rank deficient

Thm: A $m \times n$ $m \geq n$ rank $r < n$

$$A = U \Sigma V^T = \begin{matrix} m \\ \left[\begin{array}{ccc} U_1 & U_2 & U_3 \end{array} \right] \end{matrix} \begin{matrix} r & n-r \\ \left[\begin{array}{cc} \Sigma_1 & 0 \\ 0 & \Sigma_2 \\ 0 & 0 \end{array} \right] \end{matrix} \begin{matrix} n \\ \left[\begin{array}{cc} V_1 & V_2 \end{array} \right]^T \end{matrix}$$

Σ_1 full rank r
 $\Sigma_2 = 0$

The set of vectors minimizing $\|Ax - b\|_2$

$$\text{is } \left\{ x = V_1 \Sigma_1^{-1} U_1^T b + V_2 y_2, \text{ any } y_2 \in \mathbb{R}^{n-r} \right\}$$

Unique x minimizing $\|Ax - b\|_2$ and $\|x\|_2$
is gotten from $y_2 = 0$

$$x = \underline{V_1 \Sigma_1^{-1} U_1^T} b$$

Def: $A^+ = V_1 \Sigma_1^{-1} U_1^T$ is Moore-Penrose
pseudoinverse of A (includes full rank
 $r=n$ case)

(in practice Σ_2 will be all singular values
less than user-defined threshold)

So square or not, full rank or not
"best solution" is $x = A^+ b$

Proof: $\|Ax - b\|_2 = \|U \Sigma V^T x - b\|_2$
 $= \|\Sigma V^T x - U^T b\|_2$ since U orthog
 $= \|\Sigma y - U^T b\|_2$ where $y = V^T x$

$\|x\|_2 = \|y\|_2$ so okay to
minimize either one

$$= \left\| \begin{bmatrix} \Sigma_1 y_1 - U_1^T b \\ -U_2^T b \\ -U_3^T b \end{bmatrix} \right\|_2 \quad y = \begin{bmatrix} y_1 \\ y_2 \end{bmatrix}$$

minimized by $y_1 = \Sigma_1^{-1} U_1^T b$

$$\text{and } \|x\|_2^2 = \|y\|_2^2 = \|y_1\|_2^2 + \|y_2\|_2^2$$

minimized by $y_2 = 0$

$$\begin{aligned} x = Vy &= [V_1, V_2] \begin{bmatrix} y_1 \\ y_2 \end{bmatrix} = V_1 y_1 + V_2 y_2 \\ &= V_1 \Sigma_1^{-1} U_1^T b \\ &= A^+ b \end{aligned}$$

Solving LS when A (nearly) rank deficient
using truncated SVD

$$\kappa(A) = \frac{\sigma_{\max}}{\sigma_{\min}} = \infty \text{ if } A \text{ low rank}$$

$$\operatorname{argmin}_x \left\| \begin{bmatrix} 0 & 0 \\ 0 & 0 \\ 0 & 0 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} - \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix} \right\|_2 = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$$

$$\operatorname{argmin}_x \left\| \begin{bmatrix} 1 & 0 \\ 0 & \epsilon \\ 0 & 0 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} - \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix} \right\|_2 = \begin{bmatrix} 1 \\ 1/\epsilon \end{bmatrix} \text{ a ting}$$

What does a "solution" mean if
it can change discontinuously?

Often A not known exactly, just
up to some tolerance $\|A - A'\|_2 \leq \text{tol}$

What to do?

