

Comments and Corrections

Counterexample to the Vector Generalization of Costa's Entropy Power Inequality, and Partial Resolution

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Abstract—We give a counterexample to the vector generalization of Costa's entropy power inequality due to Liu *et al.* In particular, the claimed inequality can fail if the matrix-valued parameter in the convex combination does not commute with the covariance of the additive Gaussian noise. Conversely, the inequality holds if these two matrices commute.

Index Terms—Entropy power inequality, Costa's EPI.

I. INTRODUCTION AND MAIN RESULT

For a random vector X with density on \mathbb{R}^n , let $h(X)$ denote its differential entropy. Let $Z \sim N(0, \Sigma_Z)$ be a Gaussian vector in \mathbb{R}^n independent of X , and let A be a (real symmetric) positive semidefinite $n \times n$ matrix satisfying $A \leq I$ with respect to the positive semidefinite ordering, where I denotes the identity matrix. In [1, Th. 1], Liu *et al.* claim the following generalization of Costa's EPI¹ [2]:

$$e^{\frac{2}{n}h(X+A^{1/2}Z)} \geq |I - A|^{1/n} e^{\frac{2}{n}h(X)} + |A|^{1/n} e^{\frac{2}{n}h(X+Z)}. \quad (1)$$

Liu *et al.* apply (1) in their investigation of the secrecy capacity region of the degraded vector Gaussian broadcast channel with layered confidential messages.

The purpose of this note is to demonstrate that (1) can fail for $n \geq 2$, and also to offer a partial resolution. Toward the first goal, consider $n = 2$ and let us define

$$\Sigma_X = \begin{pmatrix} 200 & 100 \\ 100 & 51 \end{pmatrix}, \quad \Sigma_Z = \begin{pmatrix} 200 & 0 \\ 0 & 1 \end{pmatrix},$$

$$A^{1/2} = \frac{1}{20} \begin{pmatrix} 10 & 5 \\ 5 & 17 \end{pmatrix}. \quad (2)$$

Taking $X \sim N(0, \Sigma_X)$ and $Z \sim N(0, \Sigma_Z)$ to be independent Gaussian vectors, we have

$$\frac{1}{2\pi e} e^{\frac{2}{n}h(X+A^{1/2}Z)} = |\Sigma_X + A^{1/2}\Sigma_Z A^{1/2}|^{1/2} \approx 19.53.$$

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¹Entropies are taken to be base e throughout. For a positive semidefinite matrix M , we write $M^{1/2}$ to denote the unique positive semidefinite matrix such $M = M^{1/2}M^{1/2}$.

On the other hand,

$$\frac{1}{2\pi e} \left(|I - A|^{1/n} e^{\frac{2}{n}h(X)} + |A|^{1/n} e^{\frac{2}{n}h(X+Z)} \right)$$

$$= |I - A|^{1/2} |\Sigma_X|^{1/2} + |A|^{1/2} |\Sigma_X + \Sigma_Z|^{1/2} \approx 40.28.$$

Thus, a contradiction to (1) is obtained. We remark that there is nothing particularly unique about this counterexample, except that the matrices were chosen to violate (1) by a significant margin.

Evidently, further assumptions are needed in order for (1) to hold. To give a simple resolution, we note that it suffices for the matrices A and Σ_Z to commute.

Theorem 1: Let X be a random vector with density on \mathbb{R}^n whose entropy exists in the usual Lebesgue sense, and let $Z \sim N(0, \Sigma_Z)$ be a Gaussian vector in \mathbb{R}^n independent of X . If $A \leq I$ is positive semidefinite and commutes with Σ_Z , then

$$e^{\frac{2}{n}h(X+A^{1/2}Z)} \geq |I - A|^{1/n} e^{\frac{2}{n}h(X)} + |A|^{1/n} e^{\frac{2}{n}h(X+Z)}.$$

Proof: For brevity, we refer the reader to the original proof of [1, Th. 1], and only point out where the argument needs to be corrected. To this end, Liu *et al.*'s proof contains an incorrect application of the AM-GM inequality in the form [1, eq. (28)]:

$$|\Sigma_Z^{-1} \text{Cov}(Z|D_\gamma X + Z)(I - D_\gamma^{-2})|^{1/n}$$

$$\leq \frac{1}{n} \text{Tr}(\Sigma_Z^{-1} \text{Cov}(Z|D_\gamma X + Z)(I - D_\gamma^{-2})), \quad (3)$$

where $D_\gamma := (I + \gamma(A - I))^{1/2}$, and $\gamma \in [0, 1]$ parameterizes a path of perturbation. Indeed, a product of positive semidefinite matrices is not necessarily positive semidefinite, which can lead to failure of the AM-GM inequality in the form (3). For example, returning to the counterexample above where $X \sim N(0, \Sigma_X)$ and matrices are chosen according to (2), the eigenvalues of $\Sigma_Z^{-1} \text{Cov}(Z|D_\gamma X + Z)(I - D_\gamma^{-2})$ can be approximately computed as $\{-0.0053, -0.7273\}$ for $\gamma = 0.5$, in violation of (3).