Def: Truncated SVD:

$$\text{if } A = U \Sigma U^T$$

$$A(\text{tol}) = U \Sigma(\text{tol}) V^T$$

$$\Sigma(\text{tol}) = \text{diag}(\sigma_1(\text{tol}), \sigma_2(\text{tol}), \dots, \sigma_n(\text{tol}))$$

$$\sigma_i(\text{tol}) = \begin{cases} \sigma_i & \text{if } \sigma_i \geq \text{tol} \\ 0 & \text{if } \sigma_i < \text{tol} \end{cases}$$

$A(\text{tol})$ = lowest rank matrix within
distance tol to A

Use $A(tol)$ instead of A for LS

reduces $\kappa = \frac{\sigma_{\max}}{\sigma_{\min}}$ to $\frac{\sigma_{\max}}{tol}$

tol is a "knob" to trade off sensitivity and how well LS problem can be solved, i.e. how small you can make $\|Ax - b\|_2$

Replacing A by an "easier" matrix, called regularization, $A(tol)$ is one way, others too

Lemma: $x_1 = \operatorname{argmin}_x \|A(tol)x - b_1\|_2$

$x_2 = \operatorname{argmin}_x \|A(tol)x - b_2\|_2$

choose x_i of smallest norm

then $\|x_1 - x_2\|_2 \leq \frac{\|b_1 - b_2\|_2}{tol}$

proof: $\|x_1 - x_2\|_2 = \|(A(tol))^{\dagger} (b_1 - b_2)\|_2$

$= \|V(\Sigma(tol))^{\dagger} U^T (b_1 - b_2)\|_2$

$= \|\operatorname{diag}(\frac{1}{\sigma_1}, \frac{1}{\sigma_2}, \dots, \frac{1}{\sigma_k}, 0, \dots, 0) U^T (b_1 - b_2)\|_2^2$

$\leq \frac{1}{\sigma_k} \|b_1 - b_2\|_2$
↑
last nonzero $\sigma_k \geq tol$

$\leq \frac{1}{tol} \|b_1 - b_2\|_2$

How does $A(tol)$ depend on tol ?

piecewise continuous, changes when $\text{tol} = 0$:

An advantage of $A(\text{tol})$:

We can use it for compression of A

$A(\text{tol})$ has rank $k \Rightarrow$ need

$m \cdot k$ words for U_1

+ k words for Σ_1

+ $n \cdot k$ words for V_1

to store SVD, vs $m \cdot n$ for full svd

$k \cdot (m + n)$ can be $\ll m \cdot n$ if k small

Solving (nearly) low rank LS

using Tikhonov regularization

or ridge regression

Replace $\arg \min_x \|Ax - b\|_2^2$

by $\arg \min_x \|Ax - b\|_2^2 + \lambda \|x\|_2^2$

you choose $\lambda > 0$

λ "penalizes" very large x

tuning parameter like tol

$\arg \min_x \|Ax - b\|_2^2 + \lambda \|x\|_2^2$

$= \arg \min_x \left\| \begin{bmatrix} A \\ \sqrt{\lambda} I \end{bmatrix} x - \begin{bmatrix} b \\ 0 \end{bmatrix} \right\|_2^2$

\uparrow
full rank $\forall \lambda > 0$

$$NE \Rightarrow x = \left(\begin{bmatrix} A \\ \sqrt{\lambda} I \end{bmatrix}^T \begin{bmatrix} A \\ \sqrt{\lambda} I \end{bmatrix} \right)^{-1} \begin{bmatrix} A \\ \sqrt{\lambda} I \end{bmatrix}^T b$$

$$(*) = (A^T A + \lambda I)^{-1} A^T b$$

\Rightarrow add λ to diagonal of NE

How does λ change SVD solution?

plug $A = U \Sigma V^T$ into (*)

$$x = V (\Sigma (\Sigma^2 + \lambda I)^{-1}) U^T b$$

$$= V \operatorname{diag} \left(\frac{\sigma_i}{\sigma_i^2 + \lambda} \right) \cdot U^T b$$

usual solution if $\lambda = 0$

$$\sigma_i \gg \lambda^{1/2} \Rightarrow \frac{\sigma_i}{\sigma_i^2 + \lambda} \sim \frac{1}{\sigma_i}$$

$$\sigma_i < \lambda^{1/2} \Rightarrow \frac{\sigma_i}{\sigma_i^2 + \lambda} \leq \frac{1}{\lambda^{1/2}}$$

i.e. $\lambda^{1/2}$ and tol in $A(\operatorname{tol})$

play analogous roles

but solution a smooth function of λ

Solving low rank LS with QR

"QR with column pivoting"

Suppose we did $A = QR$ exactly

for A : $\operatorname{rank}(A) = r < n$.

What would R look like?

If leading r columns of A were independent (true for "almost all" x)

$$R = \begin{matrix} r & n-r \\ n-r \end{matrix} \left[\begin{array}{c|c} R_{11} & R_{12} \\ \hline 0 & 0 \end{array} \right] \quad \begin{array}{l} R_{11} \text{ full rank} \\ R_{22} = 0 \end{array}$$

If A nearly low rank, hope that $\|R_{22}\| < \text{tol}$, set $R_{22} = 0$

Assuming true, solve LS as follows

$$\begin{aligned} A &= Q \cdot R = \begin{matrix} n & m-n \\ m \end{matrix} \left[\begin{array}{c|c} Q_1 & Q_2 \\ \hline Q_1' & Q_2' \end{array} \right] \begin{matrix} m \\ n \end{matrix} \left[\begin{array}{c} R \\ 0 \end{array} \right] \\ &= \begin{matrix} r & n-r & m-n \\ m \end{matrix} \left[\begin{array}{c|c|c} Q_1 & Q_2 & Q_1' \\ \hline & & \end{array} \right] \begin{matrix} R_{11} & R_{12} \\ \hline 0 & 0 \end{matrix} \end{aligned}$$

$$\operatorname{argmin}_x \|Ax - b\|_2$$

$$= \operatorname{argmin}_x \left\| \begin{bmatrix} Q_1 & Q_2 & Q_1' \\ \hline & & \end{bmatrix} \begin{bmatrix} R_{11} & R_{12} \\ \hline 0 & 0 \end{bmatrix} x - b \right\|_2$$

$$= \operatorname{argmin}_x \left\| \begin{bmatrix} R_{11} & R_{12} \\ \hline 0 & 0 \end{bmatrix} x - \begin{bmatrix} Q_1^T b \\ Q_2^T b \\ Q_1'^T b \end{bmatrix} \right\|_2$$

$$x = \begin{bmatrix} x_1 \\ \hline x_2 \end{bmatrix} \begin{matrix} r \\ n-r \end{matrix}$$

$$= \operatorname{argmin}_x \left\| \begin{bmatrix} R_{11} x_1 + R_{12} x_2 - Q_1^T b \\ \hline Q_2^T b \\ \hline Q_1'^T b \end{bmatrix} \right\|_2$$

solution $x_1 = R_{11}^{-1} Q_1^T b - R_{11}^{-1} R_{12} x_2$
for all x_2

How to pick x_2 to minimize $\|x\|_2$?
What can go wrong?

Ex $A = \begin{bmatrix} e & 1 \\ 0 & 0 \end{bmatrix}$ e tiny $R_{11} = e, R_{12} = 1$

$x = \begin{bmatrix} (b_1 - x_2)/e \\ x_2 \end{bmatrix}$ e tiny \Rightarrow
very sensitive to
small changes in x_2, b_1

$A \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} = AP = \begin{bmatrix} 1 & e \\ 0 & 0 \end{bmatrix} \Rightarrow x = \begin{bmatrix} b_1 - e x_2 \\ x_2 \end{bmatrix}$

insensitive to changes in b, x_2

What would a "perfect" R factor look like?

Compare $\begin{bmatrix} R_{11} & R_{12} \\ 0 & R_{22} \end{bmatrix}$ to $\begin{bmatrix} \Sigma_1 & 0 \\ 0 & \Sigma_2 \end{bmatrix}$

Def: Rank Revealing QR factorization
(RRQR) for short is $A \cdot P = QR$

where:

R_{11} approximates Σ_1

R_{22} " Σ_2

$\|R_{11}^{-1} R_{12}\|$ small