However, if A and Σ_Z commute, then so do $\Sigma_Z^{-1/2}$ and $(I - D_\gamma^{-2})^{1/2}$ since real symmetric matrices commute if and only if they are simultaneously diagonalizable by some orthogonal matrix U . Hence,

$$\text{Tr}(\Sigma_Z^{-1} \text{Cov}(Z|D_\gamma X + Z)(I - D_\gamma^{-2}))$$

$$= \text{Tr}((I - D_\gamma^{-2})^{1/2} \Sigma_Z^{-1/2} \text{Cov}(Z|D_\gamma X + Z) \Sigma_Z^{-1/2} (I - D_\gamma^{-2})^{1/2}).$$

The argument of the second trace term is clearly positive semidefinite, and therefore (3) holds for all $\gamma \in [0, 1]$ under the additional assumption that A and Σ_Z commute, thereby repairing Liu *et al.*'s proof. \square

II. REMARKS

The critical application of inequality (1) by Liu *et al.* (see [1, p. 1877]) assumes only that A and Σ_Z are diagonal matrices, so the conclusions of [1] appear to be unaffected, aside from [1, Th. 1]. In fact, Ekrem and Ulukus [3] obtained the same secrecy capacity results using a different argument. Nevertheless, [1] has been cited numerous times in the literature, so other published results may be affected to varying degrees. As one example, a computation in [4] similarly overlooks non-commutativity of the matrices A and Σ_Z , leading to the incorrect conclusion that [1, Th. 1] is a corollary of [4, Th. 3]. Despite the error, the validity of [4, Th. 3] is not impacted, and the corrected computation, showing that [4, Th. 3] implies Theorem 1 above, can be found in [5, Sec. III.A]. Other works may be more seriously affected, but we do not attempt to give an accounting of consequences here.

In closing, we remark that the additional assumption that A and Σ_Z commute is a relatively strong one. It can easily be seen, using the simultaneous diagonalization property of A and Σ_Z by a common orthogonal matrix, that Theorem 1 has a completely equivalent statement where $Z \sim N(0, I)$ and A is restricted to be a diagonal matrix with diagonal entries $0 \leq a_i \leq 1$, $i = 1, \dots, n$. Theorem 1 should be viewed as an extension of Costa's original 1985 result (which assumed identical parameters a_i) in this sense. As pointed out to the authors by an anonymous referee, a standard information-theoretic argument may be used to establish Theorem 1 as a corollary of Costa's original inequality, without appealing to the perturbation framework used in [1]. This argument has been included in the appendix.

APPENDIX

What follows is a proof of Theorem 1 using Costa's entropy power inequality [2]. In particular, we shall prove Theorem 1 in the equivalent setting noted in the closing statement above, where $Z \sim N(0, I)$ and A is a diagonal matrix with diagonal entries $0 \leq a_i \leq 1$, $i = 1, \dots, n$. The argument below was provided by an anonymous referee.

$$Y_1 := X + A^{1/2}Z_1, \text{ and } Y_2 := Y_1 + (I - A)^{1/2}Z_2,$$

where Z_1, Z_2 are independent copies of $Z \sim N(0, I)$, also independent of X . Since A is assumed diagonal with entries a_1, \dots, a_n , Theorem 1 may thus be written as

$$\prod_{i=1}^n (1 - a_i)^{1/n} e^{\frac{2}{n}(h(X) - h(Y_1))} + \prod_{i=1}^n a_i^{1/n} e^{\frac{2}{n}(h(Y_2) - h(Y_1))} \leq 1. \quad (4)$$

We consider the exponential terms separately. By the Csiszár sum identity,

$$\begin{aligned} h(X) - h(Y_1) &= \sum_{i=1}^n \left(h(X_i | X^{i-1}, Y_{1,i+1}^n) \right. \\ &\quad \left. - h(Y_{1,i} | X^{i-1}, Y_{1,i+1}^n) \right) \\ &= \sum_{i=1}^n \left(h(X_i | V_i) - h(Y_{1,i} | V_i) \right), \end{aligned}$$

where $V_i := (X^{i-1}, Y_{1,i+1}^n)$. Similarly,

$$\begin{aligned} h(Y_2) - h(Y_1) &= \sum_{i=1}^n \left(h(Y_{2,i} | Y_2^{i-1}, Y_{1,i+1}^n) \right. \\ &\quad \left. - h(Y_{1,i} | Y_2^{i-1}, Y_{1,i+1}^n) \right) \\ &\leq \sum_{i=1}^n \left(h(Y_{2,i} | V_i) - h(Y_{1,i} | V_i) \right), \end{aligned}$$

where the last inequality follows from

$$I(X^{i-1}; Y_{2,i} | Y_2^{i-1}, Y_{1,i+1}^n) \leq I(X^{i-1}; Y_{1,i} | Y_2^{i-1}, Y_{1,i+1}^n),$$

which holds by Markovity induced by construction of Y_1, Y_2 . Hence, (4) holds if

$$\begin{aligned} e^{\frac{1}{n} \sum_{i=1}^n \log(1 - a_i) + 2 h(X_i | V_i) - 2 h(Y_{1,i} | V_i)} \\ + e^{\frac{1}{n} \sum_{i=1}^n \log(a_i) + 2 h(Y_{2,i} | V_i) - 2 h(Y_{1,i} | V_i)} \leq 1. \end{aligned}$$

By convexity of $x \mapsto e^x$, it is sufficient to show that

$$(1 - a_i)e^{2 h(X_i | V_i) - 2 h(Y_{1,i} | V_i)} + a_i e^{2 h(Y_{2,i} | V_i) - 2 h(Y_{1,i} | V_i)} \leq 1$$

for each i . However, this is just the conditional form of Costa's scalar inequality for concavity of entropy power, which holds again by convexity of $x \mapsto e^x$, and the easily verified fact that $V_i \rightarrow X_i \rightarrow Y_{1,i} \rightarrow Y_{2,i}$.

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